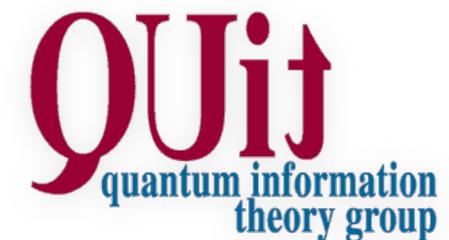


Operational probabilistic theories and cellular automata: how I learned to stop worrying and love C^* algebras

School on Advanced Topics in Quantum Information and Foundations

Quantum Information Unit and the Yukawa Institute for Theoretical Physics, Kyoto University



Paolo Perinotti - February 8-12 2021

Lecture 2

Quantum Cellular automata

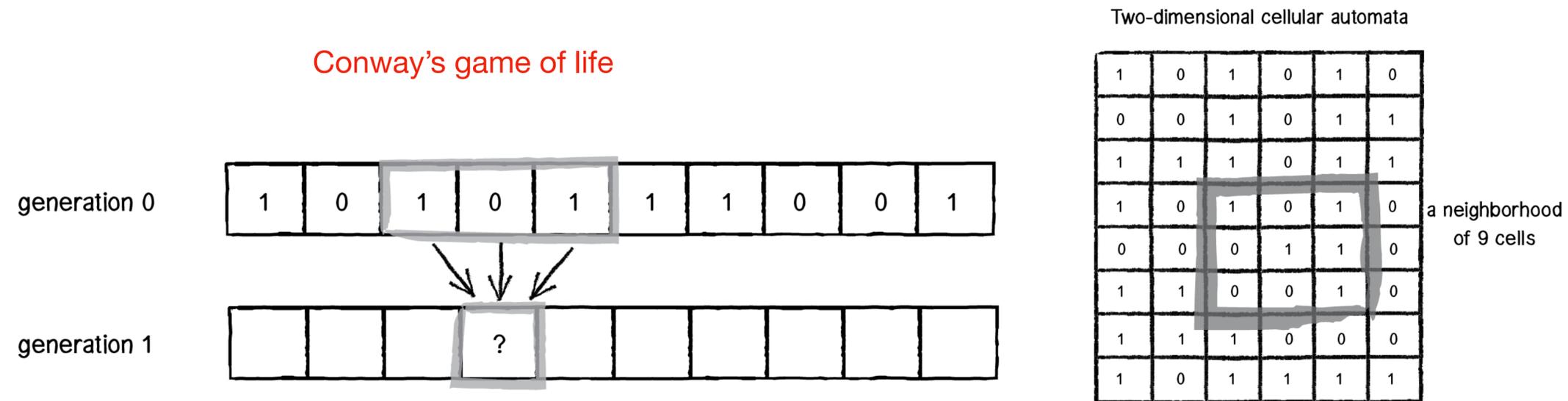
Summary

- Cellular Automata
- QCA
 - finite quantum systems
 - Infinite systems - inductive limit
 - Infinite systems - topological closure
- Neighbourhood
- The problem of quantisation
- Margolus decompositions
- Fermionic CA and quantum walks



Cellular Automata

J. Von Neumann and A. W. Burks, "Theory of self-reproducing automata" 1966

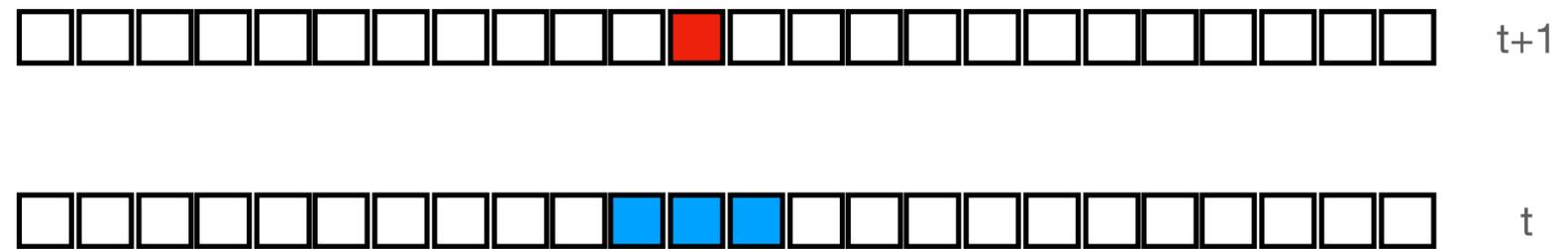


Definition (Cellular Automata). A cellular automaton (CA) is a 4-tuple $(L, \Sigma, \mathcal{N}, f)$ consisting of (1) a d -dimensional lattice of cells L indexed $i \in \mathbb{Z}^d$, (2) a finite set of states Σ , (3) a finite neighborhood scheme $\mathcal{N} \subset \mathbb{Z}^d$, and (4) a local transition function $f : \Sigma^{\mathcal{N}} \rightarrow \Sigma$.

Quantum cellular automata

Problems

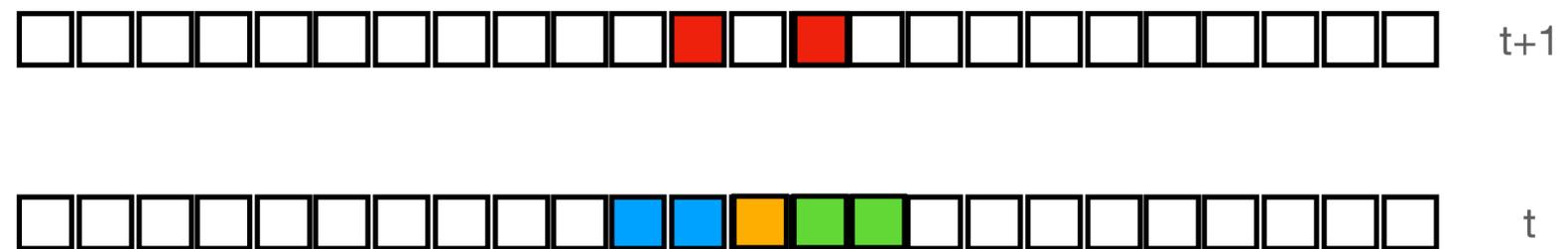
- What is a local update rule?



Quantum cellular automata

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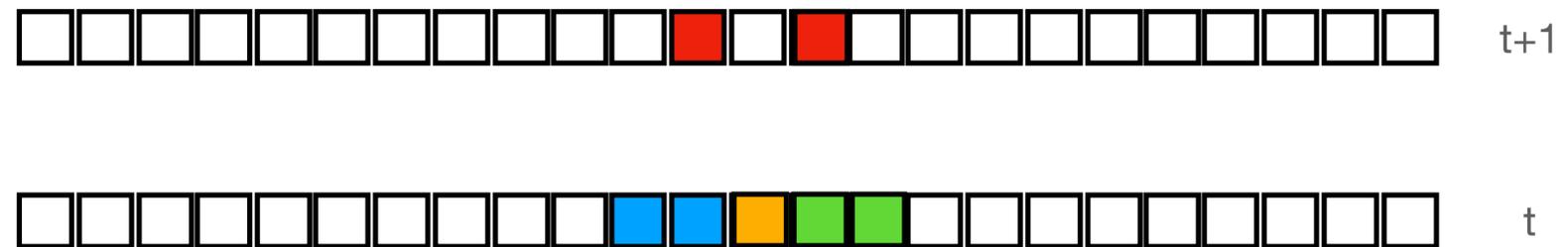
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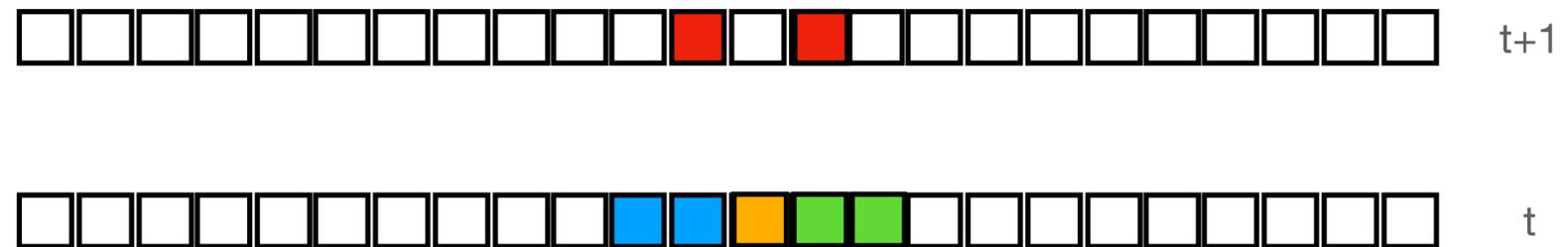
- What is a local update rule?
- No-cloning theorem
- Non commutativity of local operations



Quantum cellular automata

Problems

- What is a local update rule?
- No-cloning theorem
- Non commutativity of local operations
- Heart of the problem: definition through action on states

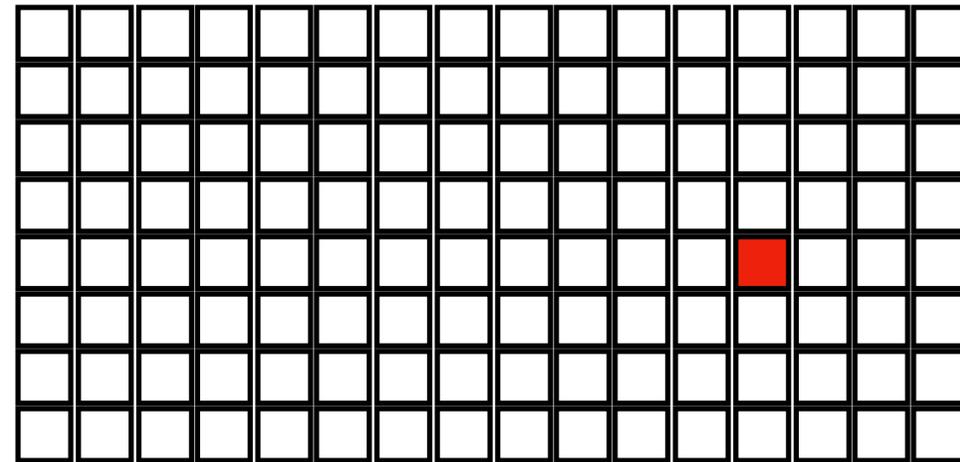


Quantum cellular automata

Finite case

$$\blacksquare = \mathcal{H}_x$$

$$\text{C.A.: } U : \bigotimes_x \mathcal{H}_x \rightarrow \bigotimes_x \mathcal{H}_x$$

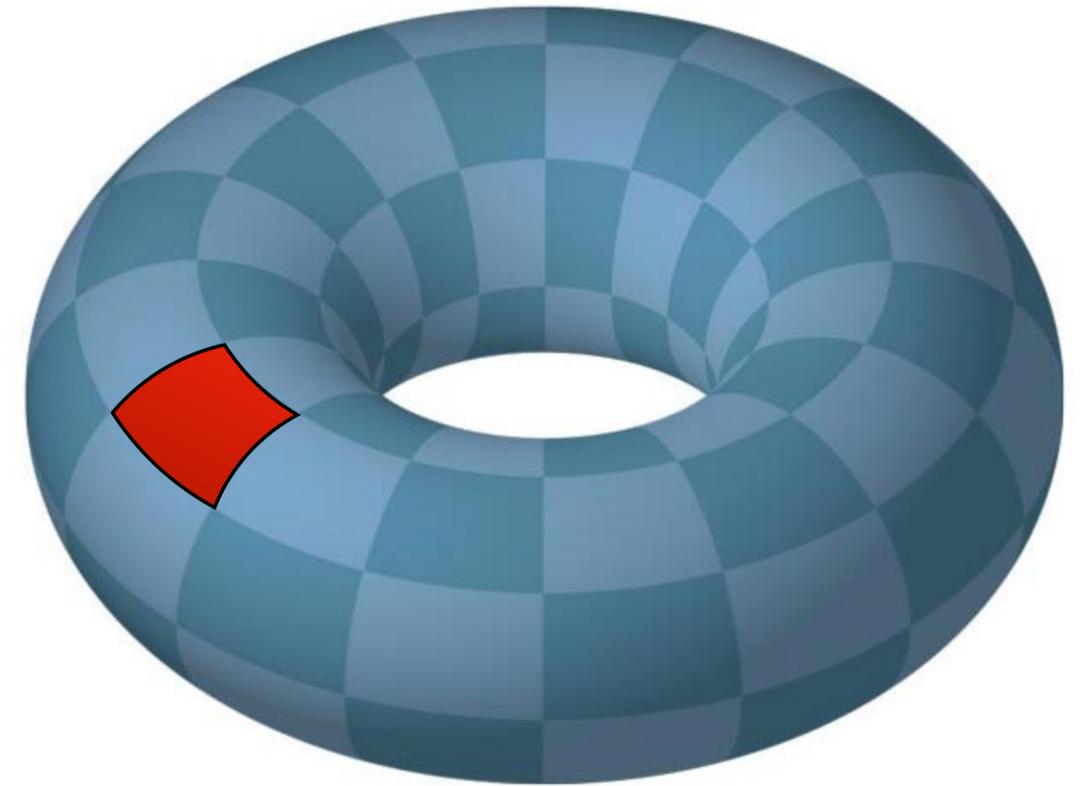
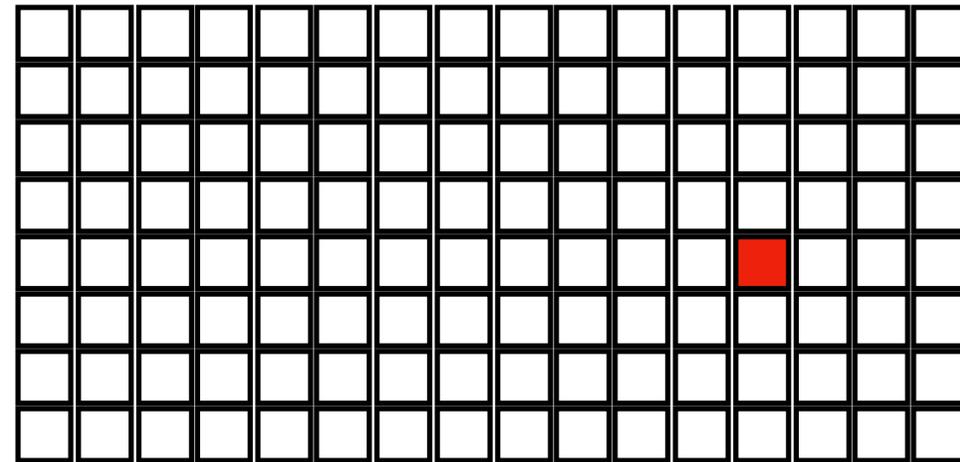


Quantum cellular automata

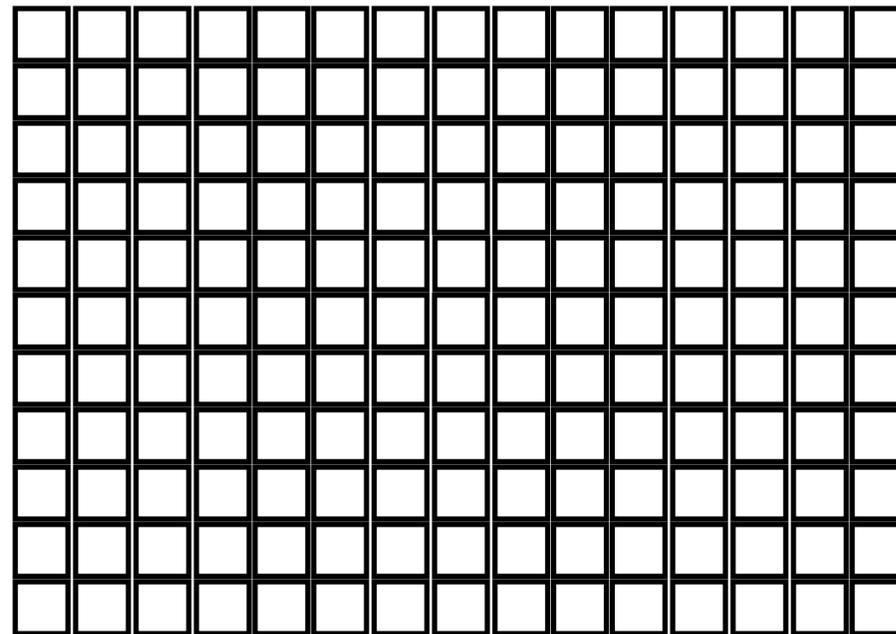
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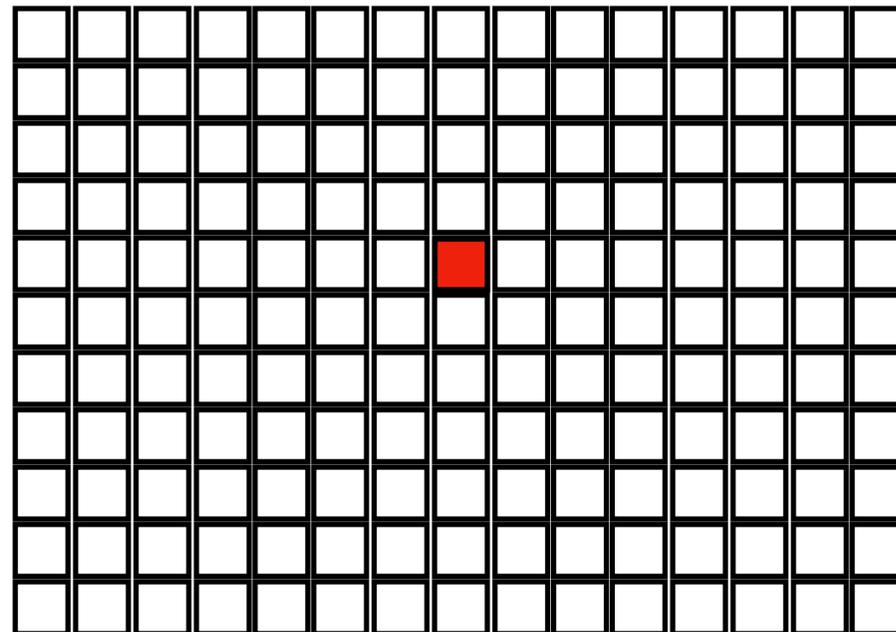


Local system and local rule



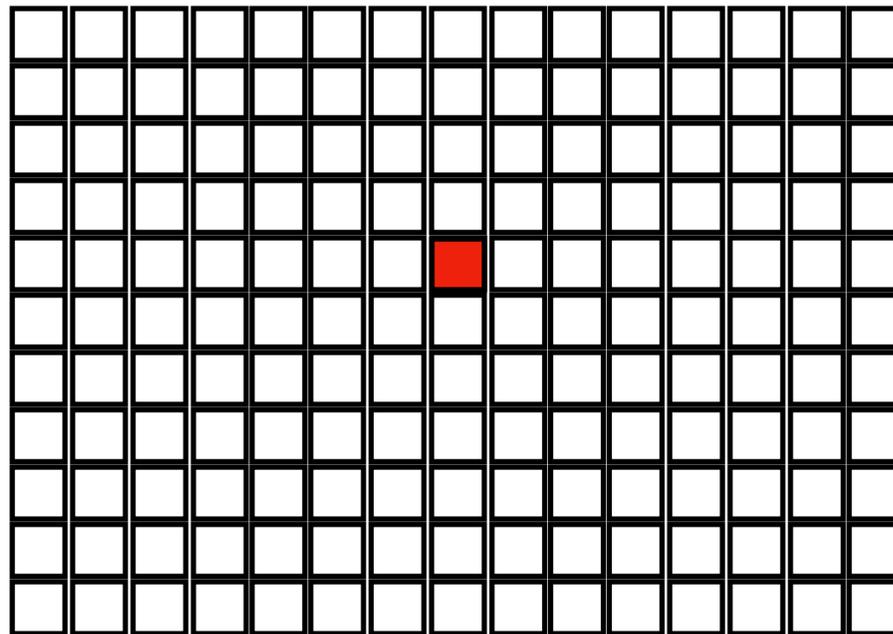
Local system and local rule

$$p(A_x = a \mid \rho) = \text{Tr}[P(a)_x \rho] = \text{Tr}[\{P(a) \otimes I_{\bar{x}}\} \rho]$$



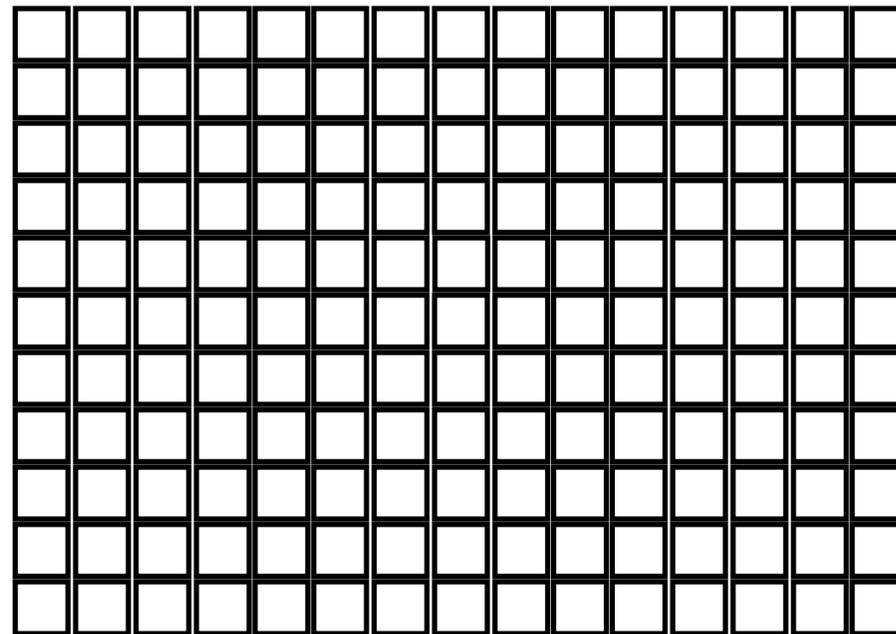
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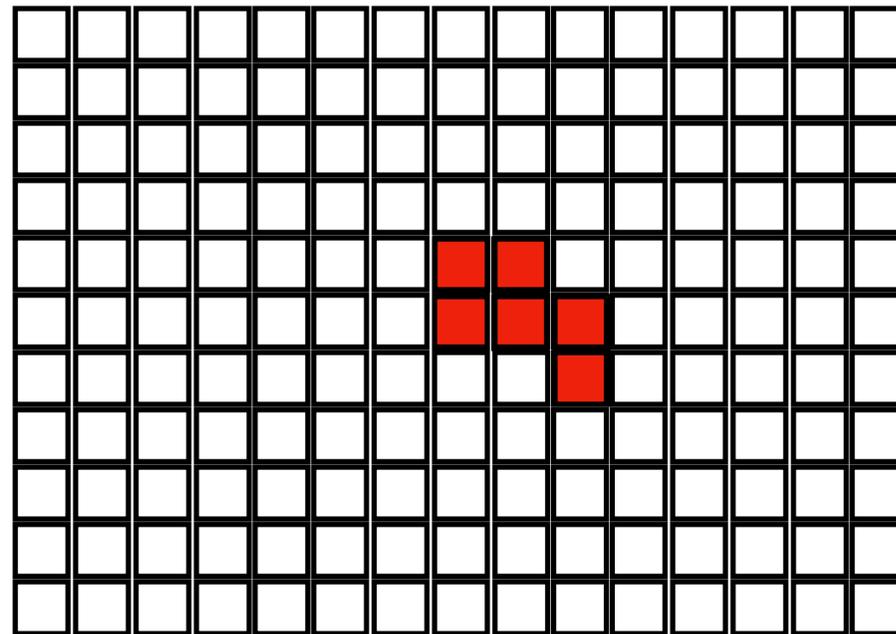
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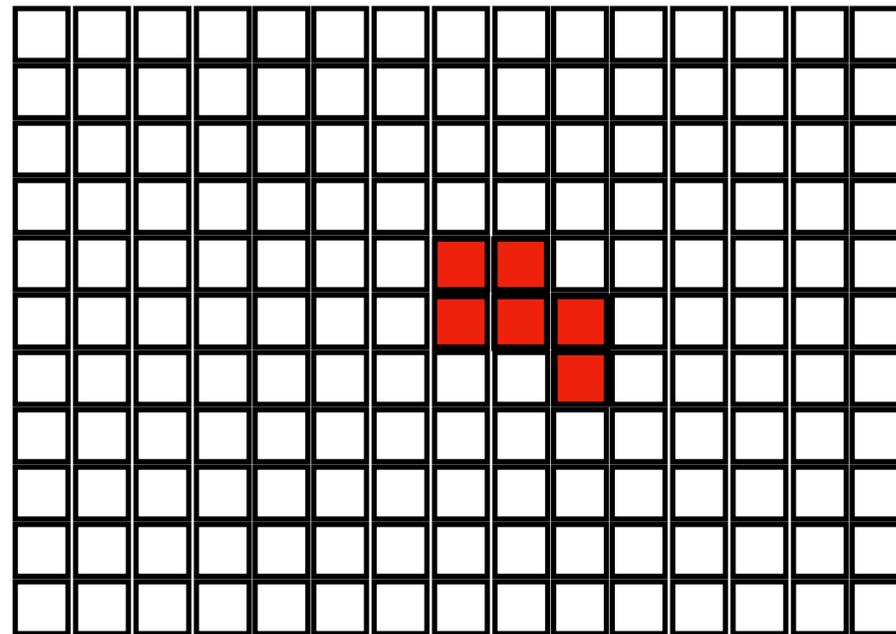
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$$A_R \in \text{Span} \left(\bigotimes_{x \in R} A_x \right)$$

Heisenberg picture

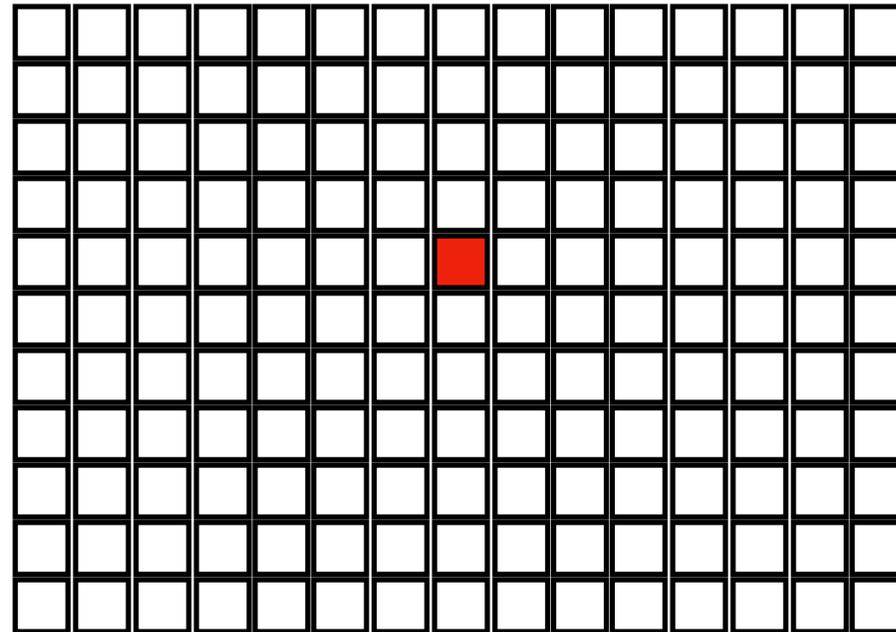
$$p(A_x = a \mid U\rho U^\dagger) = \text{Tr}[P(a)_x U\rho U^\dagger] = \text{Tr}[U^\dagger P(a)_x U\rho]$$

Quantum cellular automata

Neighbourhood of a cell

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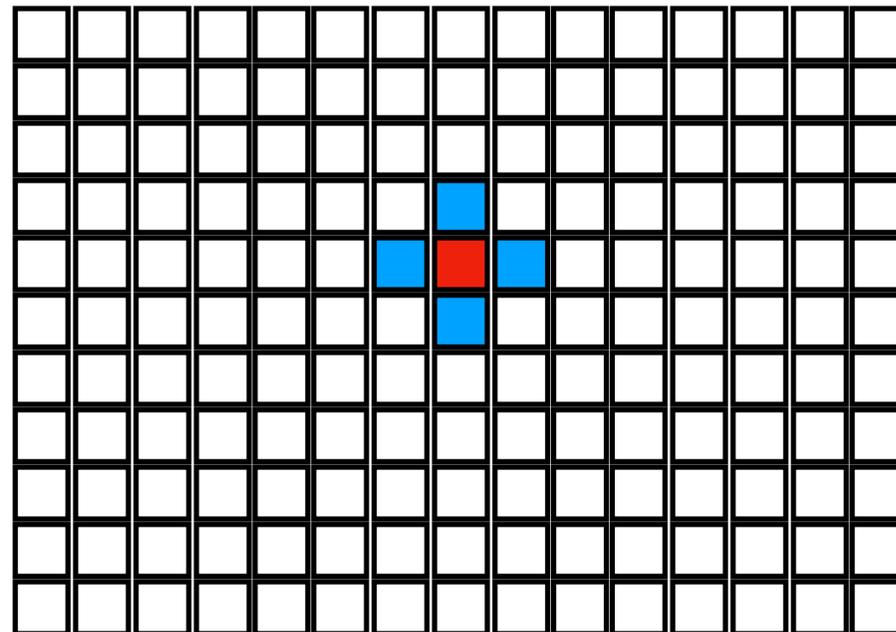
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$$U^{-1} A_{x_0} U = A_{N-(x_0)} \otimes I_{\bar{x}_0}$$



Quantum cellular automata

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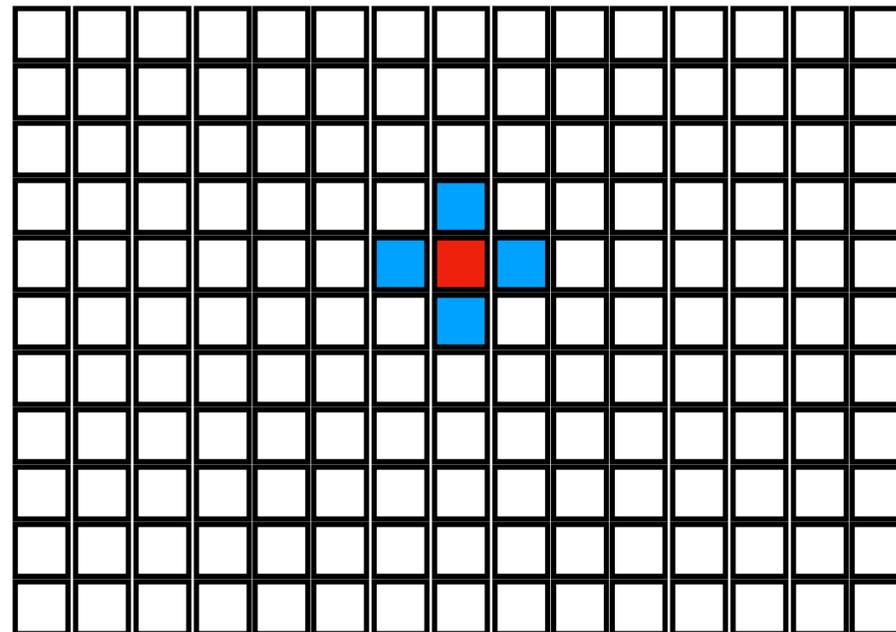
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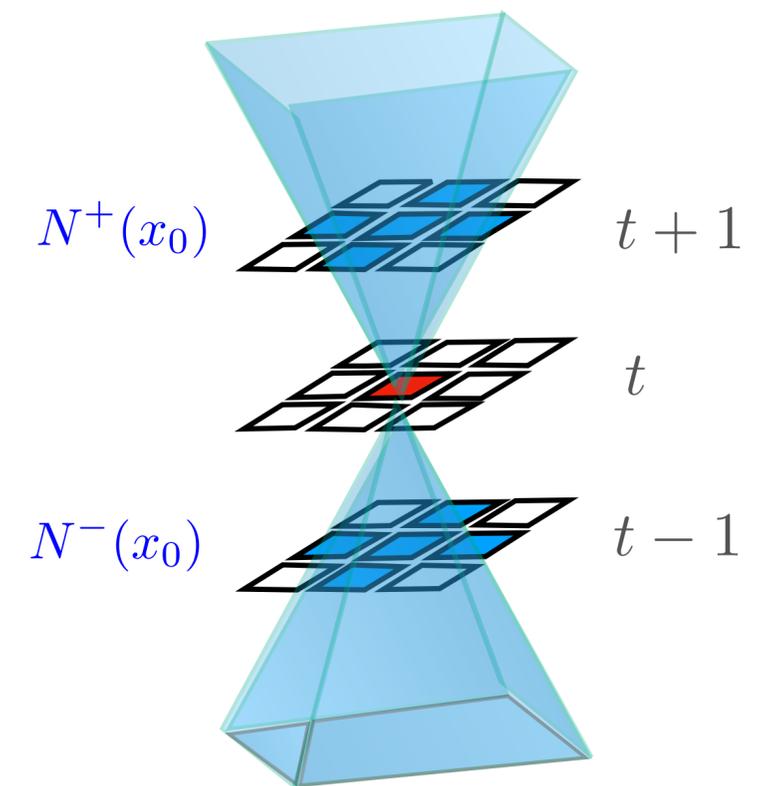
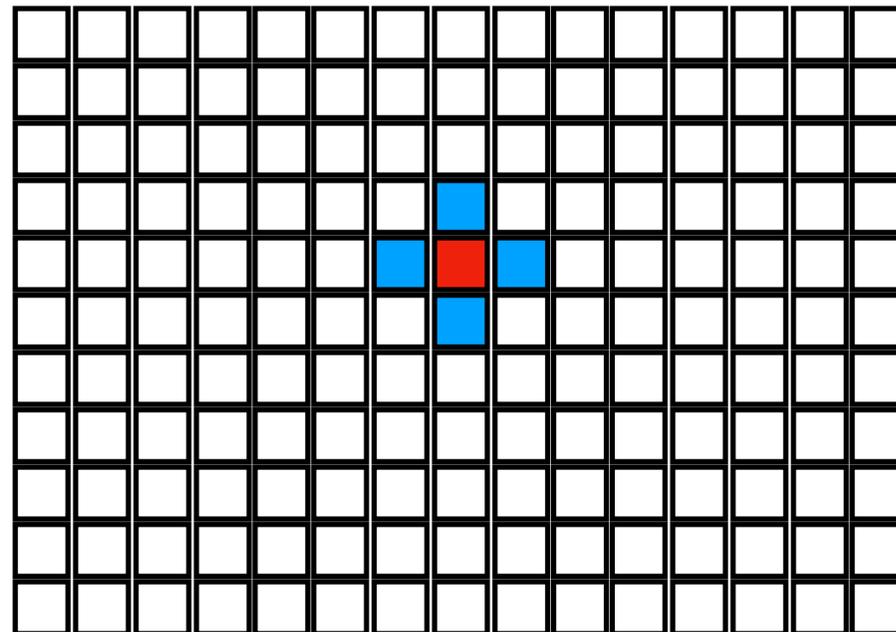
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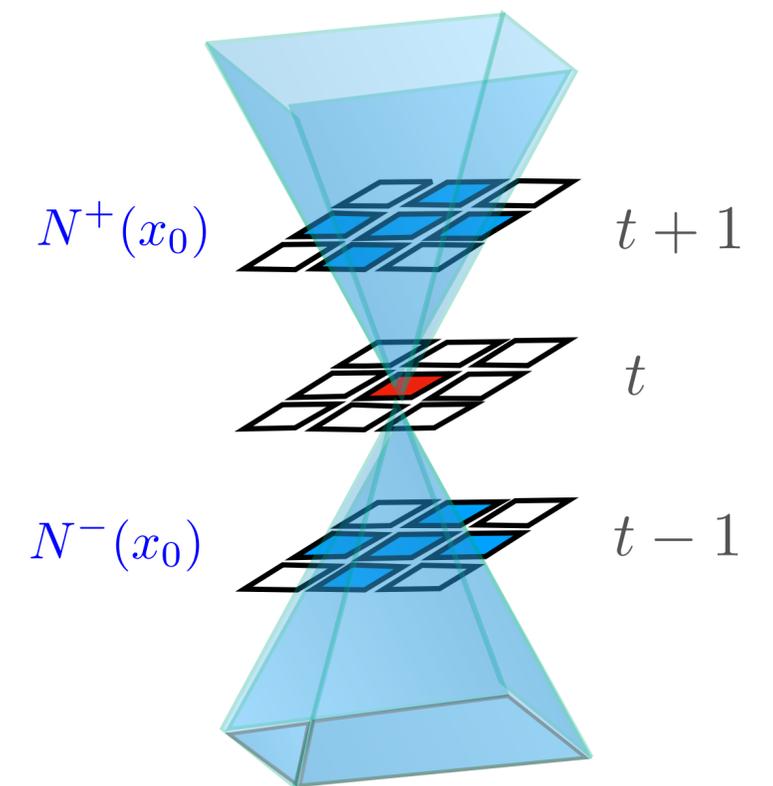
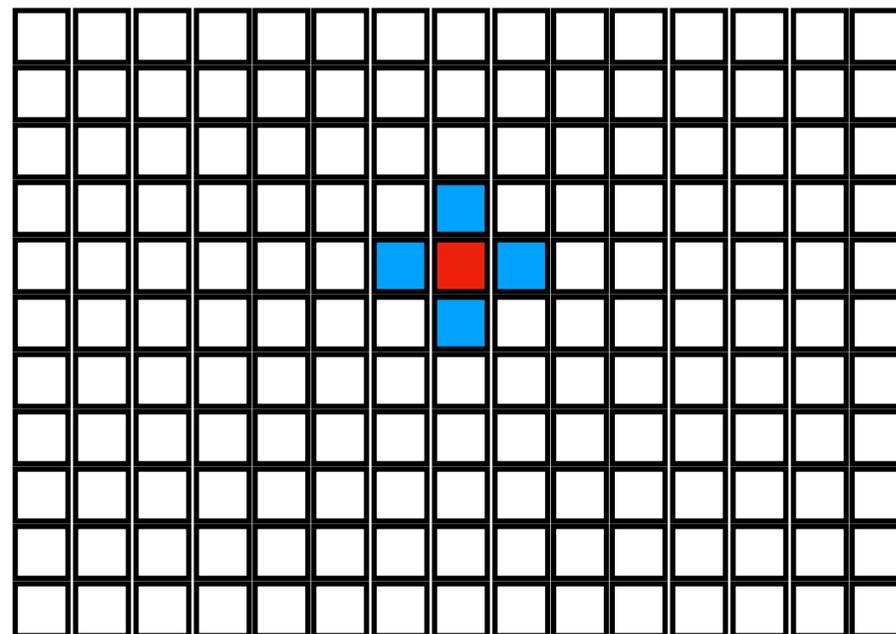
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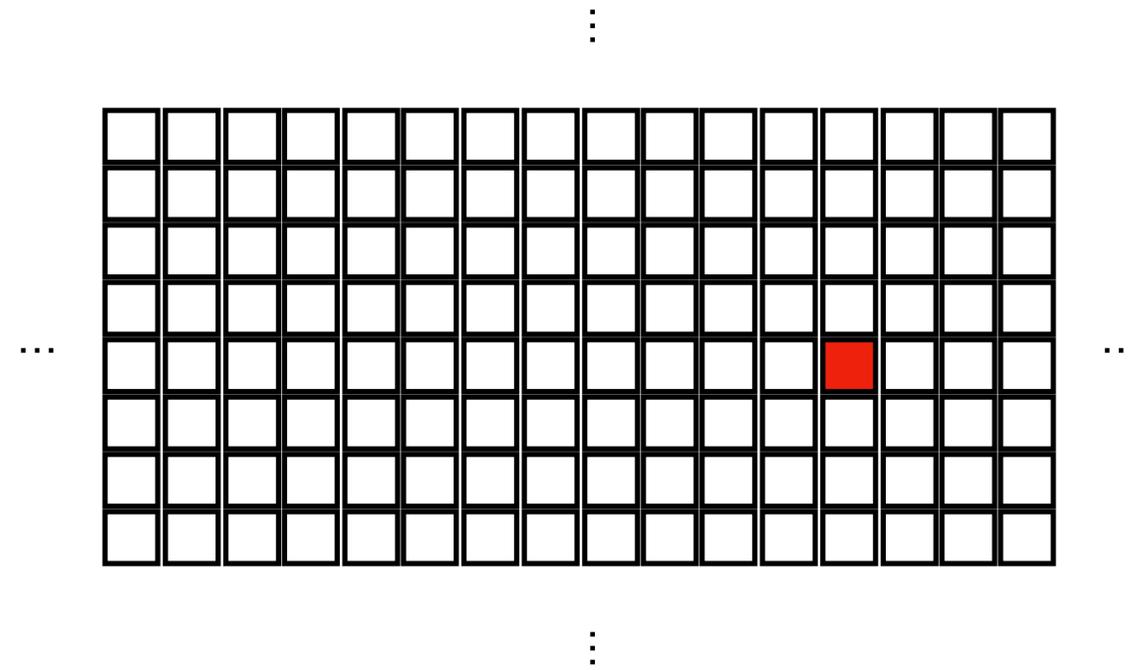


More precisely, $N^-(x_0)$ is the smallest set of cells such that $U^{-1} A_{x_0} U = A_{N^-(x_0)} \otimes I_{\bar{x}_0}$

Quantum cellular automata

Infinite case

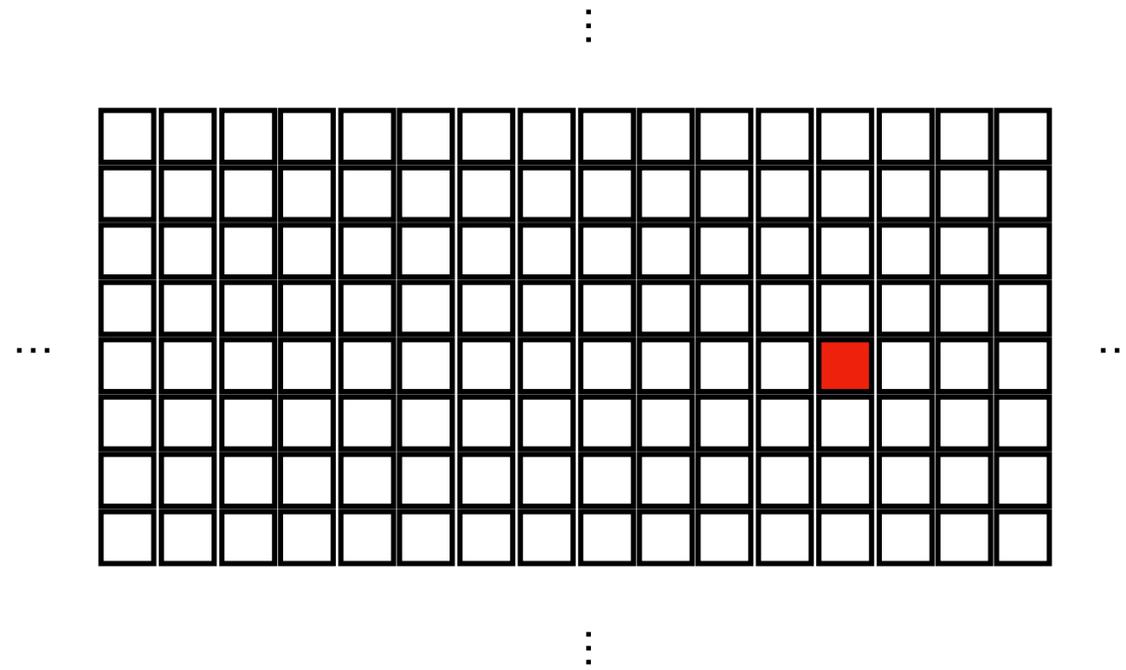
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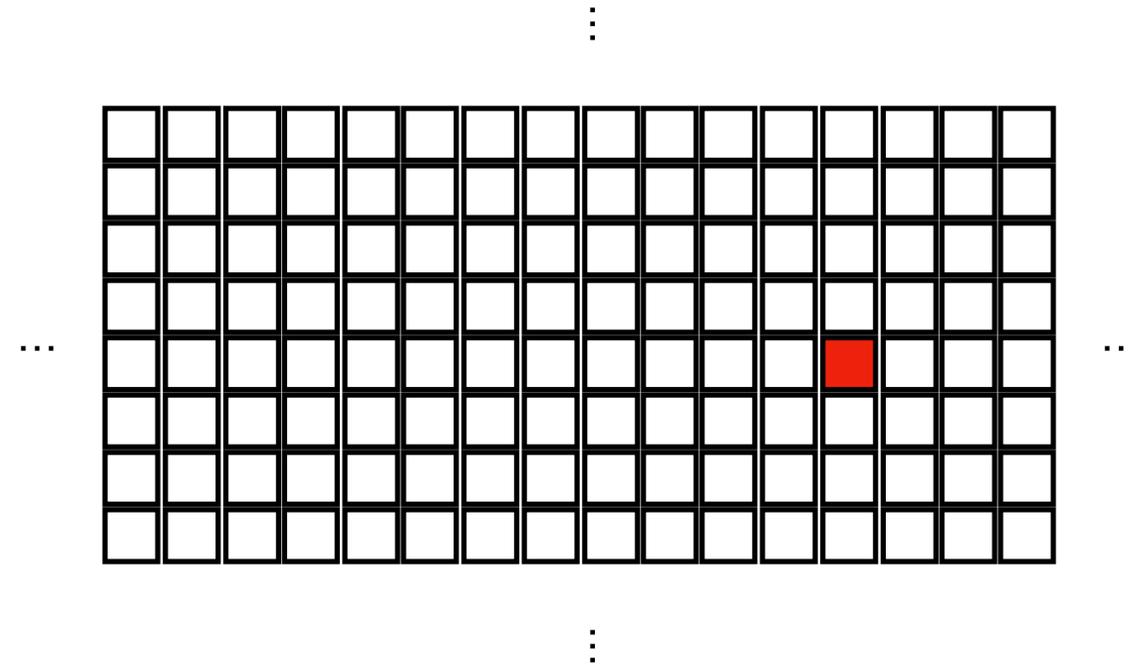


$$A = \bigotimes_x A_x$$

Infinite system

Inductive limit

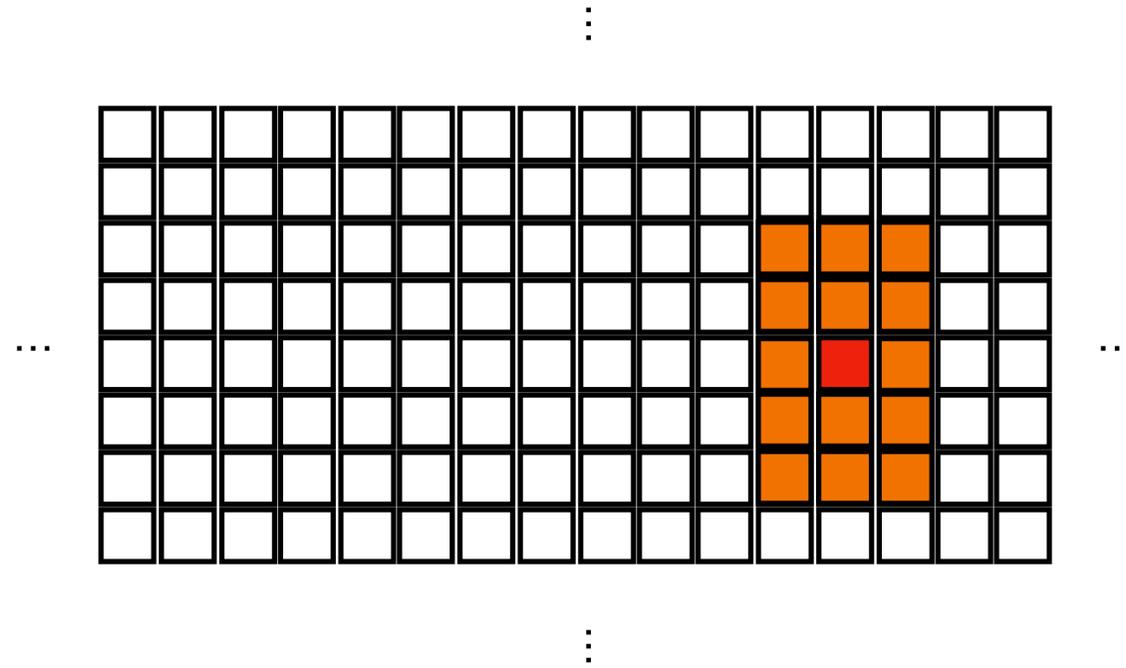
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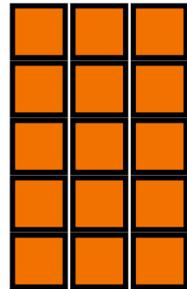
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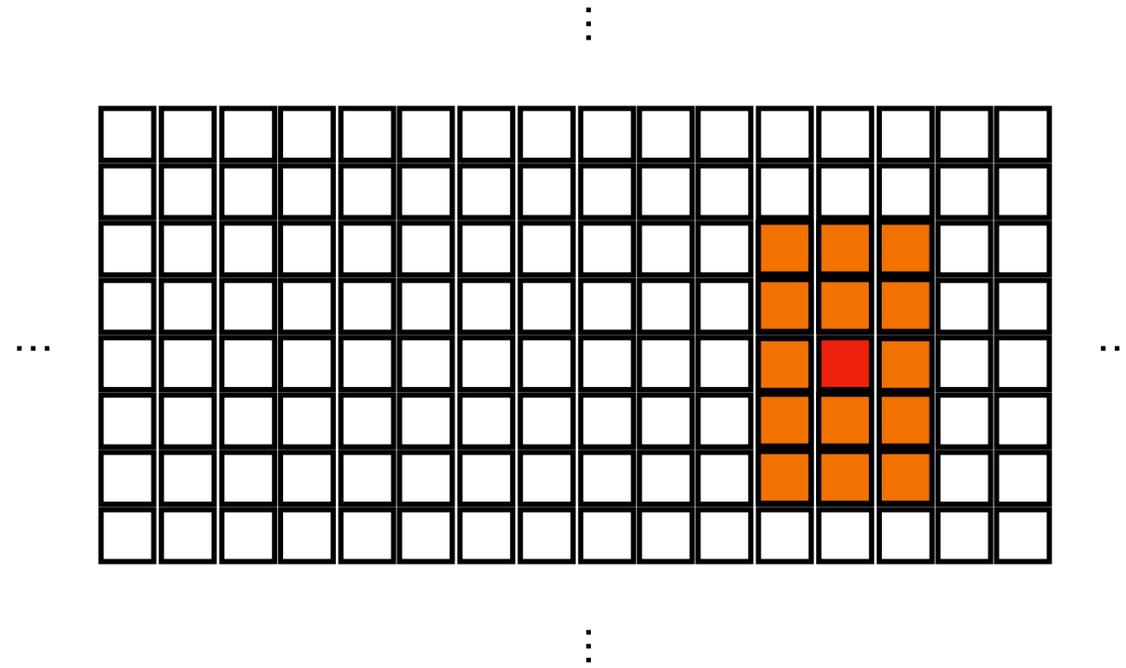
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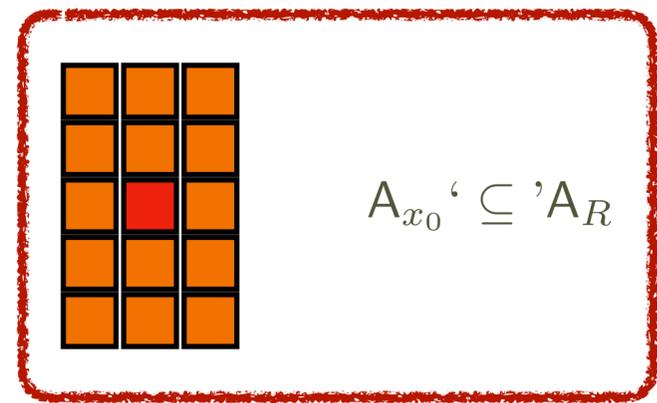
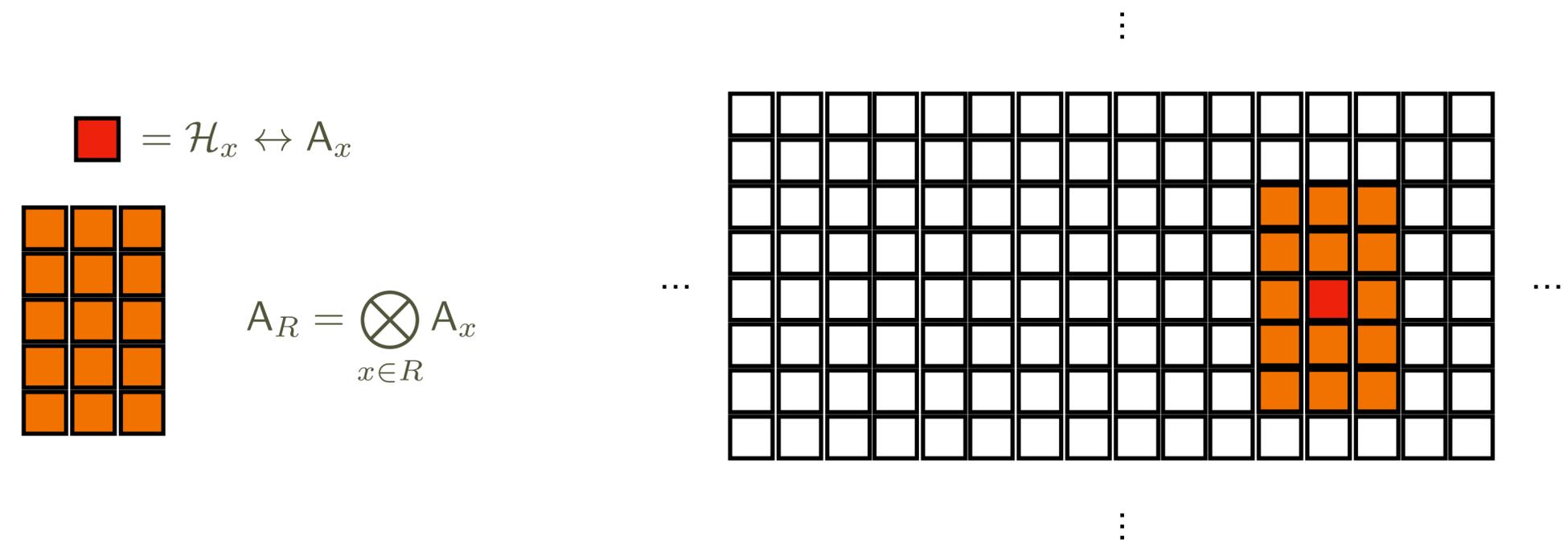


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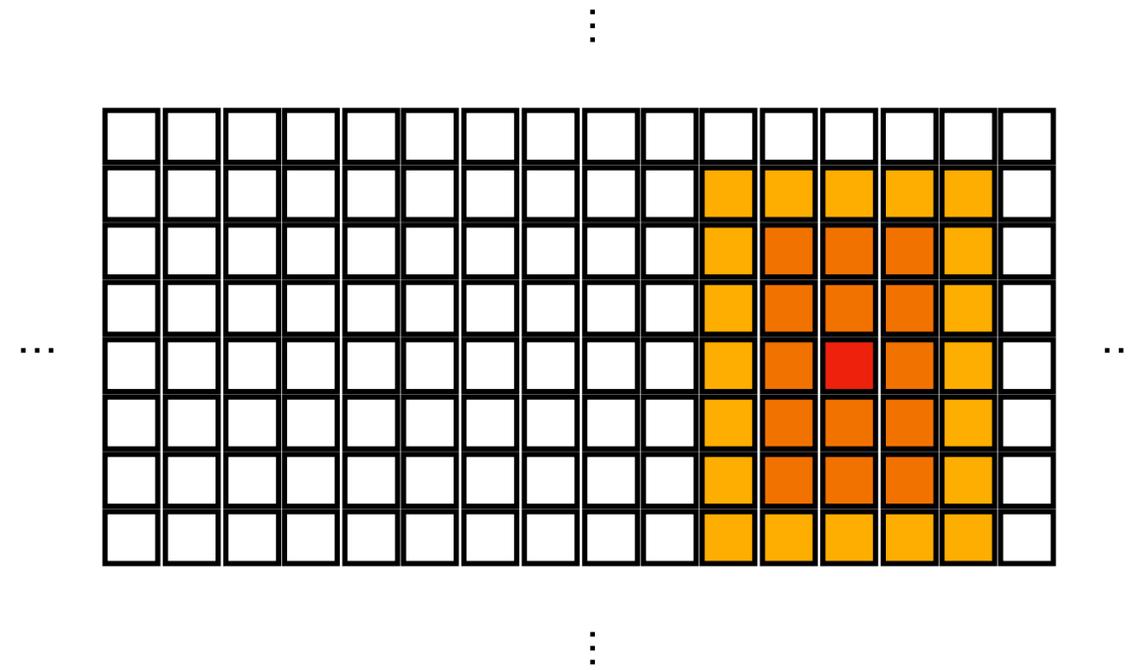
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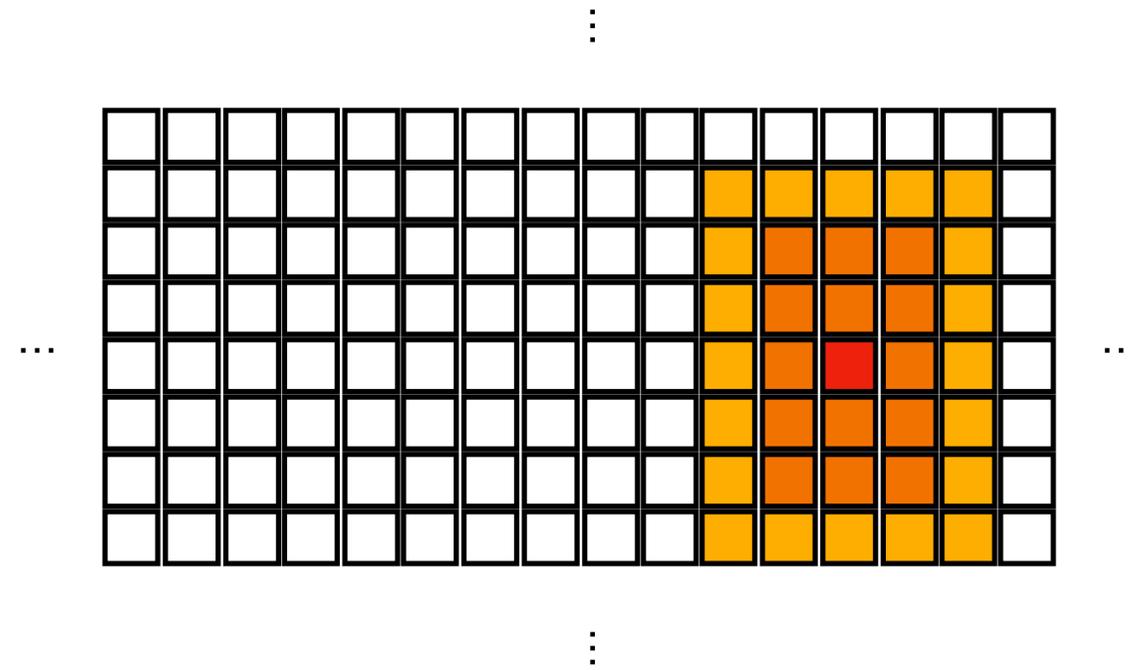
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$$R_0 \subseteq R_1 \subseteq R_2 \subseteq \dots$$

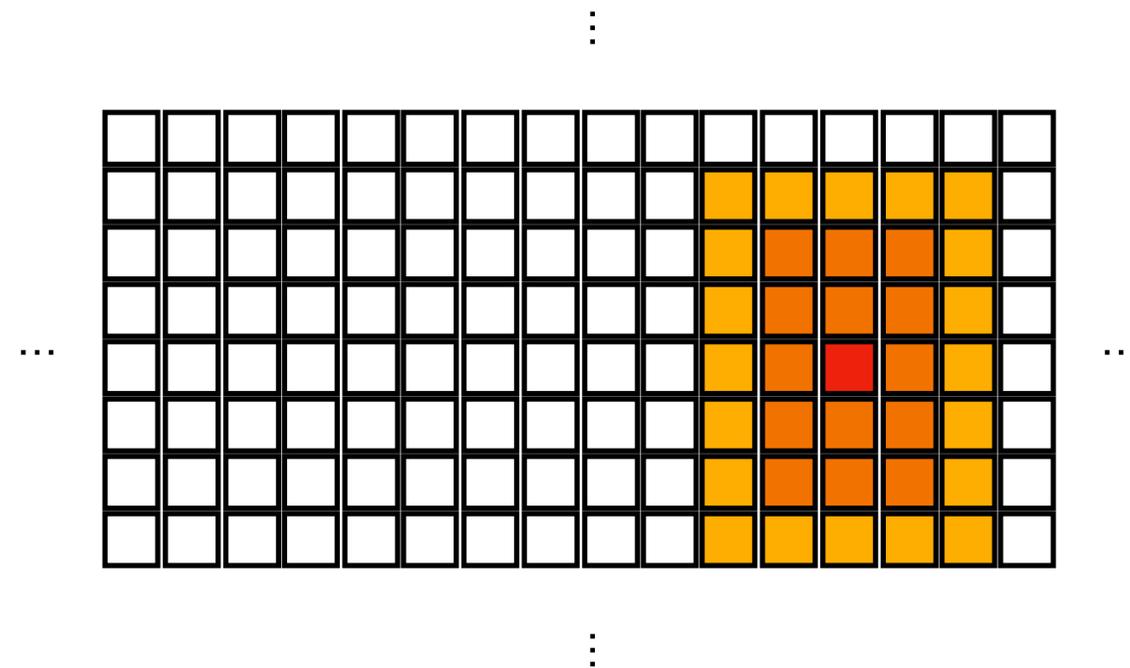
$$A_{R_0} \subseteq A_{R_1} \subseteq A_{R_2} \subseteq \dots$$

$$f_{ij} : A_{R_i} \rightarrow A_{R_j}$$

$$f_{ij}(B) = B \otimes I_{R_j \setminus R_i}$$

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$$B_j \sim B_i \text{ if } \exists k \geq i, j \text{ s.t. } f_{ik}(B_i) = f_{jk}(B_j)$$

$$A_L := \bigsqcup_i A_{R_i} / \sim$$

Infinite system

Topological limit

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- A is the **quasi-local algebra**

Quasi-local algebra

Topological limit

- The algebra A contains the local algebra as a subalgebra

$$A_L \subseteq A \quad B_i \in A_L \Leftrightarrow B_{i_n} \in A, \forall n \in \mathbb{N} \quad B_{i_n} = B_i$$

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- Define the adjoint of $B = \lim_{n \rightarrow \infty} B_{i_n}$ as

$$B^\dagger := \lim_{n \rightarrow \infty} B_{i_n}^\dagger$$

Quasi-local algebra

C^* -algebra structure

- It is easily proved that A is a c^* -algebra, i.e. a Banach vector space with

Quasi-local algebra

C*-algebra structure

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$$(AB)C = A(BC)$$

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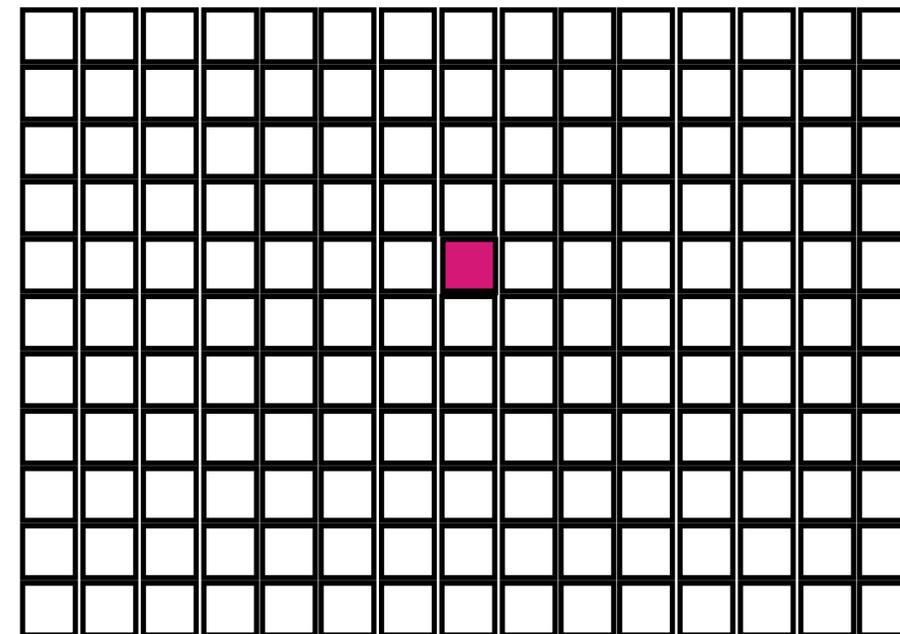
$$(AB)^* = B^*A^*$$

- and the norm satisfies $\|AB\| \leq \|A\| \|B\|$, $\|A^*A\| = \|A\|^2$

Quantum Cellular Automaton

Infinite case

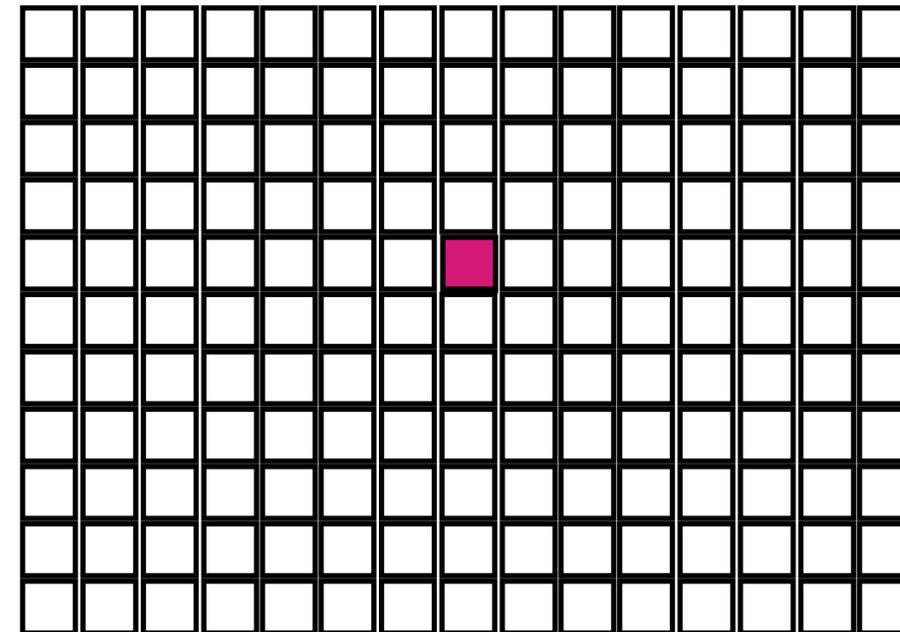
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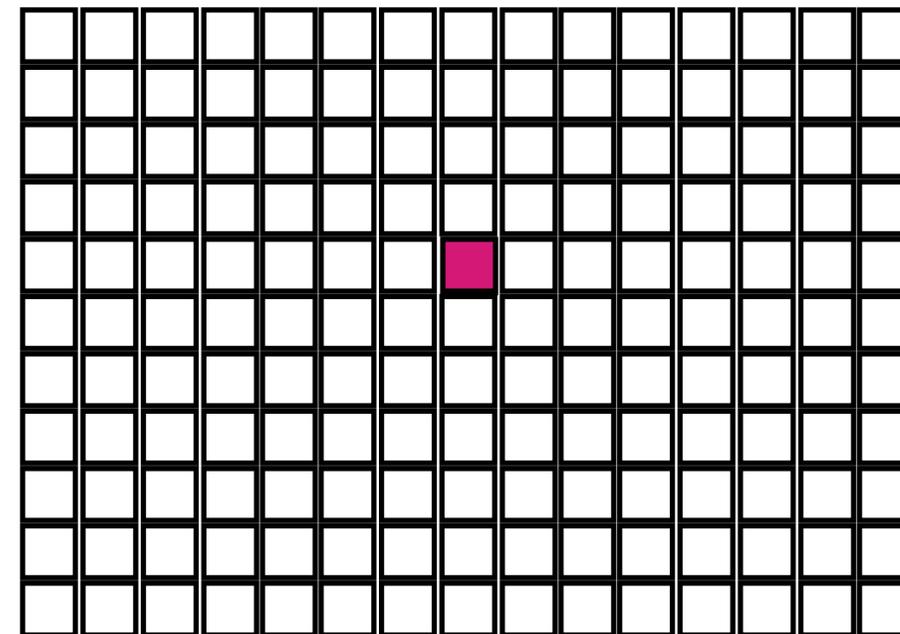
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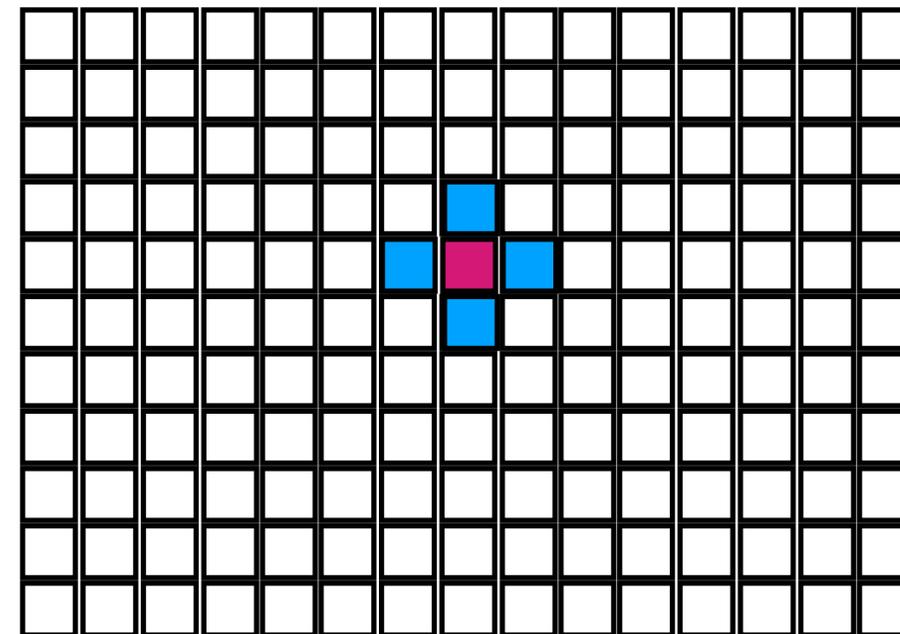
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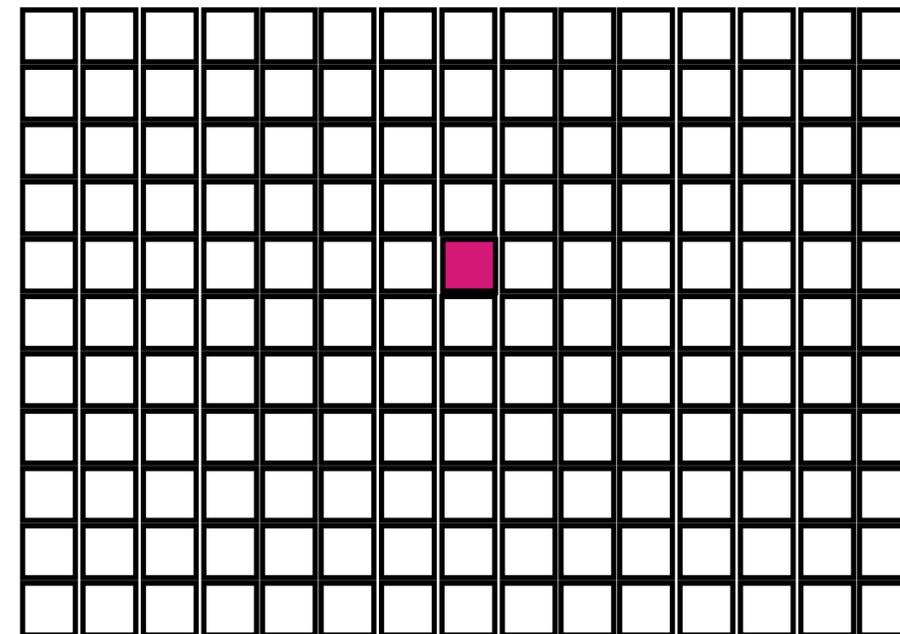
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- Translationally invariant

$$\mathcal{V}[\mathcal{T}_z(A_x)] = \mathcal{T}_z[\mathcal{V}(A_x)]$$



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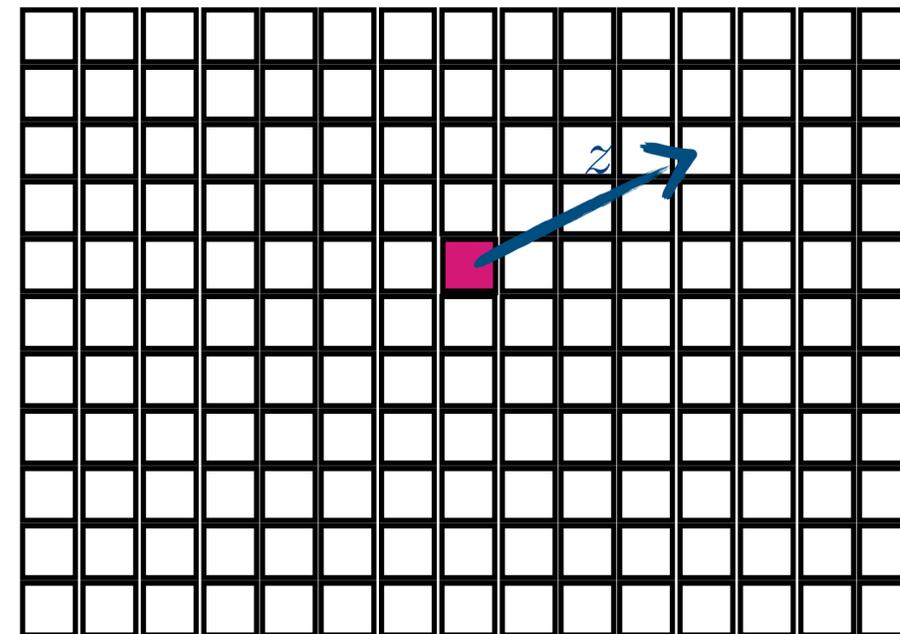
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Quantum Cellular Automaton

Infinite case

- The following definition was given for QCA on \mathbb{Z}^d
 - A QCA is an homomorphism \mathcal{V} of the quasi local algebra \mathcal{A} such that

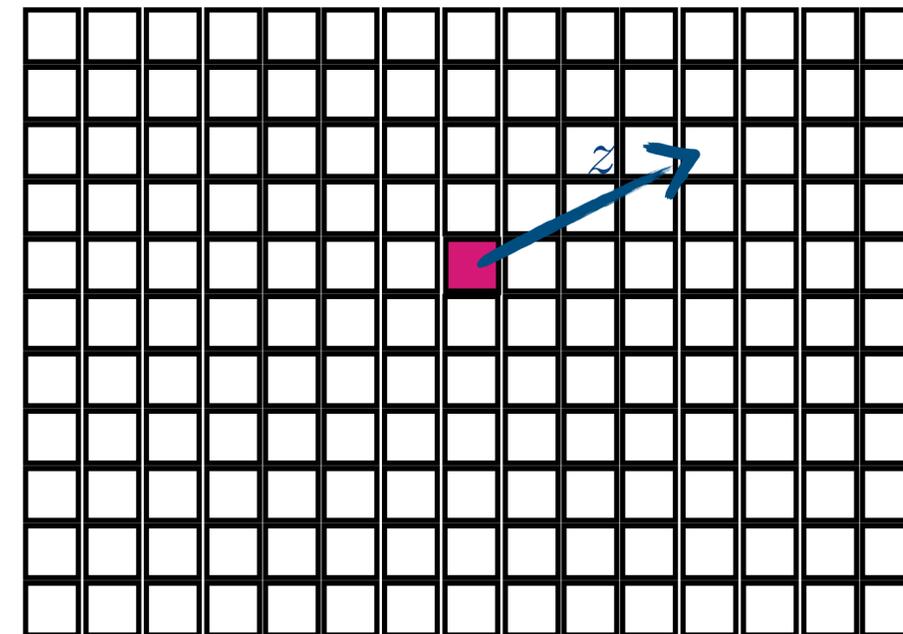
- \mathcal{V} satisfies locality

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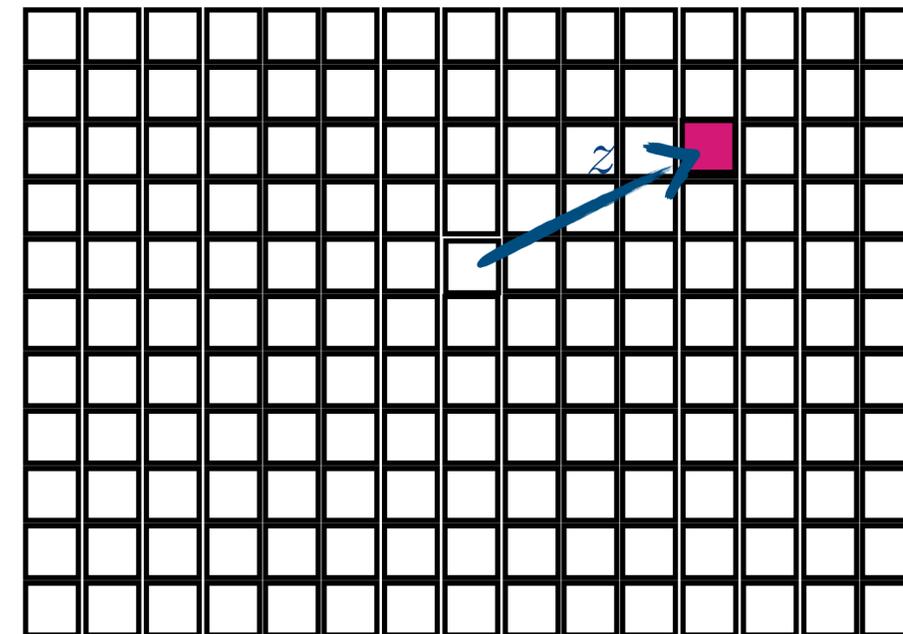
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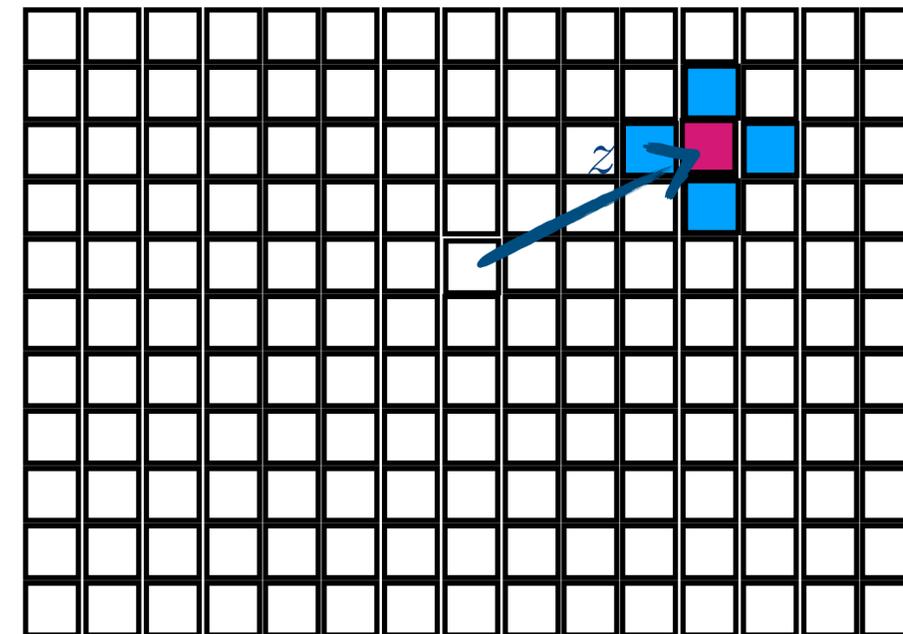
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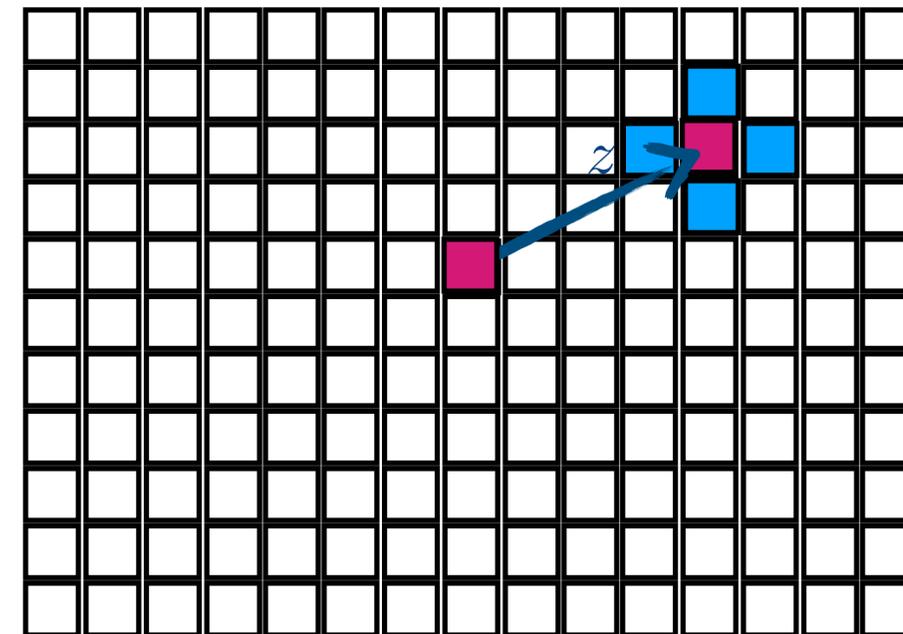
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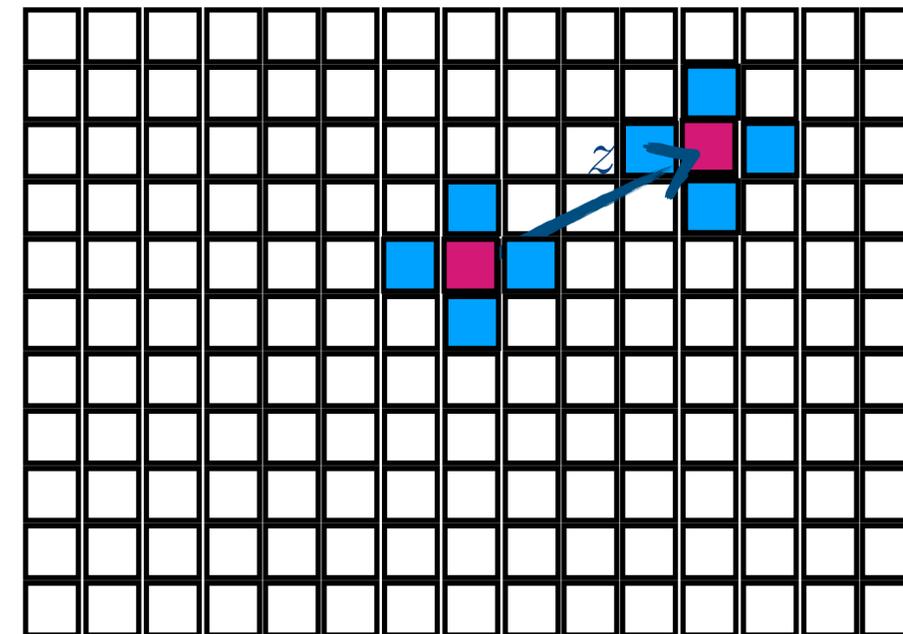
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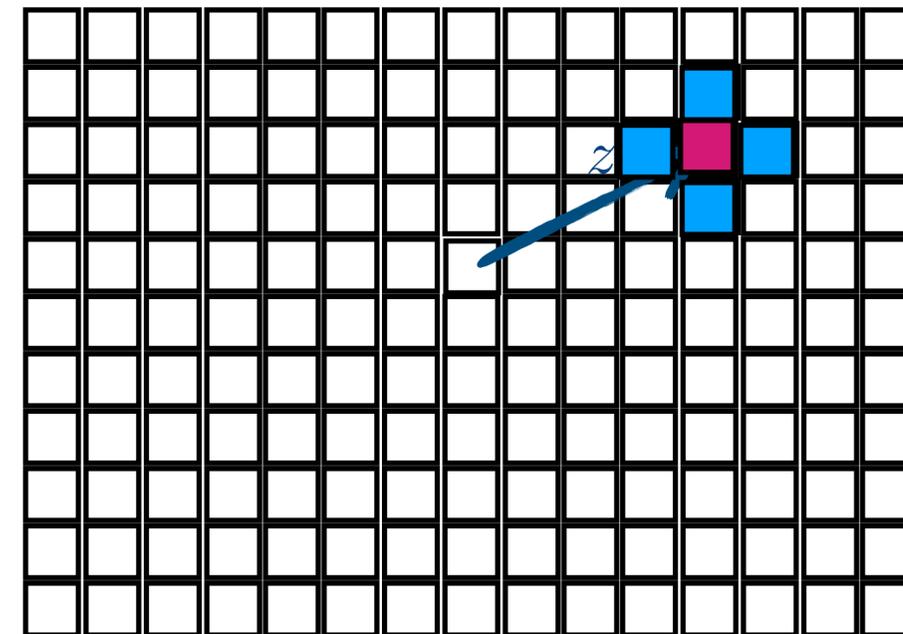
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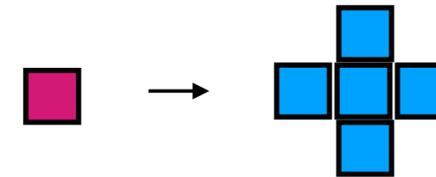


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Local update rule

Consequence of the definition

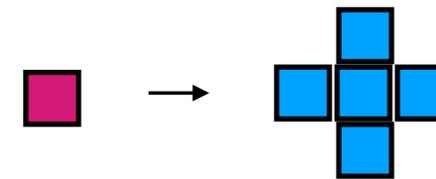
- The QCA induces a local rule $\mathcal{V}_x : A_x \rightarrow A_{N+(x)}$
 $\mathcal{V}_x(B_x) := \mathcal{V}(B_x)$



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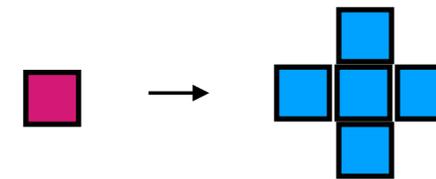
- The local rule at $x=0$ completely specifies the QCA

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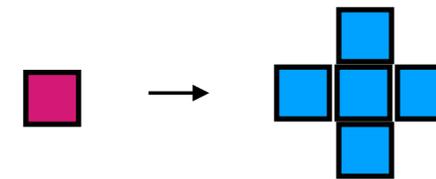
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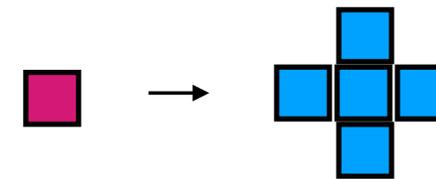
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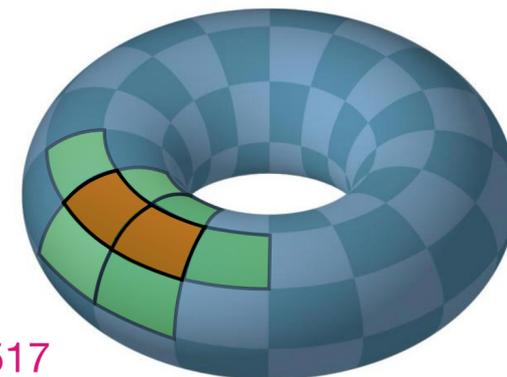
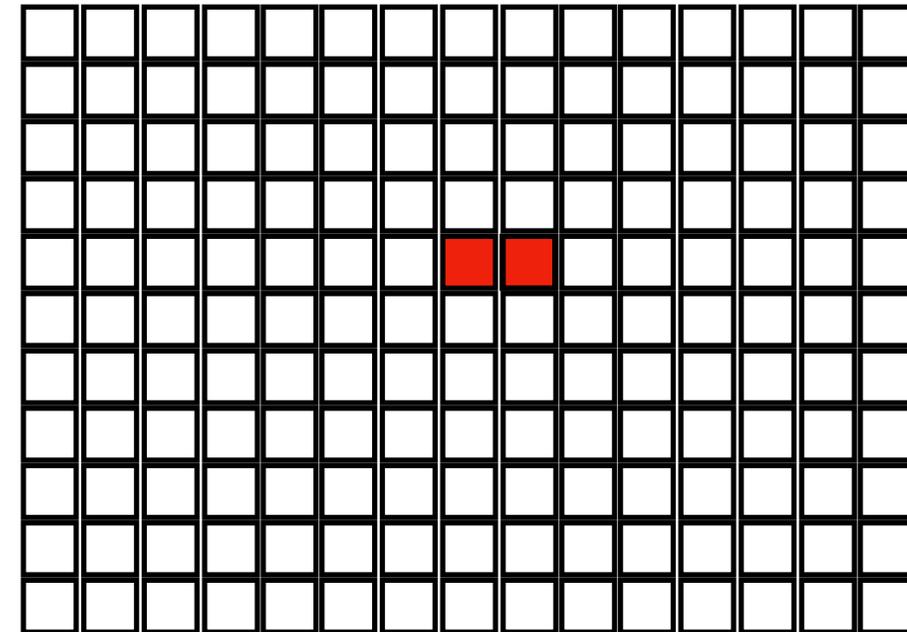
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- The image of every **local** effect can be obtained using \mathcal{V}_0 since $B_R = \sum_i \bigotimes_{x \in R} B_x^{(i)}$
- Boundedness: image of limits (of the **quasi-local algebra**) is determined

Wrapping lemma

Reduces the local rule to that of a finite QCA

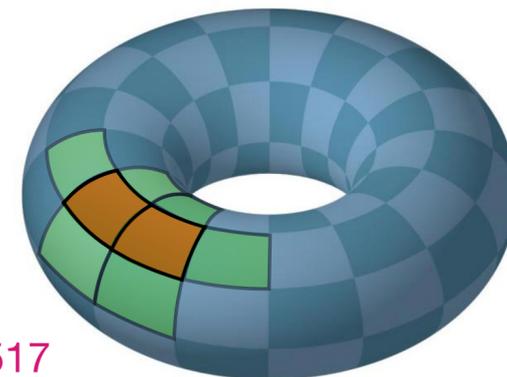
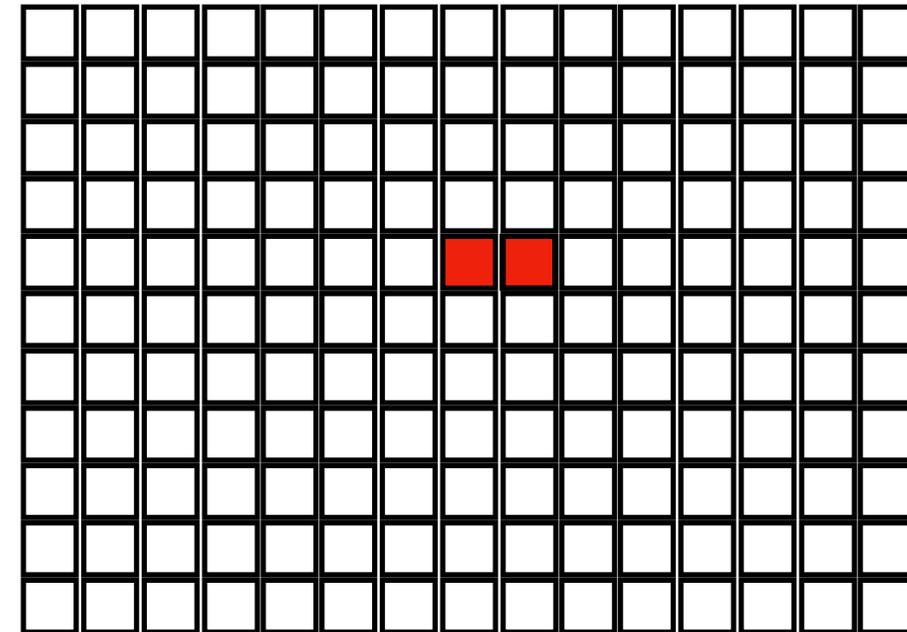
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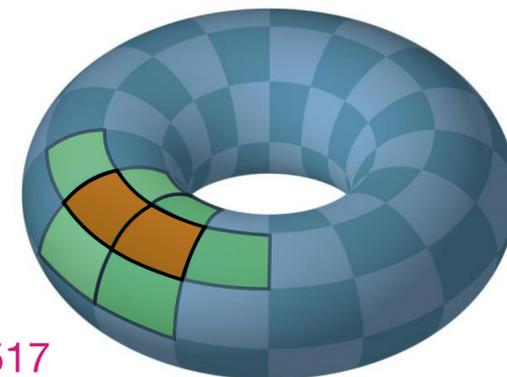
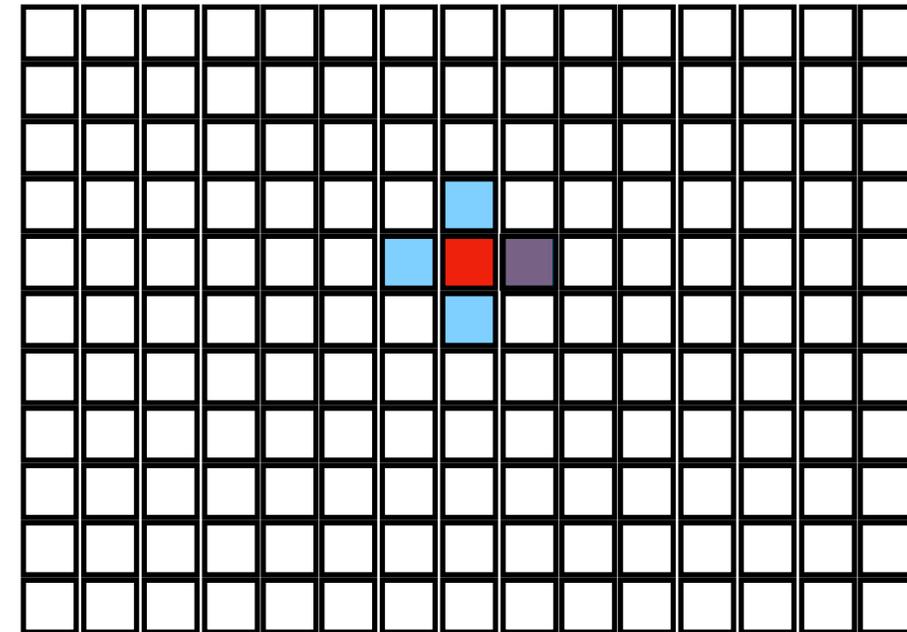
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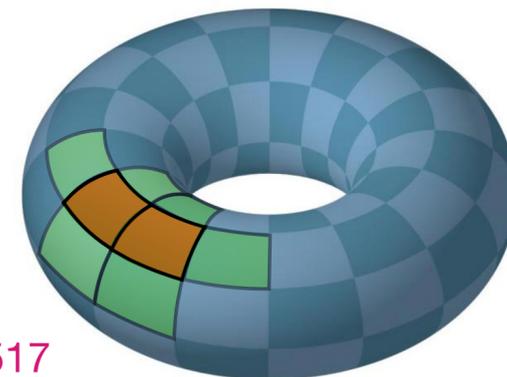
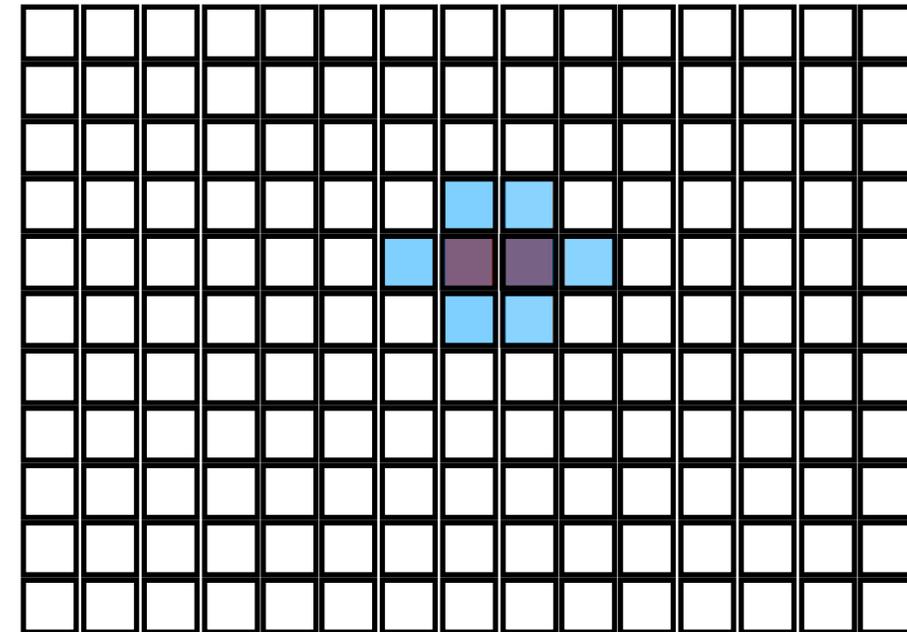
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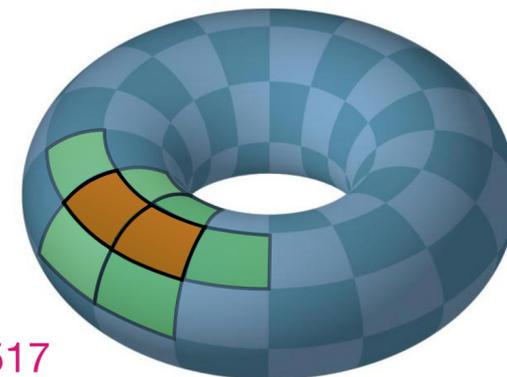
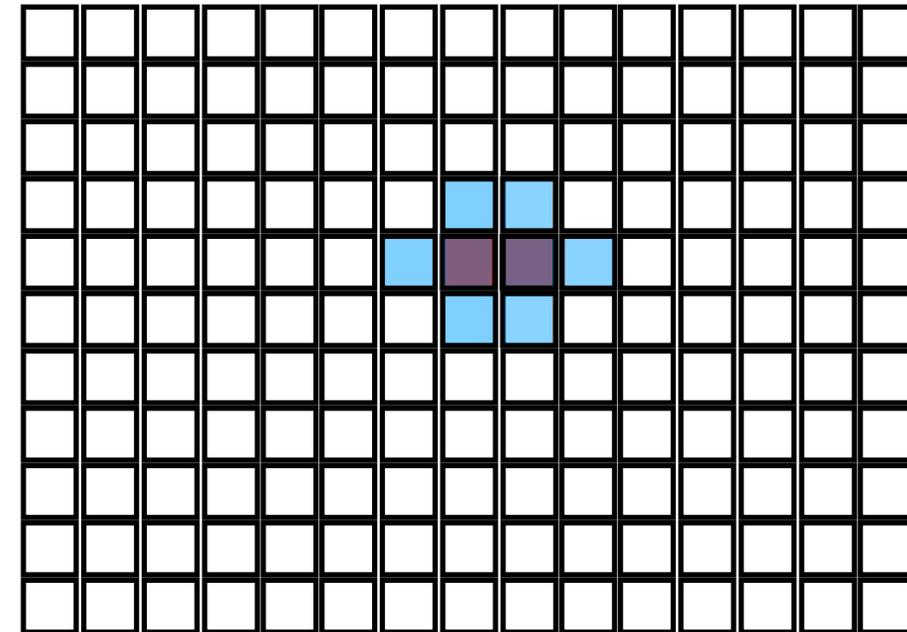


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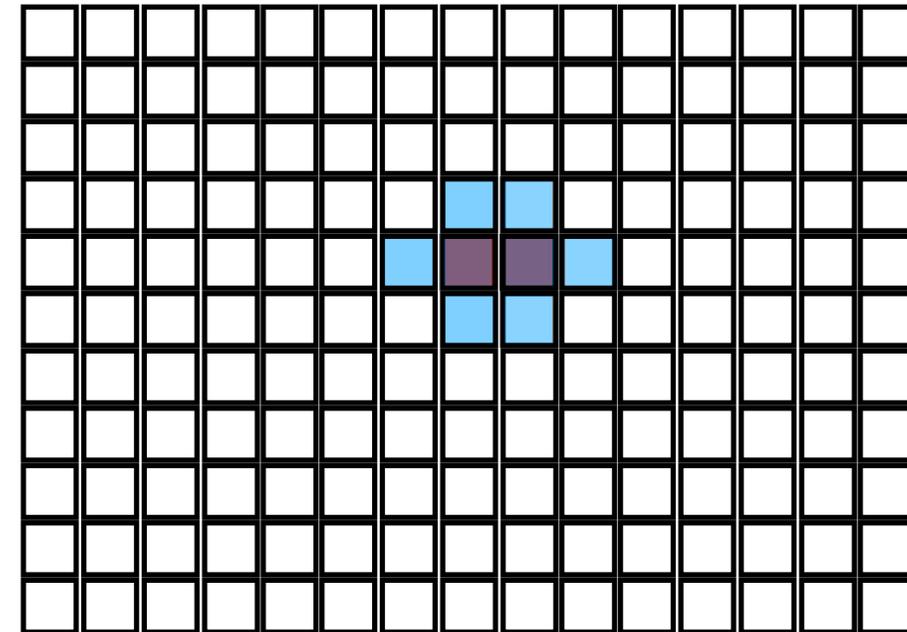


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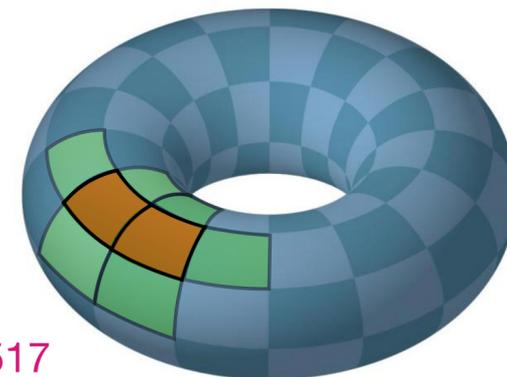
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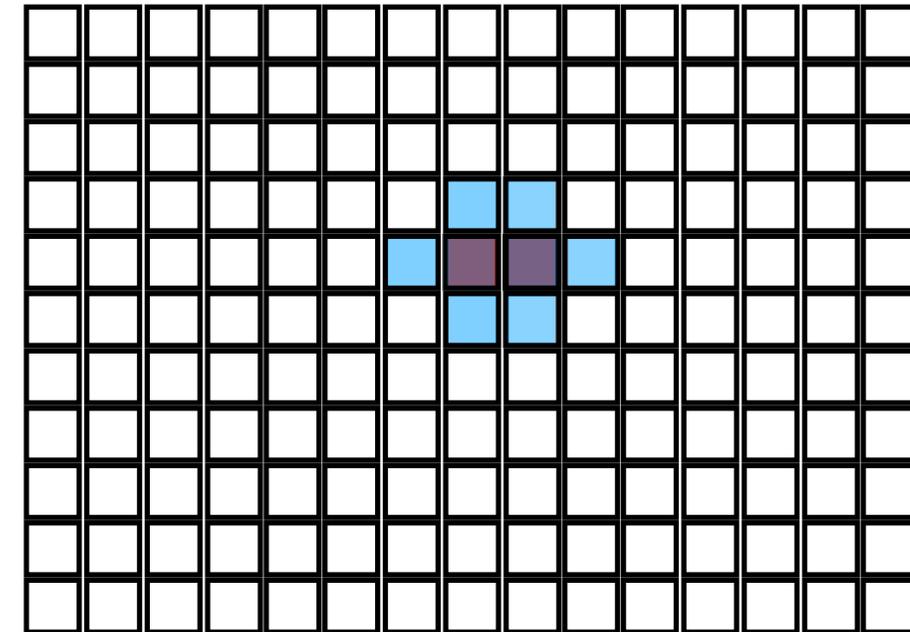


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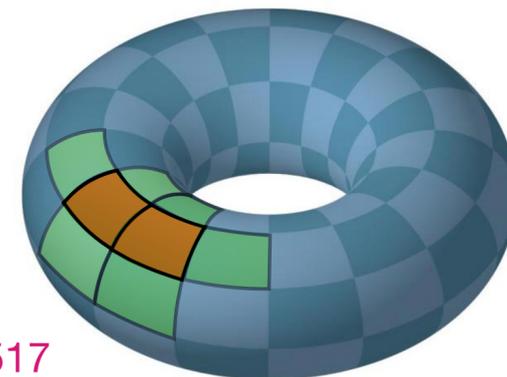
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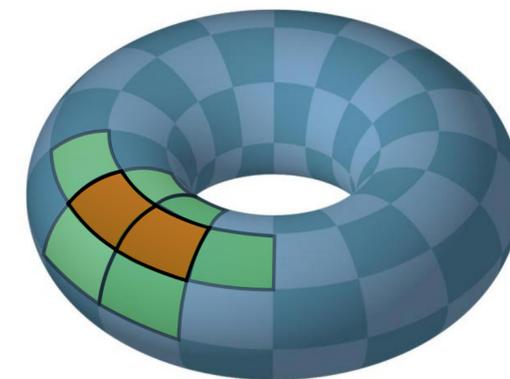
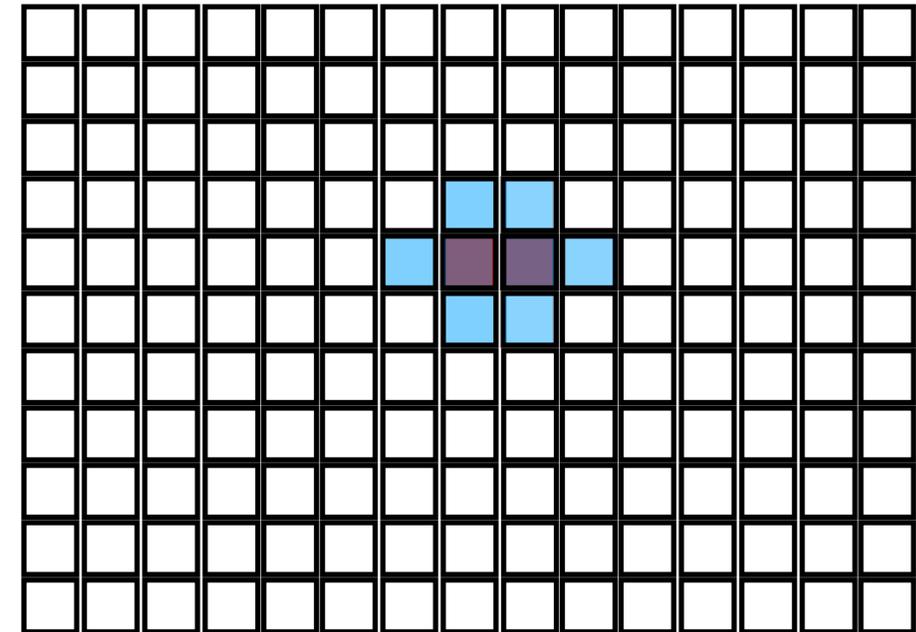
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Wrapping lemma

Unitarity of QCAs

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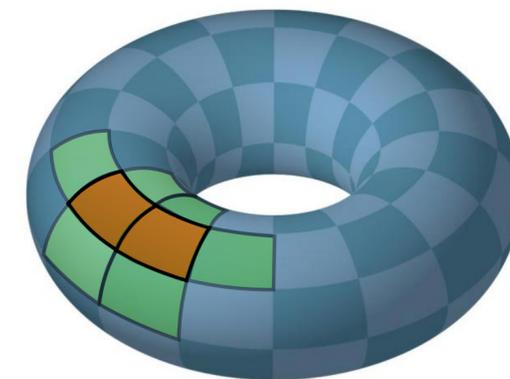
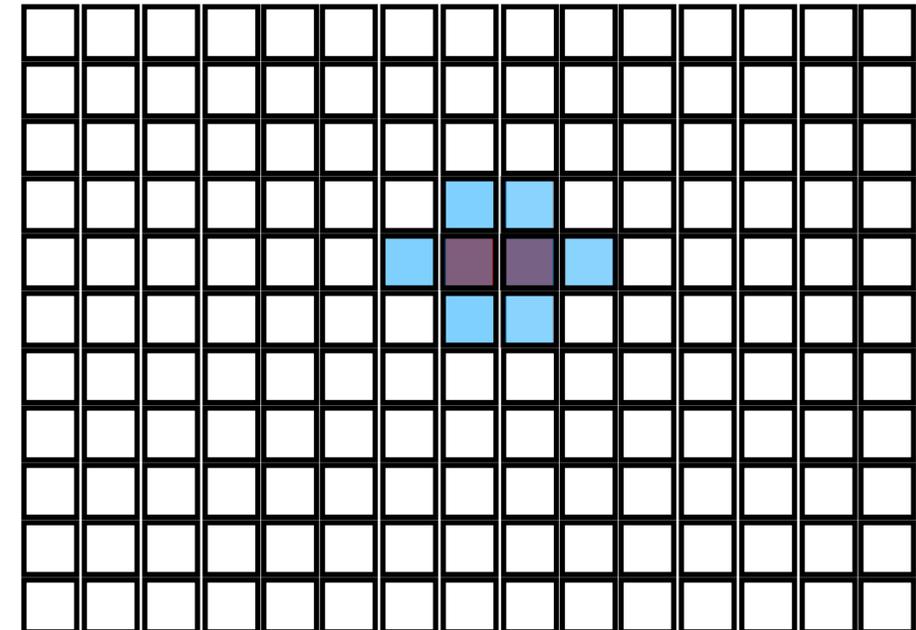


Wrapping lemma

Unitarity of QCAs

- The commutation condition is local
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- By Stinespring theorem in the finite case

$$\begin{aligned}\mathcal{V}(B_R) &= U(I \otimes B_R)U^\dagger \\ &= \sum_i \prod_{x \in R} U_x(I \otimes B_x^{(i)})U_x^\dagger\end{aligned}$$



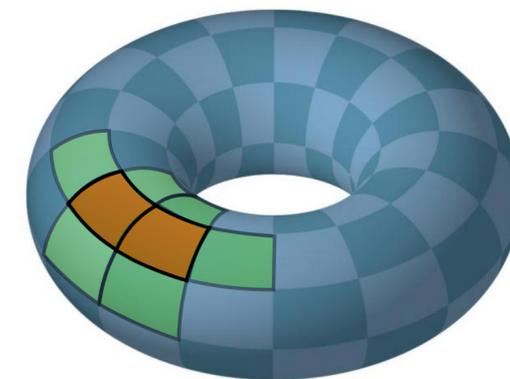
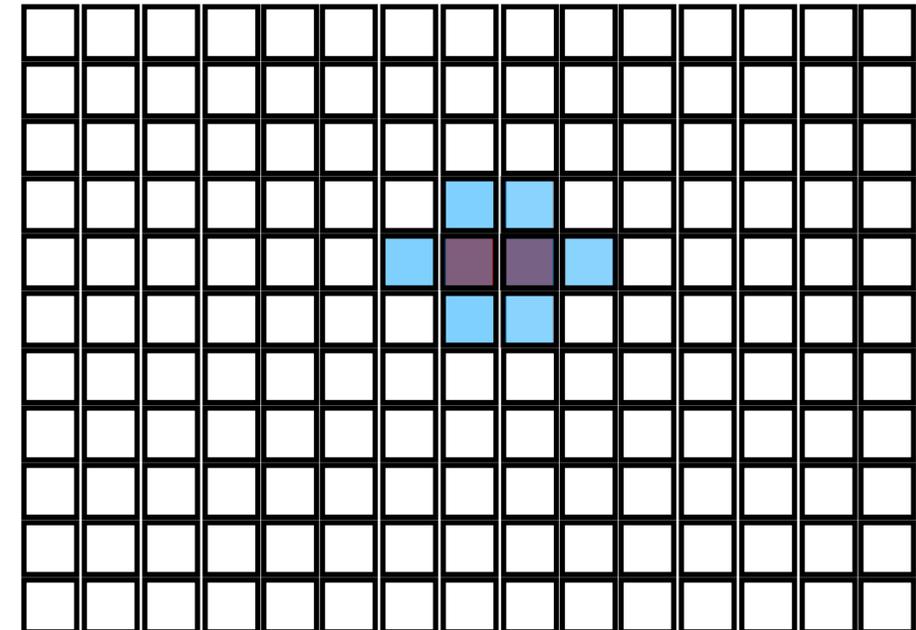
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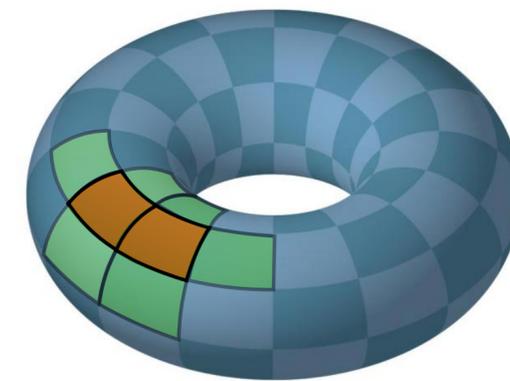
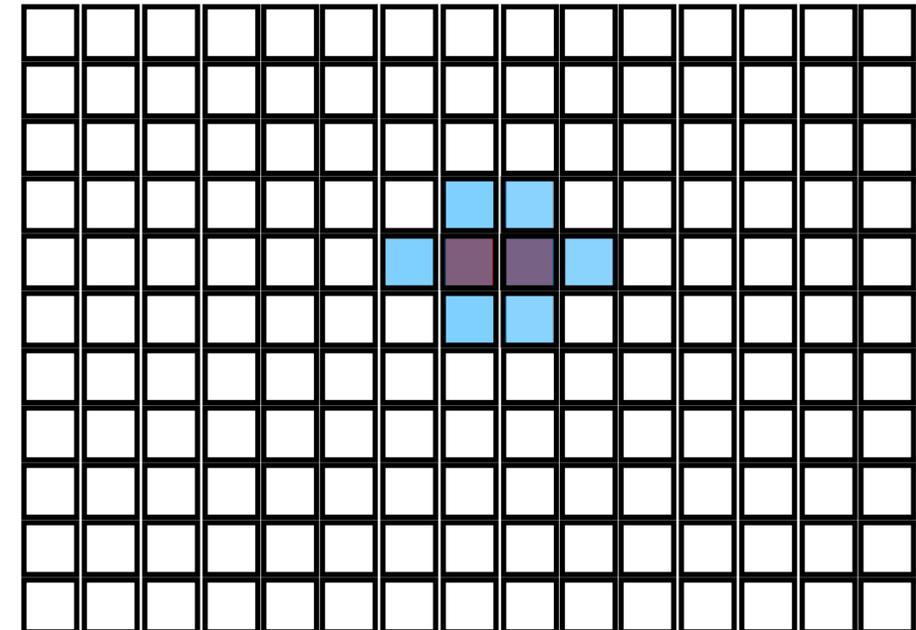
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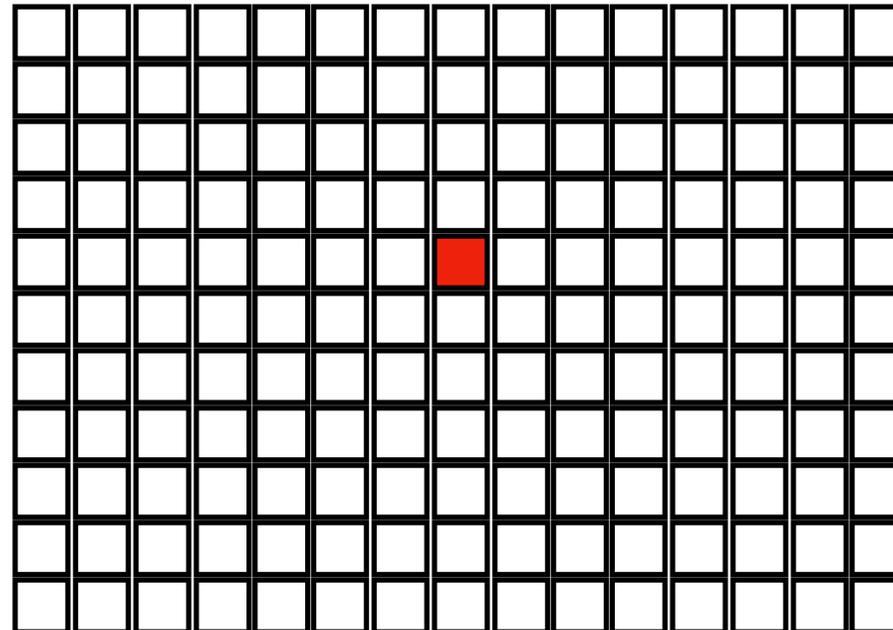
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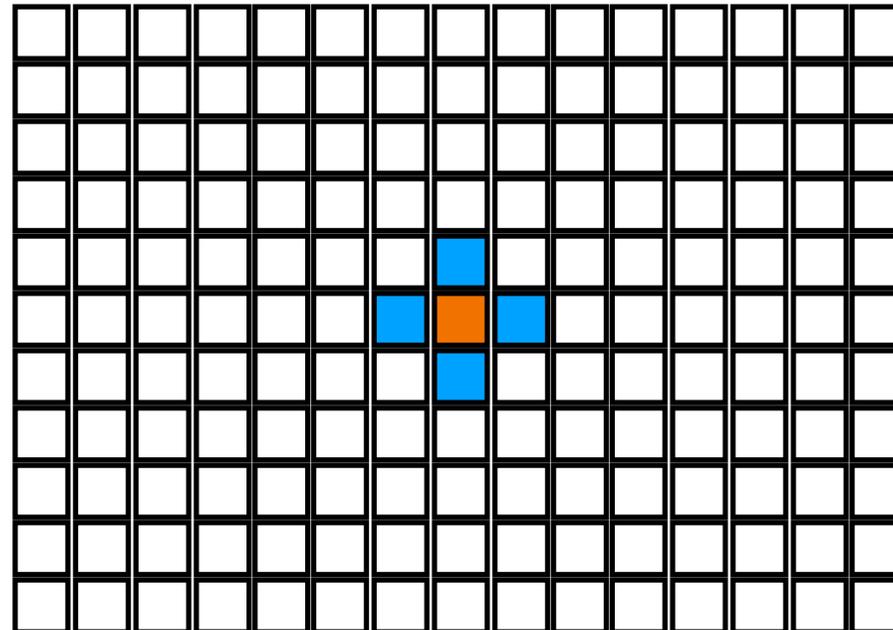
Inverse of a QCA

- The inverse of ν on \mathbb{Z}^d is a QCA itself



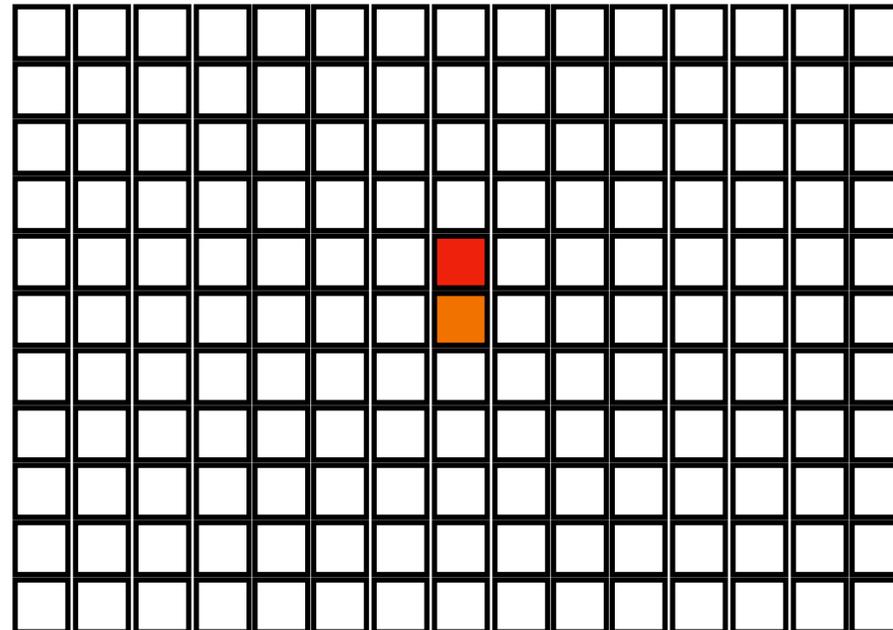
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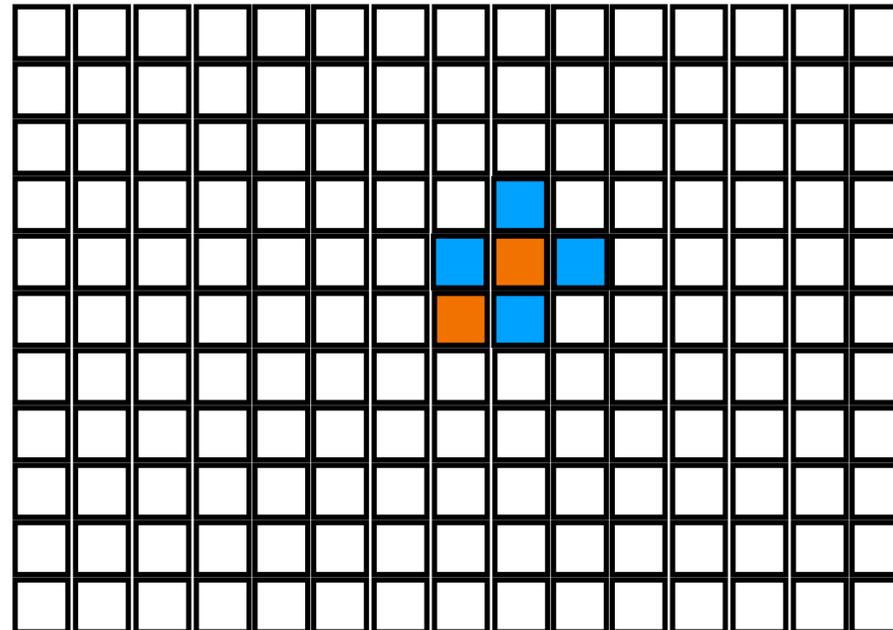
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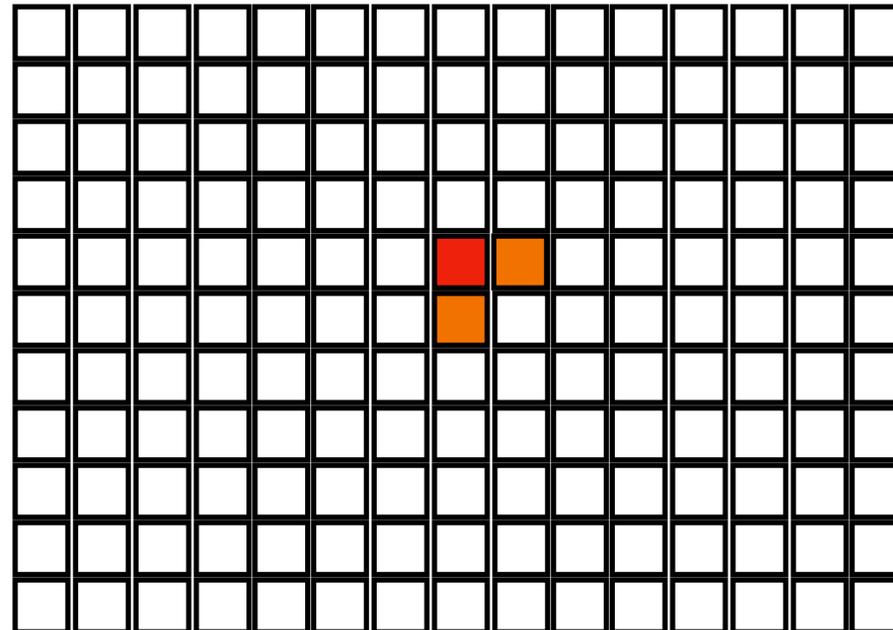
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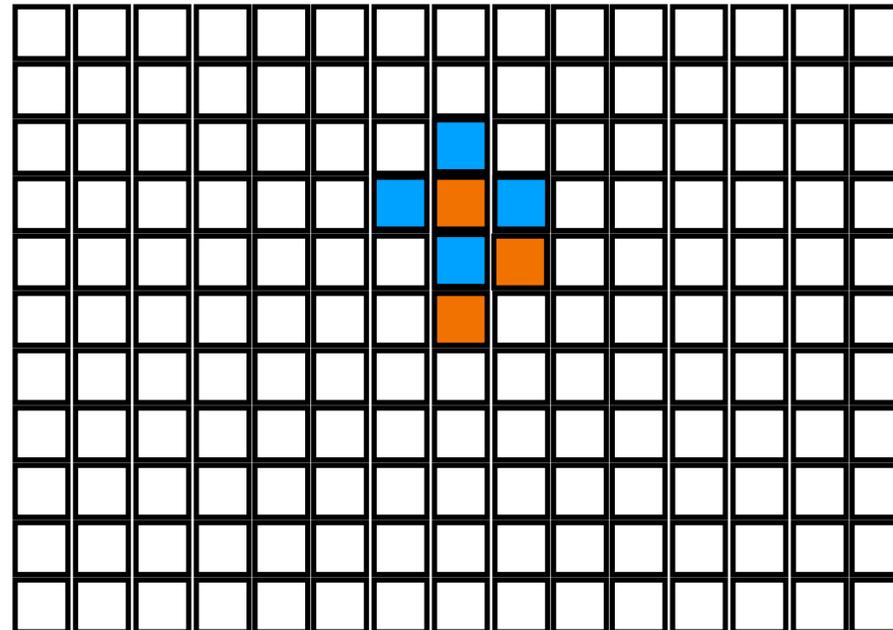
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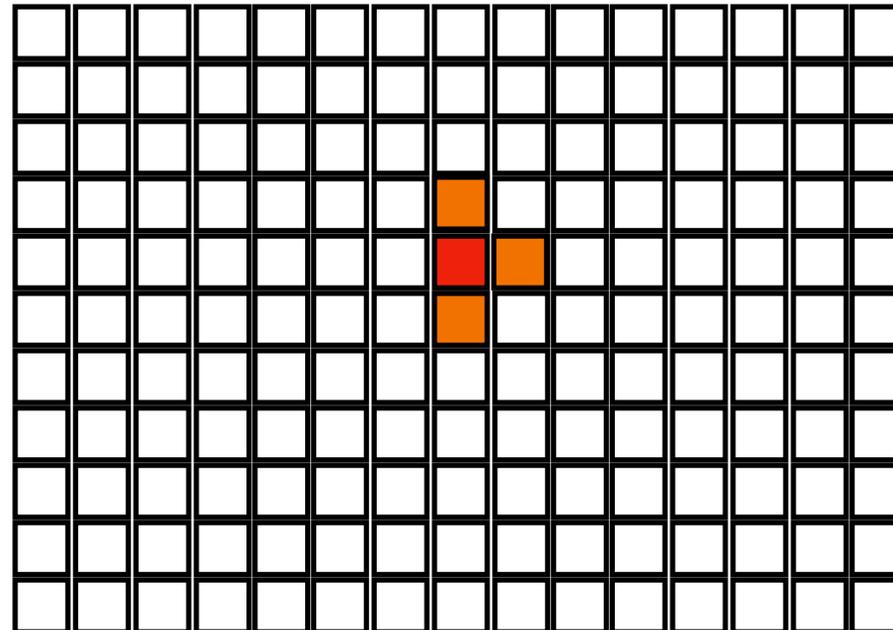
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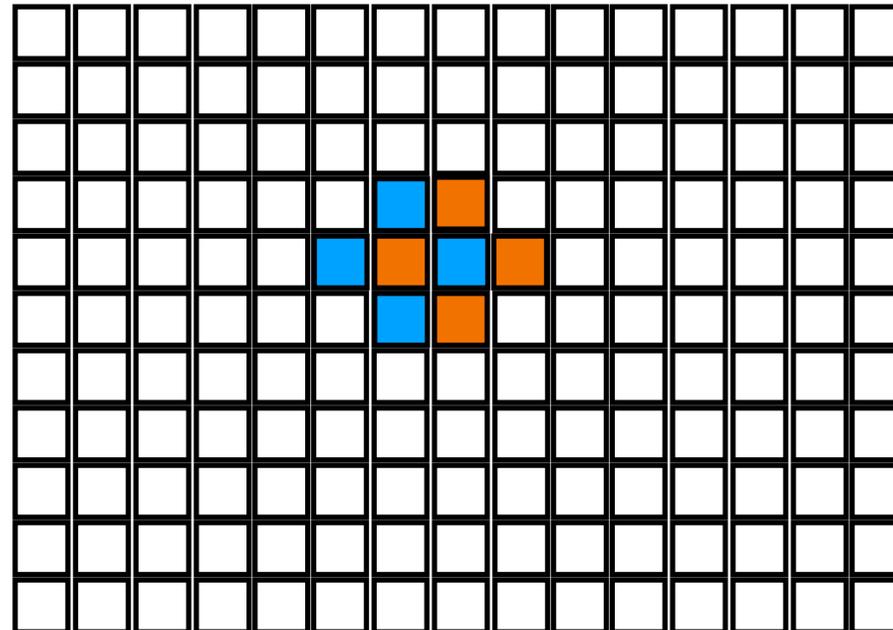
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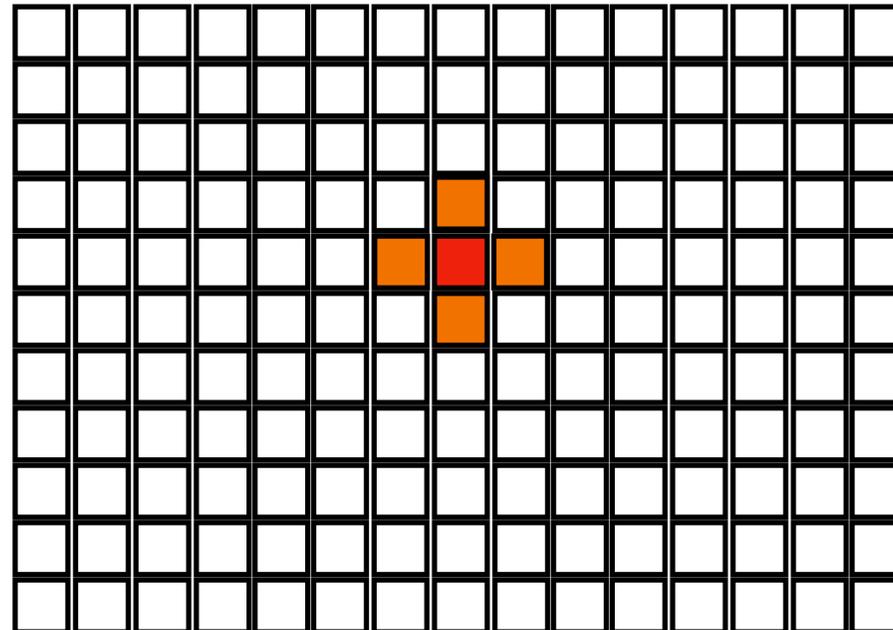
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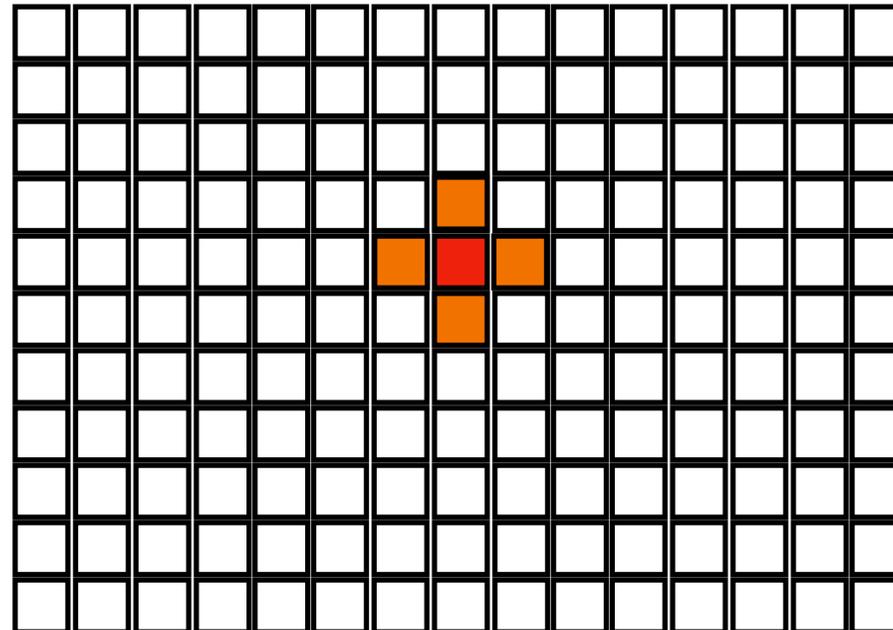
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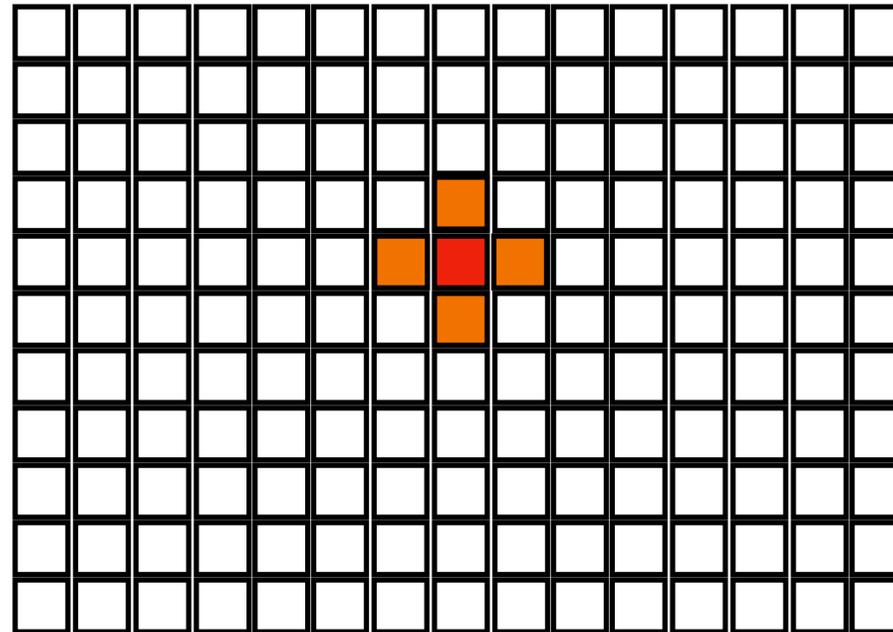
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$$N_{\nu^{-1}}^+(x) = N_{\nu}^-(x)$$

Inverse of a QCA

- The inverse of \mathcal{V} on \mathbb{Z}^d is a QCA itself

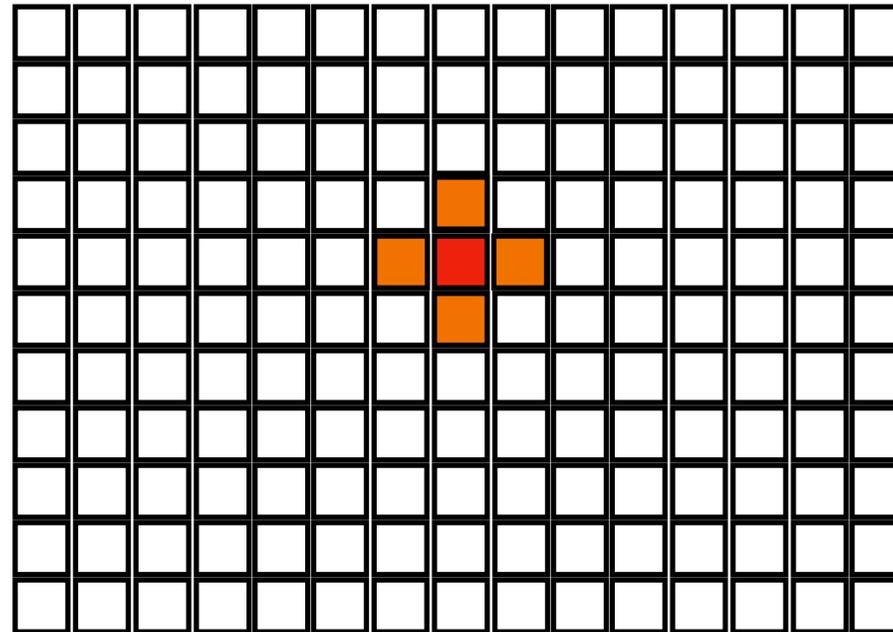


$$N_{\mathcal{V}^{-1}}^+(x) = N_{\mathcal{V}}^-(x)$$

- \mathcal{V}^{-1} satisfies locality

Inverse of a QCA

- The inverse of \mathcal{V} on \mathbb{Z}^d is a QCA itself



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- \mathcal{V}^{-1} satisfies locality
- \mathcal{V}^{-1} is translationally invariant

Problem of quantisation

The neighbourhood of the inverse

- There are classical CA where $N_{V^{-1}}^+(x) \neq N_V^-(x)$

Problem of quantisation

The neighbourhood of the inverse

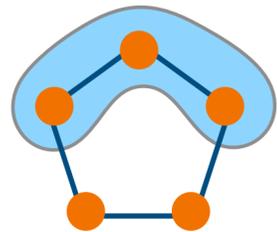
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- Example:
(cells are bits)

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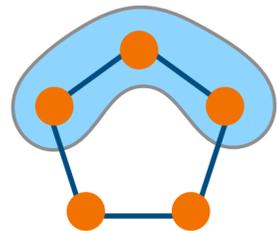
$$(Vb)_{i,t+1} = b_{i\ominus 1,t} \oplus b_{i,t} \oplus b_{i\oplus 1,t}$$

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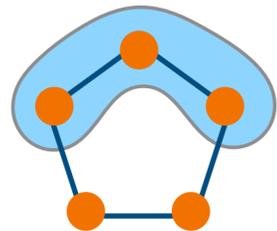
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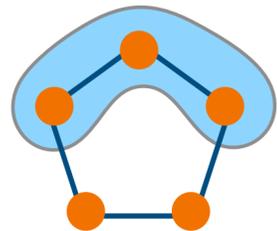
$$\begin{aligned} (V^3b)_{i,t+3} &= (Vb)_{i\ominus 2,t+2} \oplus (Vb)_{i,t+2} \oplus (Vb)_{i\oplus 2,t+2} \\ &= (b_{i\oplus 2,t} \oplus b_{i\ominus 2,t} \oplus b_{i\ominus 1,t}) \oplus (b_{i\ominus 1,t} \oplus b_{i,t} \oplus b_{i\oplus 1,t}) \oplus (b_{i\oplus 1,t} \oplus b_{i\oplus 2,t} \oplus b_{i\ominus 2,t}) \\ &= b_{i,t} \end{aligned}$$

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$$V^3 = I \Leftrightarrow V^{-1} = V^2$$

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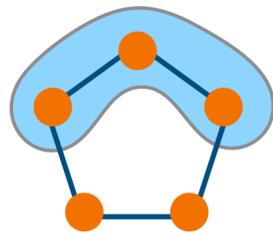
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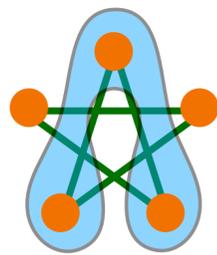
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V

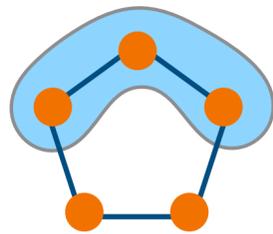


V^{-1}

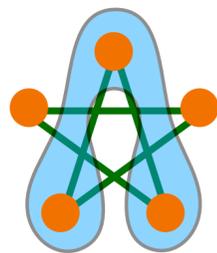
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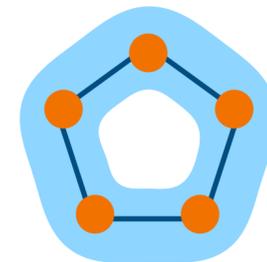
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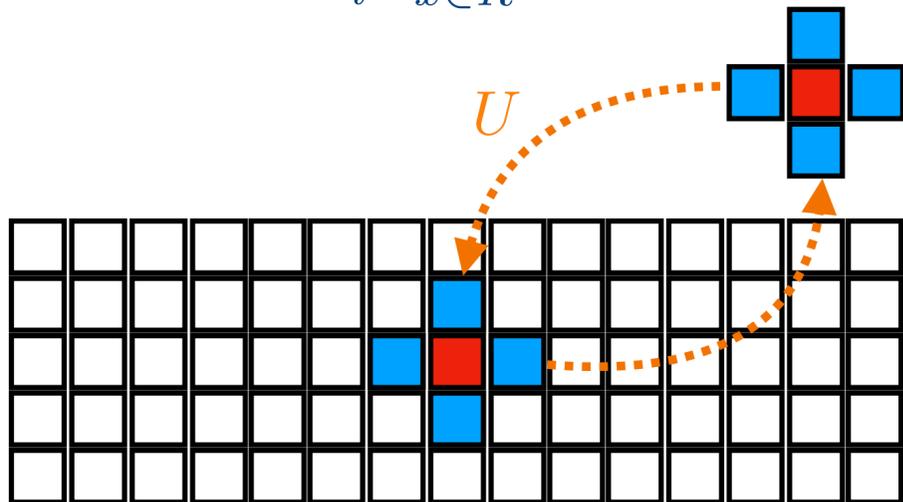


V, V^{-1}

Margolus decomposition

- Let us consider a QCA \mathcal{V} on \mathbb{Z}

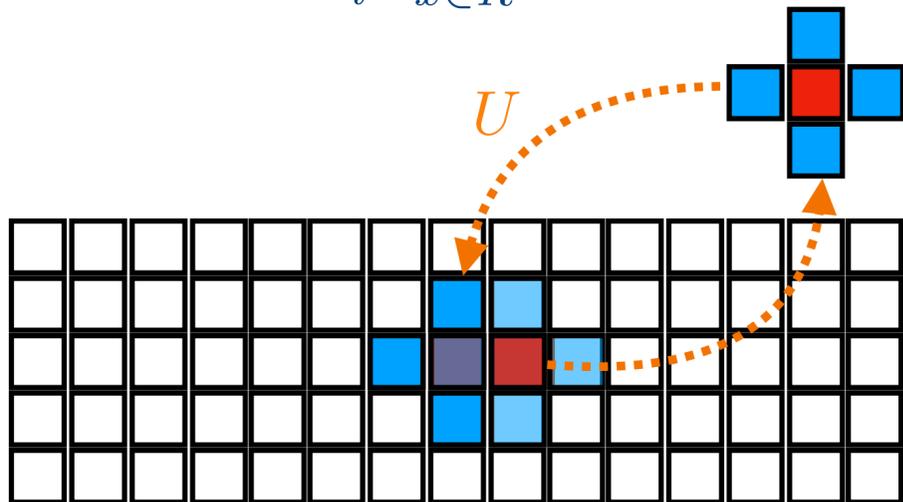
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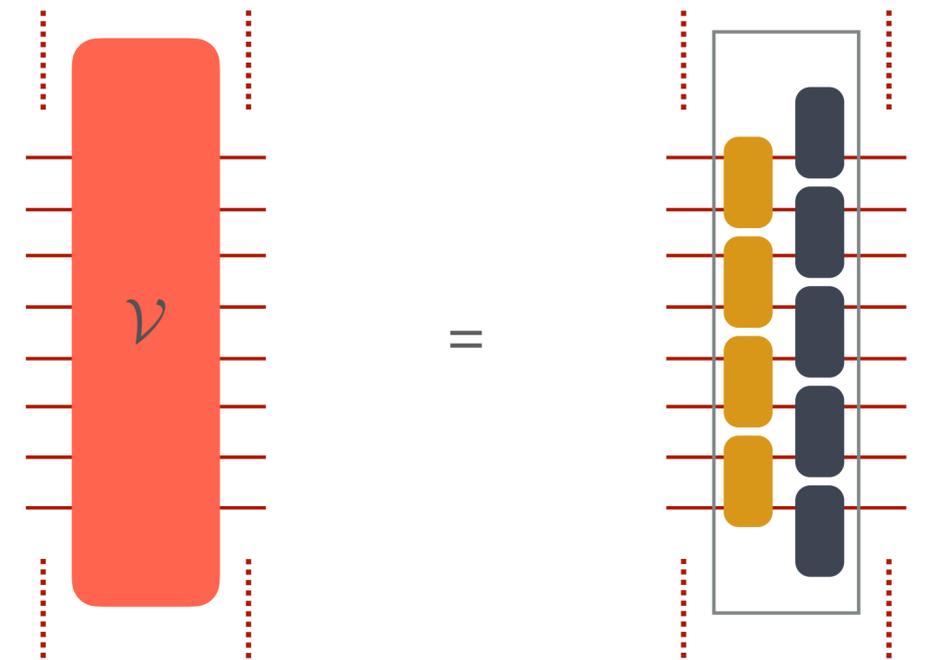
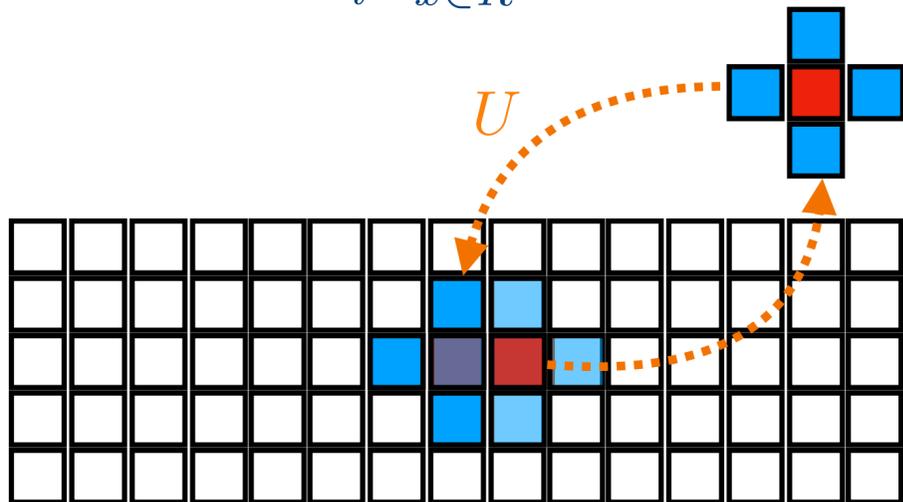
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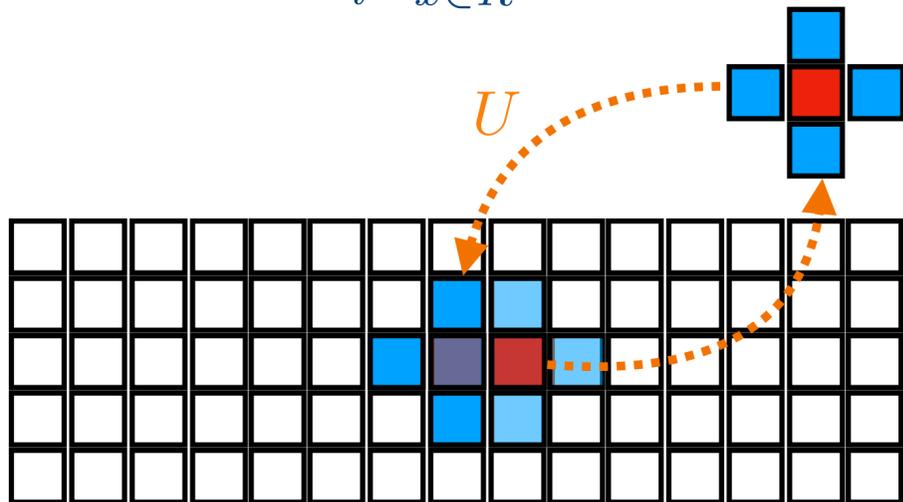
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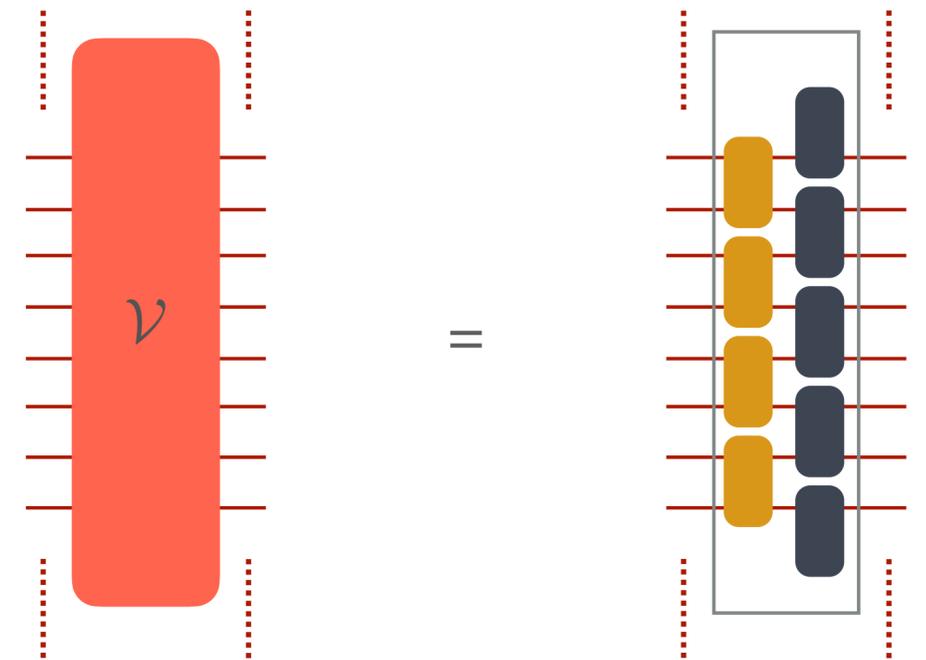
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$$[I \otimes U, U \otimes I] = 0$$



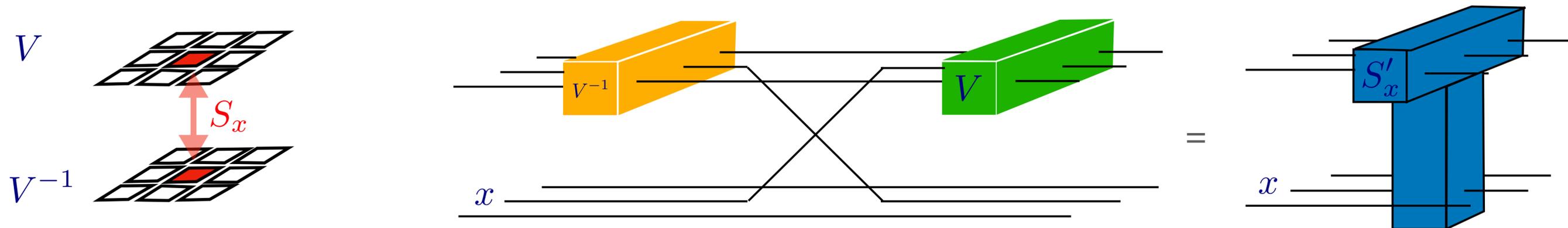
$$\mathcal{V}_\infty(B_R) = \prod_{x \in R} U_x(I \otimes B_R) \prod_{y \in R} U_y^\dagger$$



Unitarity and locality imply Margolus decomposability

- Let \mathcal{V} be a QCA on \mathbb{Z}^d with cell $H \simeq \mathbb{C}^r$
- Consider the QCA $\mathcal{V} \otimes \mathcal{V}^{-1}$ on \mathbb{Z}^d with cell $H \simeq \mathbb{C}^r \otimes \mathbb{C}^r$
- Since $V \otimes V^{-1} = (V \otimes I)S(V^{-1} \otimes I)S$, where $S|\phi\rangle|\psi\rangle := |\psi\rangle|\phi\rangle$, and $S = \bigotimes_x S_x$

$$V \otimes V^{-1} = \prod_x (S'_x) \bigotimes_y S_y \quad S'_x := (V \otimes I)S_x(V^{-1} \otimes I)$$



Fermionic theory

- The theory is meant to provide a realisation of the fermion algebra

$$\{\varphi_i^\dagger, \varphi_j\} = \delta_{ij}I, \quad \{\varphi_i, \varphi_j\} = 0$$

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- Example of basis state: $|0010110\rangle = \varphi_3^\dagger \varphi_5^\dagger \varphi_6^\dagger |0000000\rangle$

Theories without local discriminability

Example 2: Fermionic quantum theory

- The representation depends on the chosen ordering of the LFMs

$$J(\varphi_i) := I_1 \otimes \cdots \otimes I_{i-1} \otimes \sigma_i^- \otimes \sigma^z_{i+1} \cdots \otimes \sigma^z_N$$

$$J(XY) := J(X)J(Y) \quad J(X^\dagger) := J(X)^\dagger$$

$$J(aX + bY) := aJ(X) + bJ(Y)$$

*Bravyi and Kitaev, Annals of Physics **298**, 210–226 (2002)

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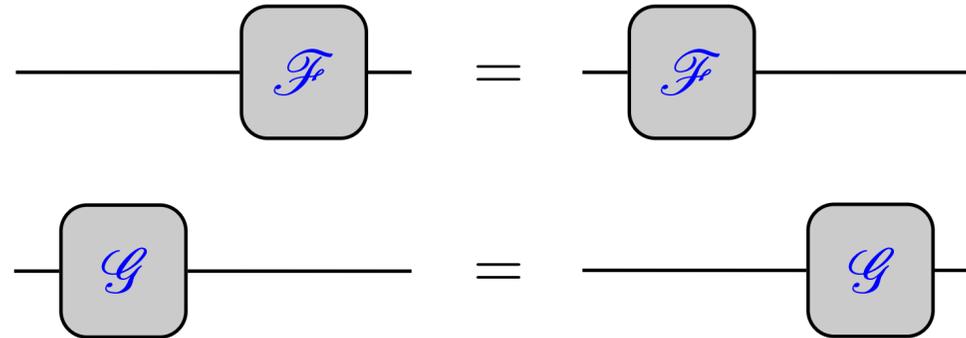
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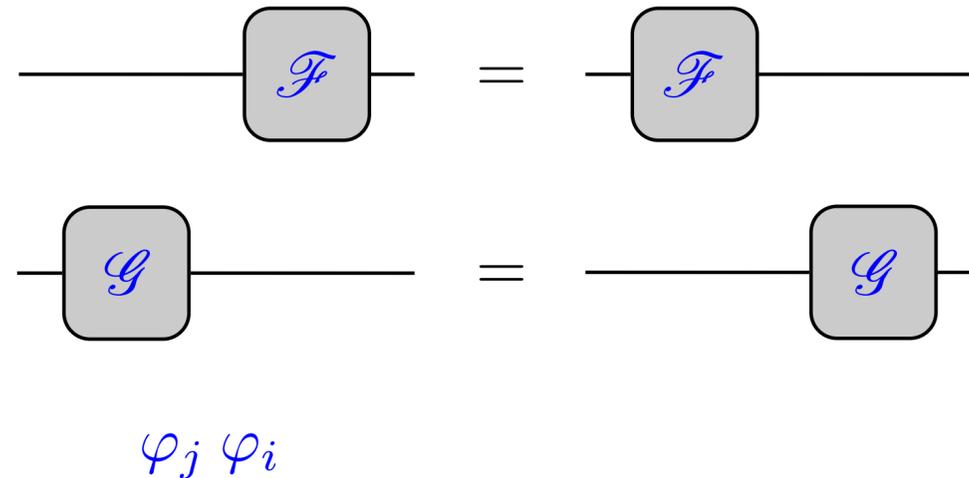
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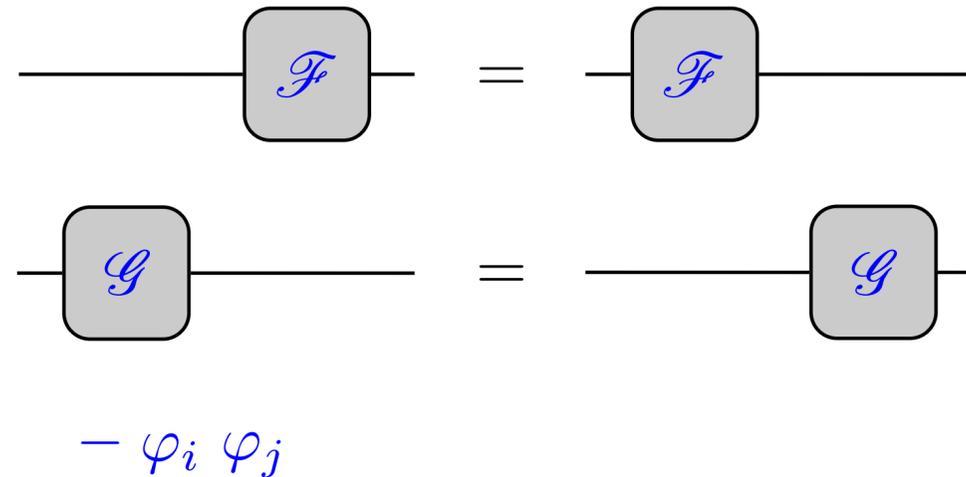
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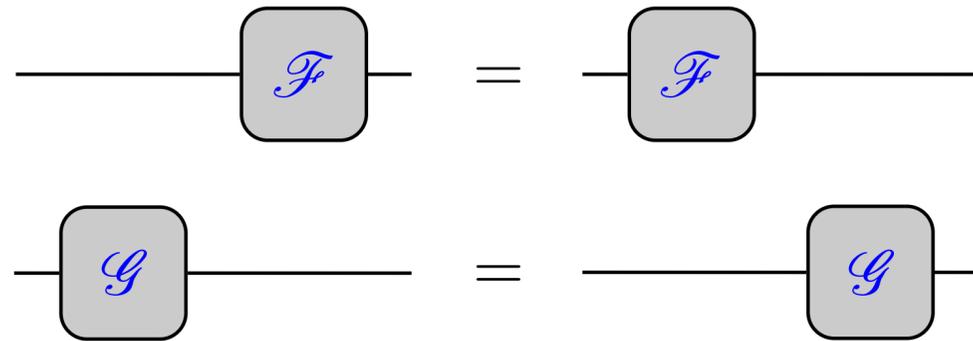
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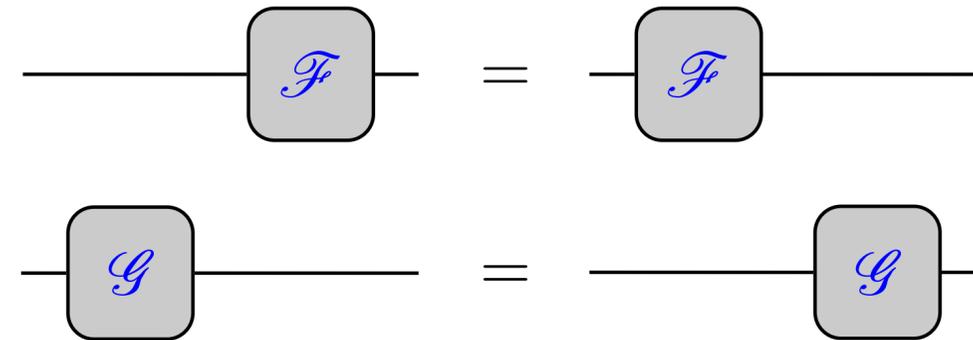


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$$\varphi_j (\varphi_i^\dagger \varphi_i + \varphi_i)$$

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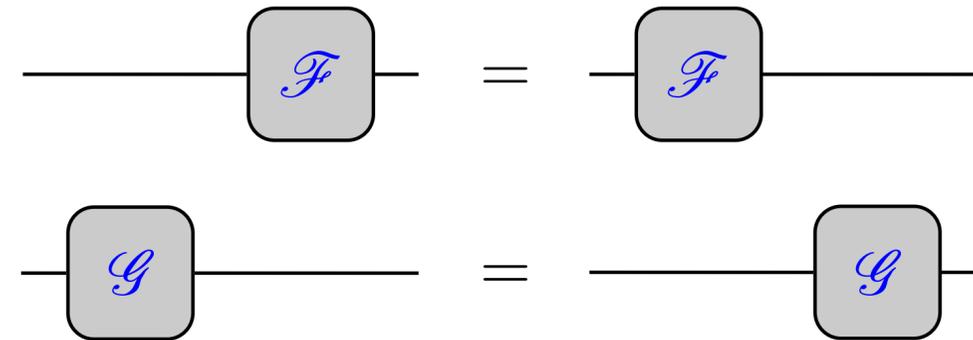


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 \end{aligned}$$

Theories without local discriminability

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$$\begin{array}{c}
 \text{---} \boxed{F} \text{---} = \text{---} \boxed{F} \text{---} \\
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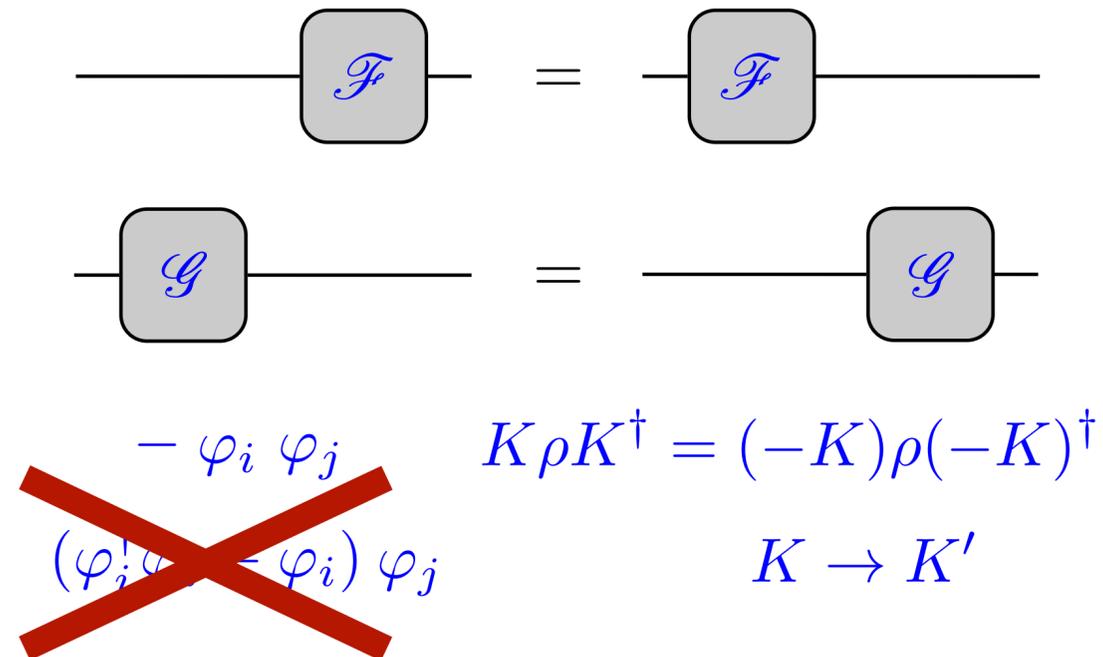
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- States and effects are combinations **even products** of field operators

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- The states of one LFM are the states of a classical bit $|0\rangle\langle 0|, |1\rangle\langle 1|$
- Parity superselection \rightarrow block-diagonal structure for states

$$J(\rho) = \left(\begin{array}{c|c} p\rho_O & 0 \\ \hline 0 & (1-p)\rho_E \end{array} \right)$$

Fermionic CA

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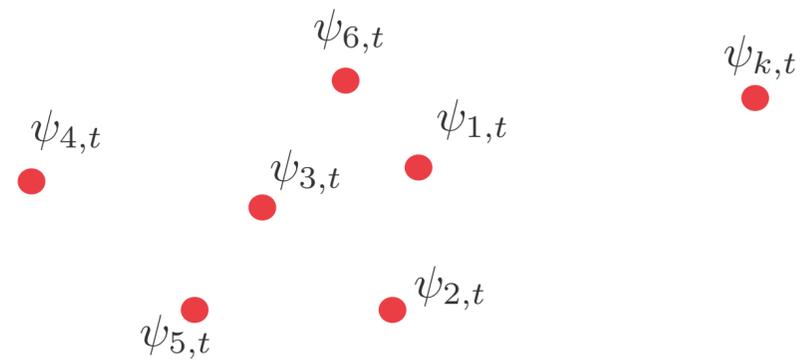
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 - Wrapping lemma: true also in the infinite case

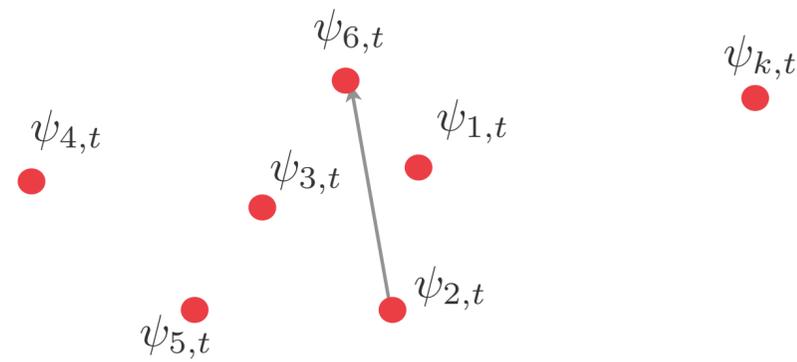
Linear FCAs and quantum walks



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$$A_{ij} \in M_{s_i \times s_j}$$

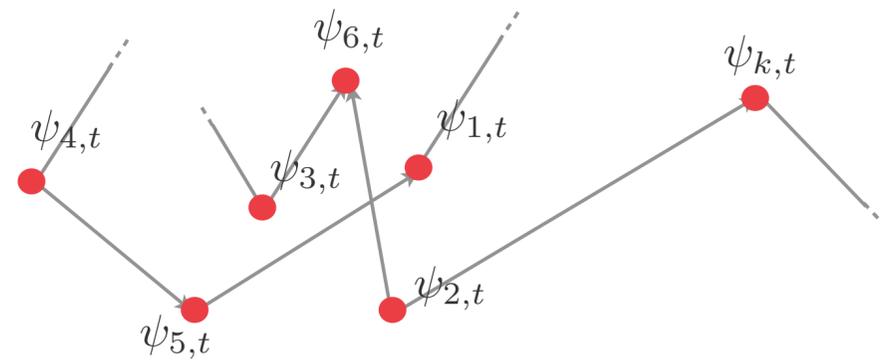
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