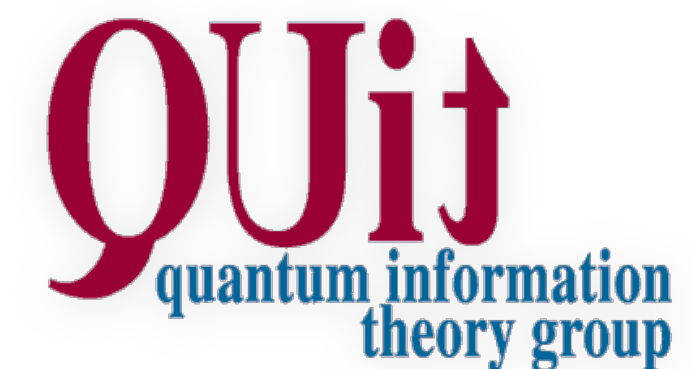


# Operational probabilistic theories and cellular automata: how I learned to stop worrying and love $C^*$ algebras

School on Advanced Topics in Quantum Information and Foundations

Quantum Information Unit and the Yukawa Institute for Theoretical Physics, Kyoto University



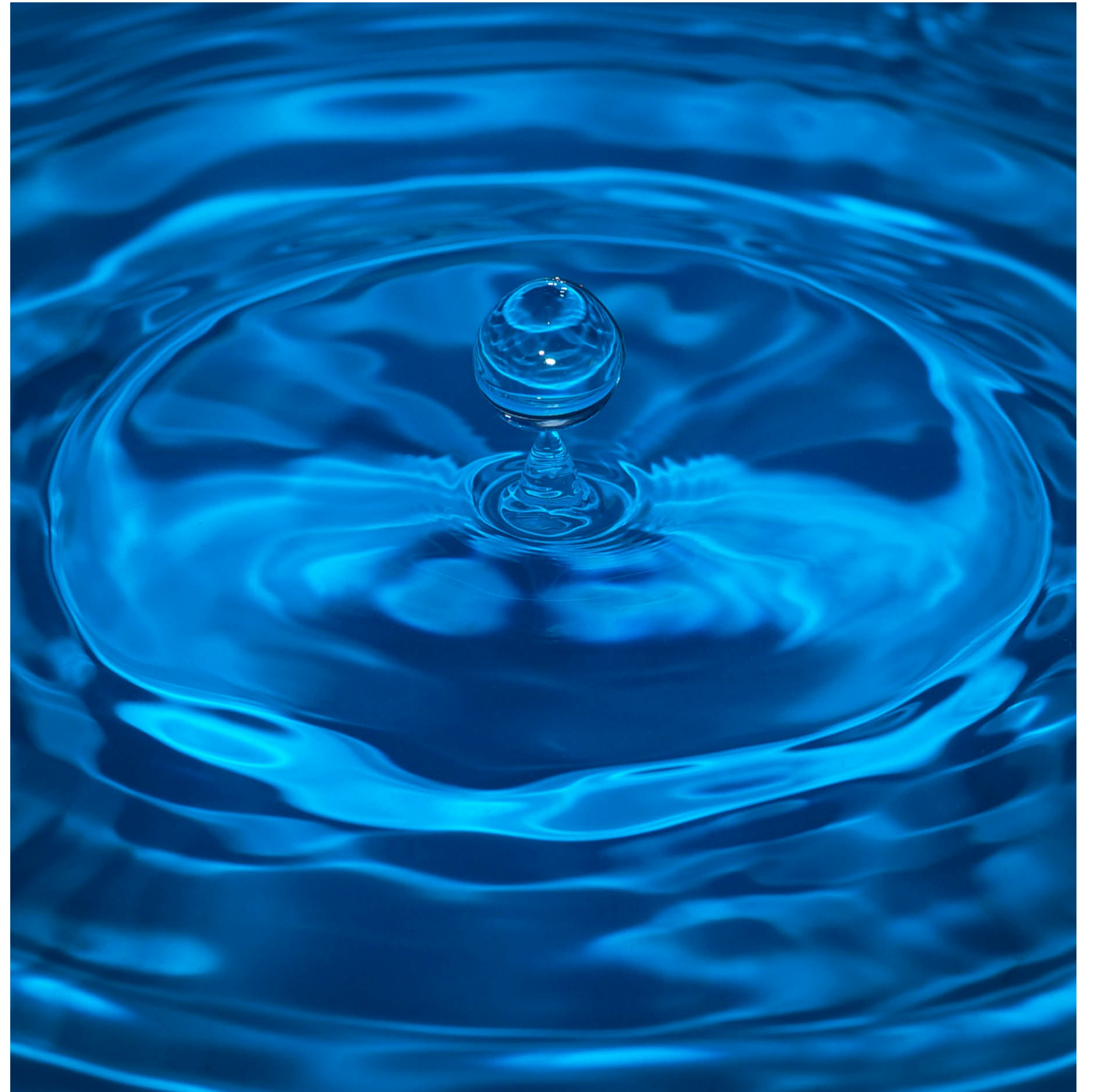
Paolo Perinotti - February 8-12 2021

# Lecture 4

## Causal influence in OPTs

# Summary

- Networks and causal cones
- Signalling
- Propagation of interventions
- The comb structure
- Classical and Quantum theory
- No interaction without disturbance



# The neighbourhood

## Causal influence

- Let us consider the local algebra  
 $[[A_g C \rightarrow A_g C]]_{Q\mathbb{R}}$

# The neighbourhood

## Causal influence

- Let us consider the local algebra  
 $[[A_g C \rightarrow A_g C]]_{Q\mathbb{R}}$
- It is transformed in a subalgebra of  
some region  $R$

$$\mathcal{V}[[A_g C \rightarrow A_g C]]_{Q\mathbb{R}}\mathcal{V}^{-1} \subseteq [[A_R C \rightarrow A_R C]]_{Q\mathbb{R}}$$

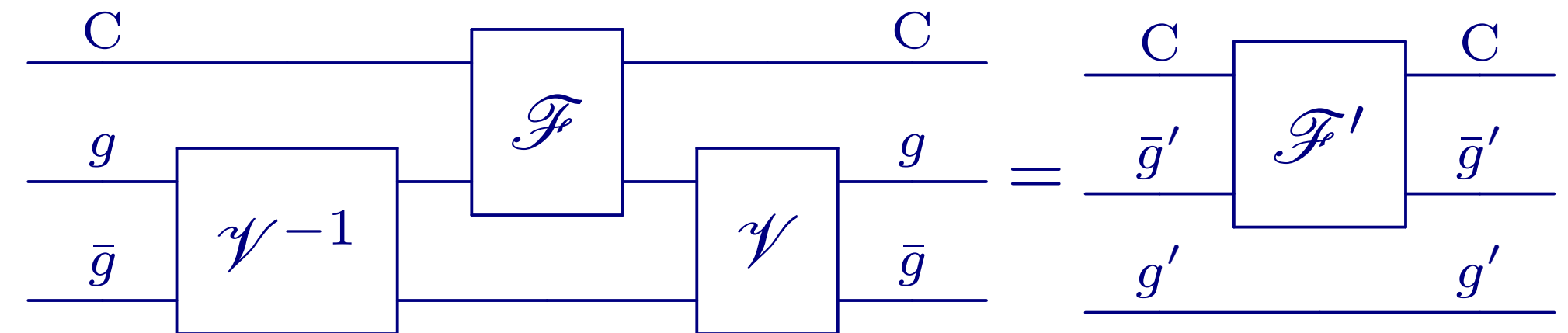
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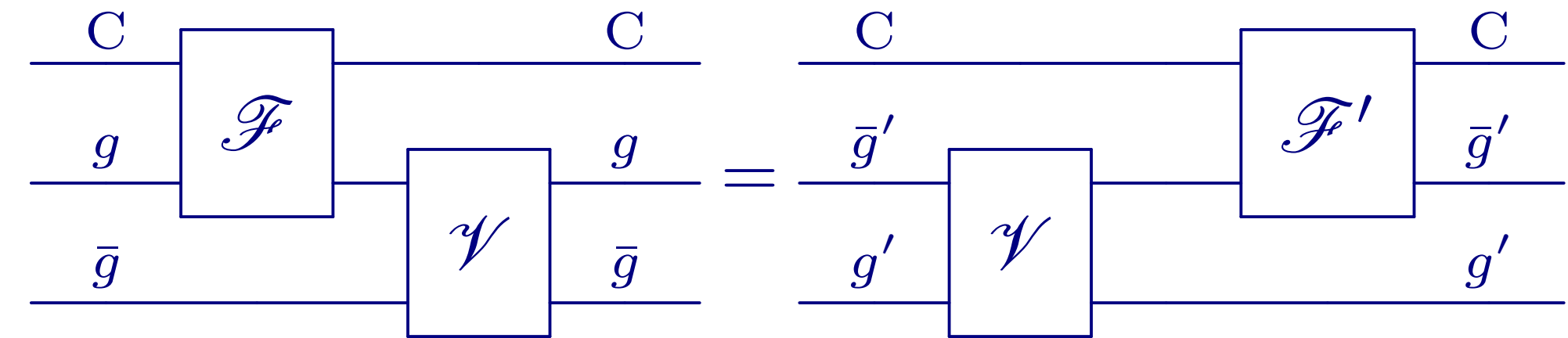
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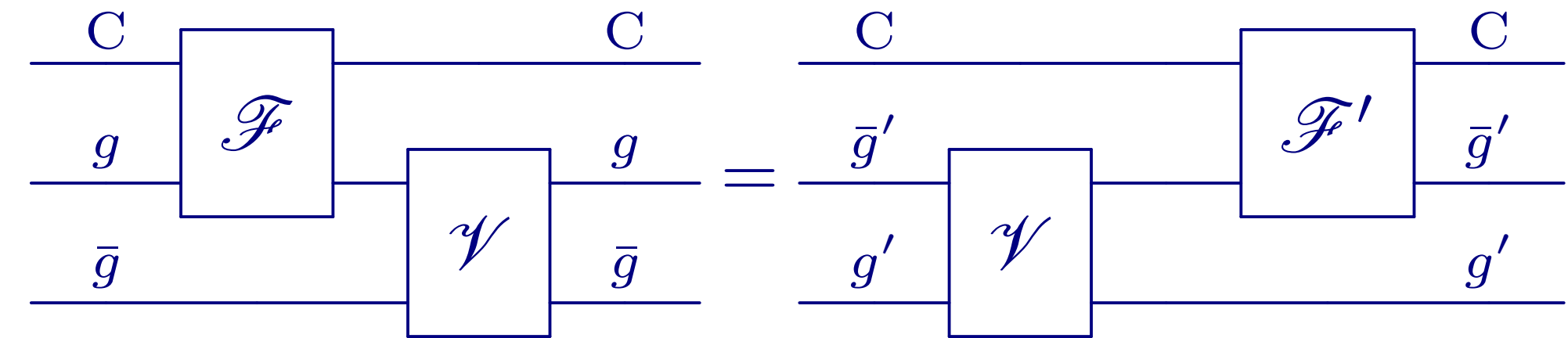
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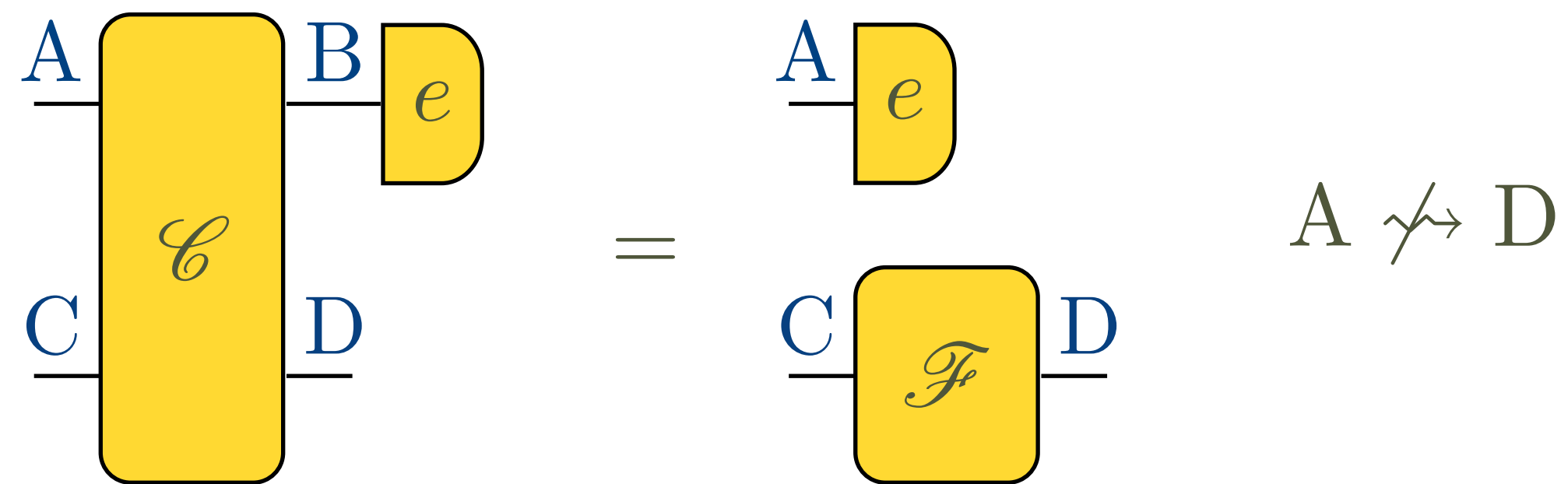
- We say that  $g$  does not causally influence  $g'$  if for any  $C$ ,  $g' \notin R$
- The forward neighbourhood  $N_g^+$  of  $g$  is the set of all  $g'$  such that  $g$  causally influences  $g'$





# Causal influence in quantum literature

The traditional approach: start from no-signalling



No intervention on the state of  $A$  can influence the state of  $C$

In quantum theory

$$\text{Tr}_B[R_{\mathcal{C}}] = I_A \otimes R_{\mathcal{F}}$$

$R_{\mathcal{C}}$  denoting the Choi operator corresponding to  $\mathcal{C}$

# Defining (no) causal influence in OPTs

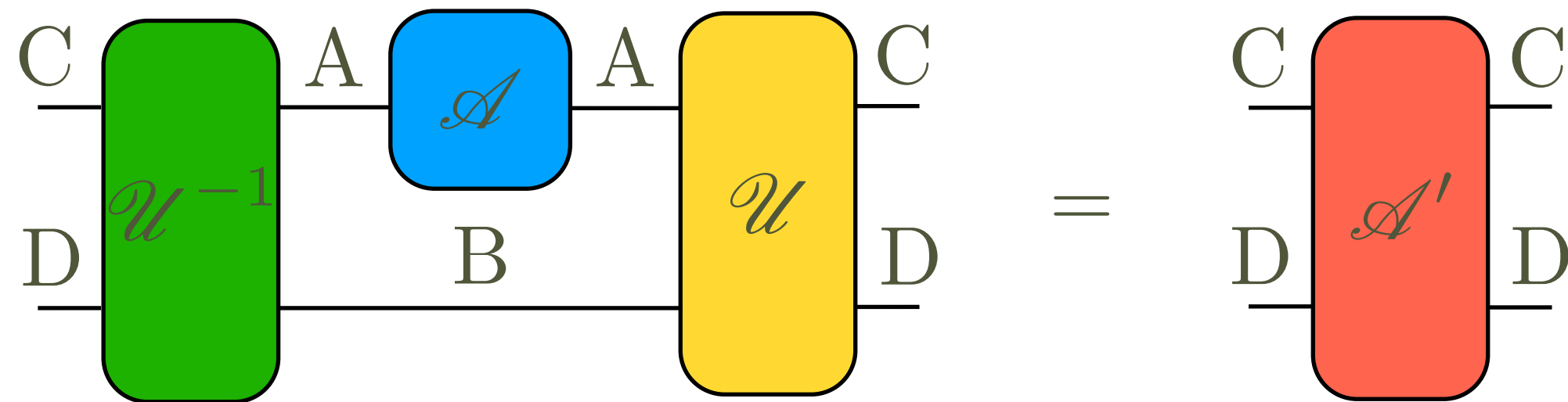
## The raw idea

- The definition is inspired by the notion of neighbourhood in QCAs
- It holds for **reversible** transformations

# Defining (no) causal influence in OPTs

## The raw idea

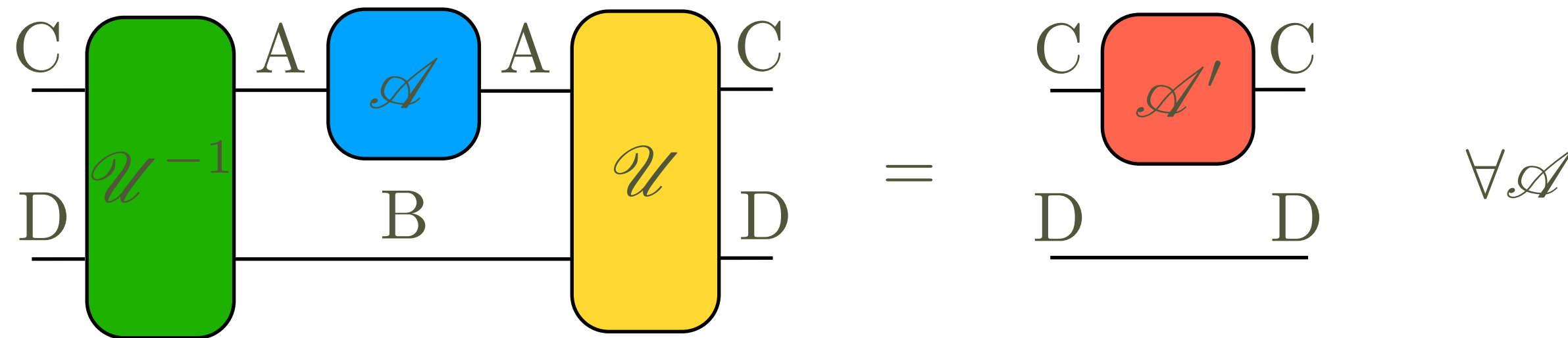
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# Defining (no) causal influence in OPTs

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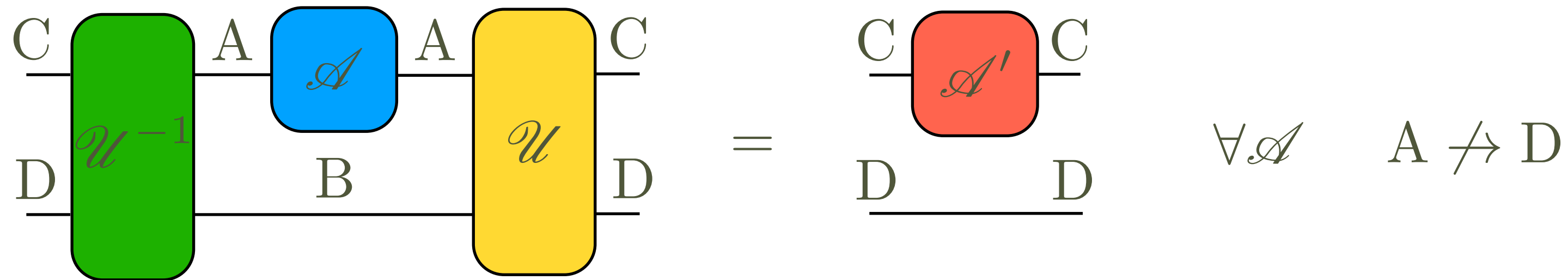
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# Defining (no) causal influence in OPTs

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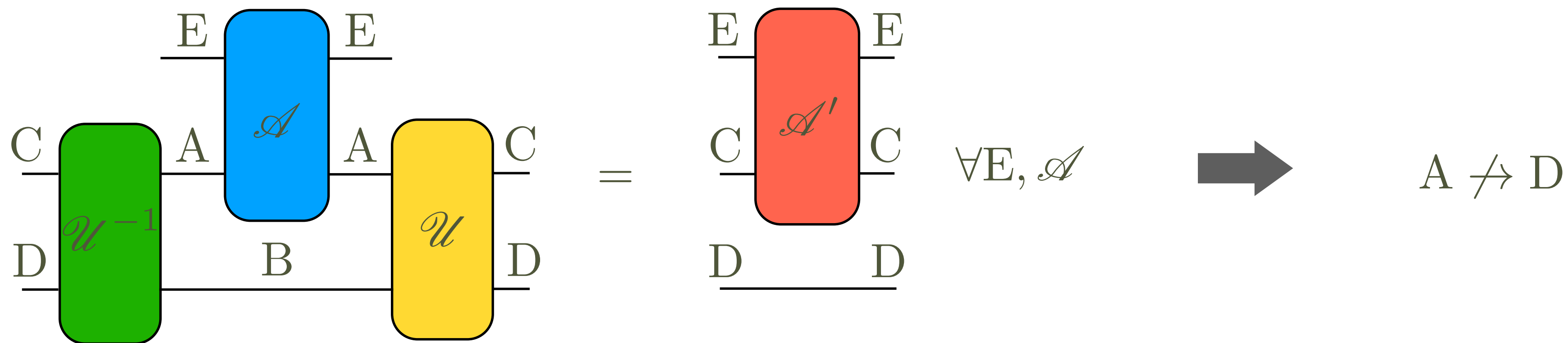
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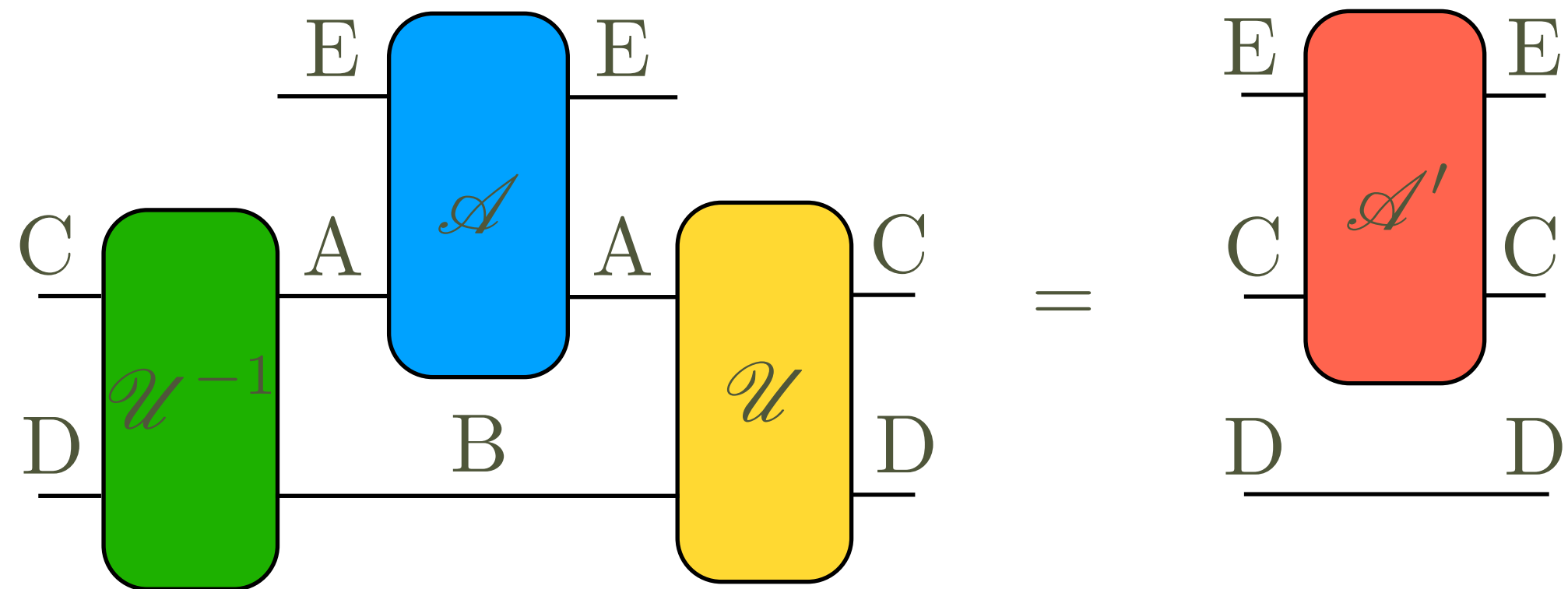
## The precise notion

- Without local discriminability (local tomography/tomographic locality) we need to take into account interventions involving ancillary systems



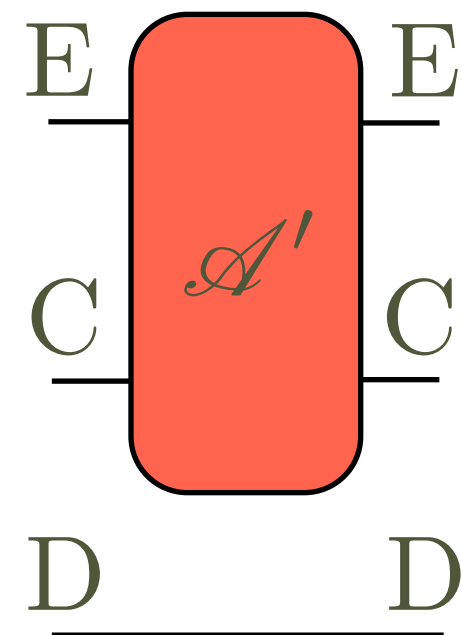
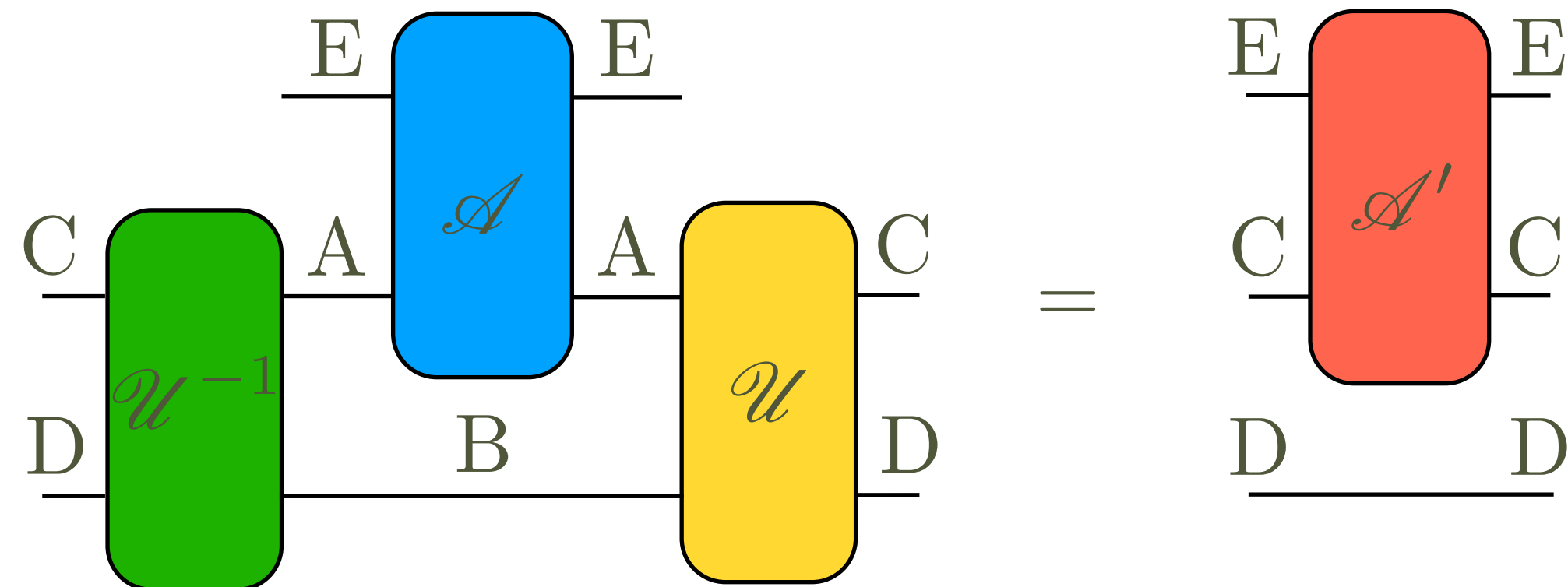
# Explanation of the definition

Suppose that under  $\mathcal{U}$  one has  $A \not\rightarrow D$

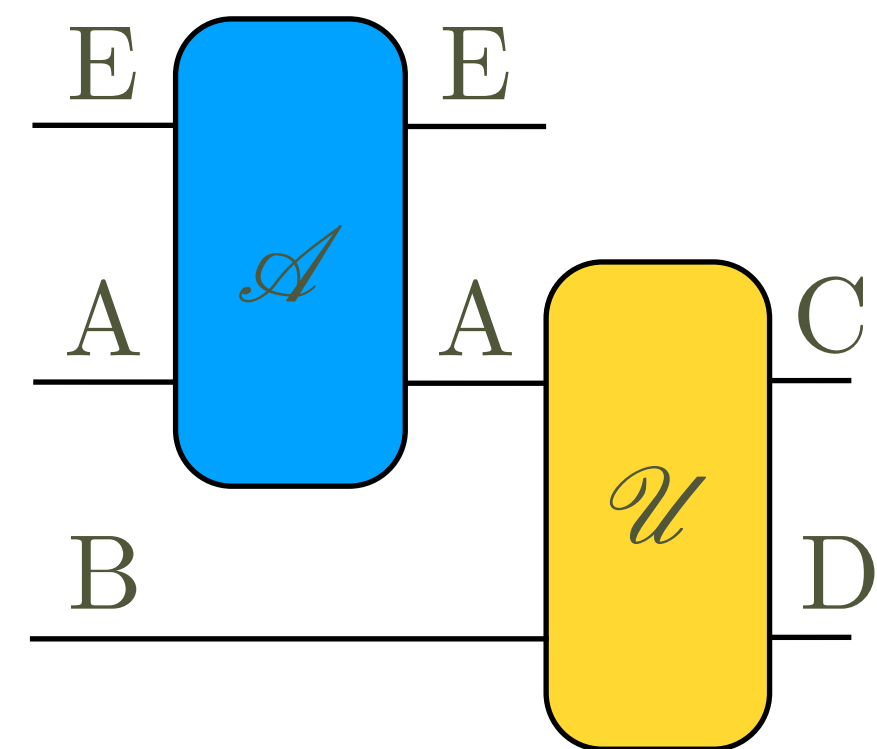


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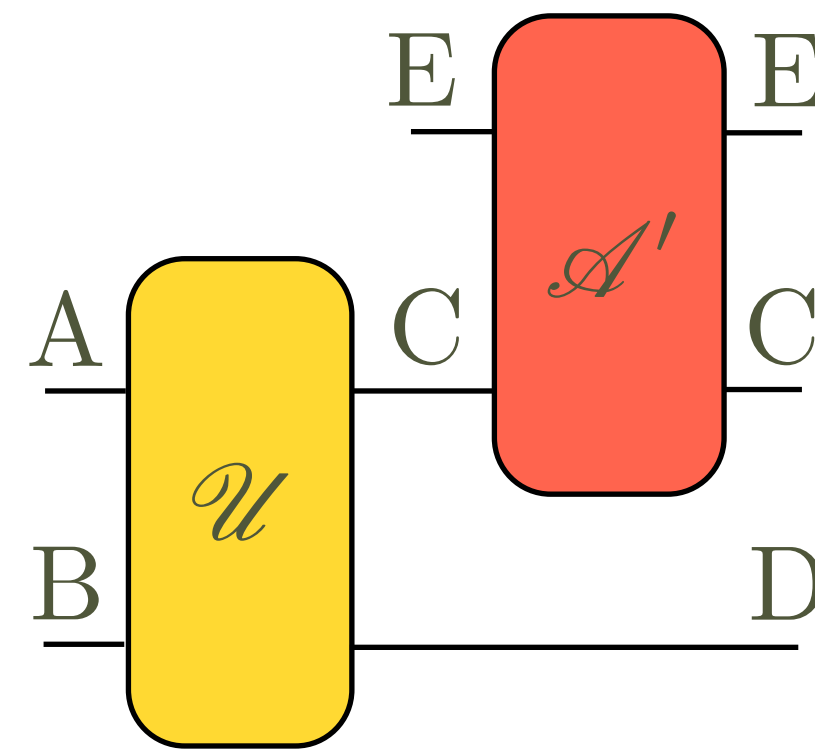
Suppose that under  $\mathcal{U}$  one has  $A \not\rightarrow D$



Equivalently:



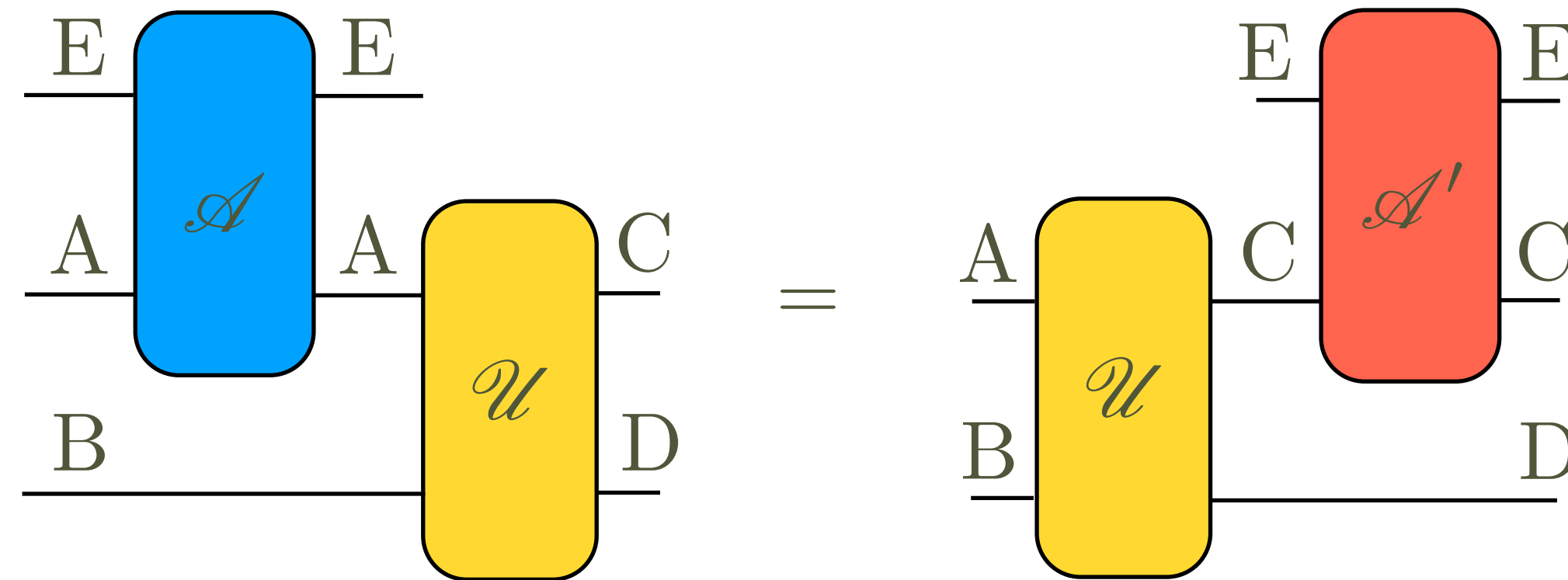
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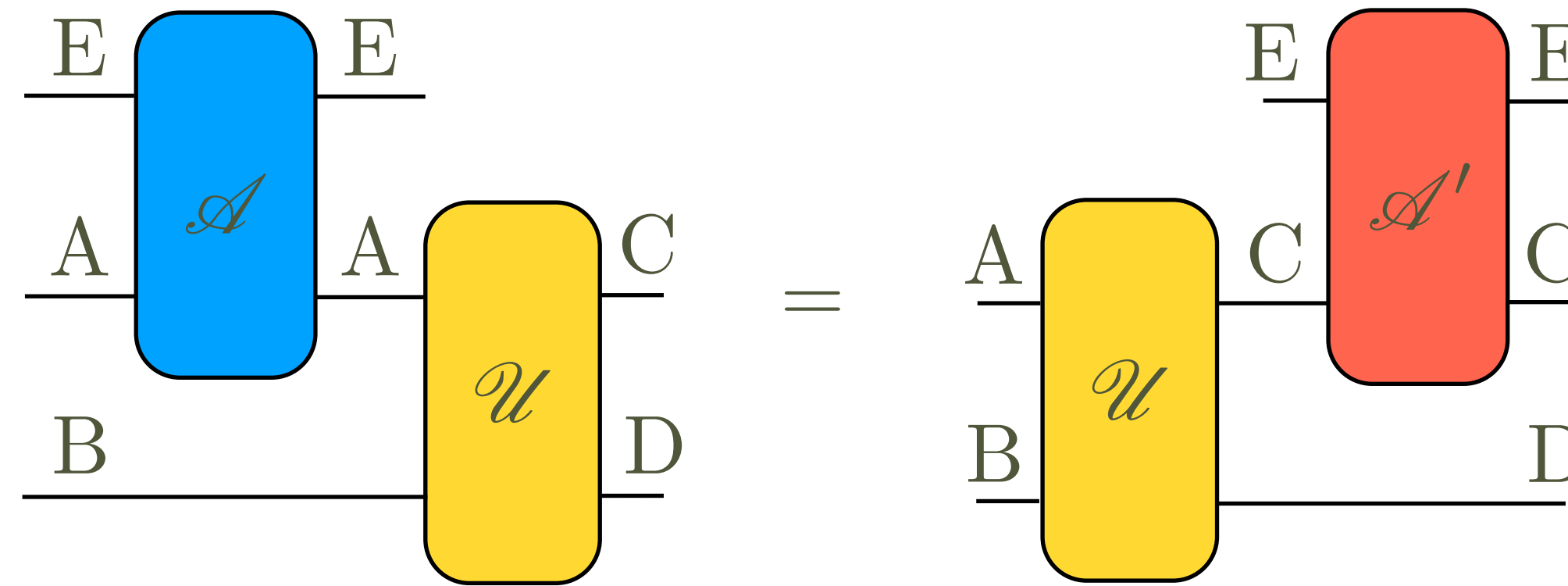
# Necessary condition for no C.I.: no-signalling

Let  $A \not\leftrightarrow D$  :

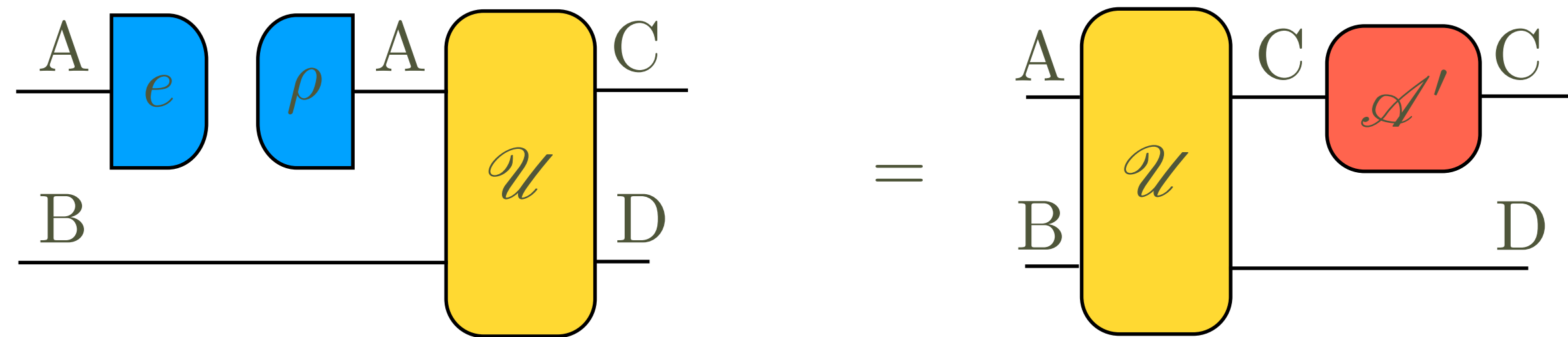


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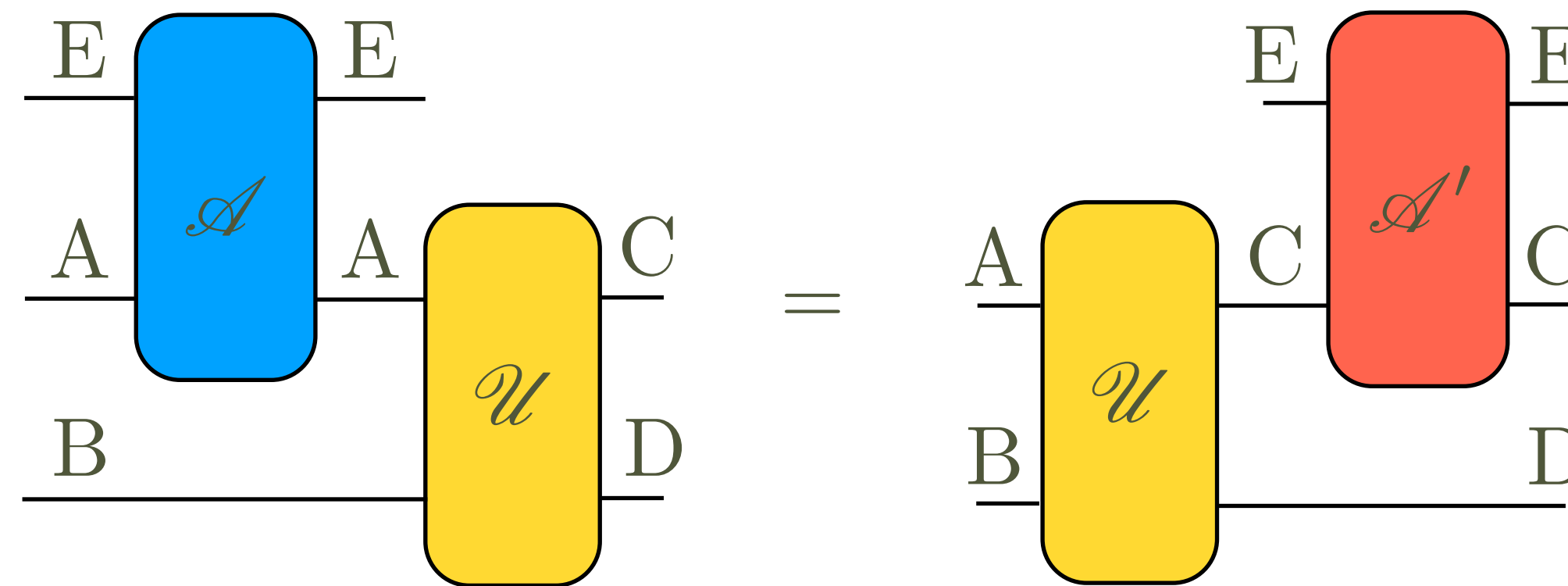


Consider  $E = I$ ,  $\begin{array}{c} A \\ \text{---} \end{array} \mathcal{A} \begin{array}{c} \text{---} \\ A \end{array} = \begin{array}{c} A \\ \text{---} \end{array} e \begin{array}{c} \text{---} \\ \rho \end{array} \begin{array}{c} \text{---} \\ A \end{array}$

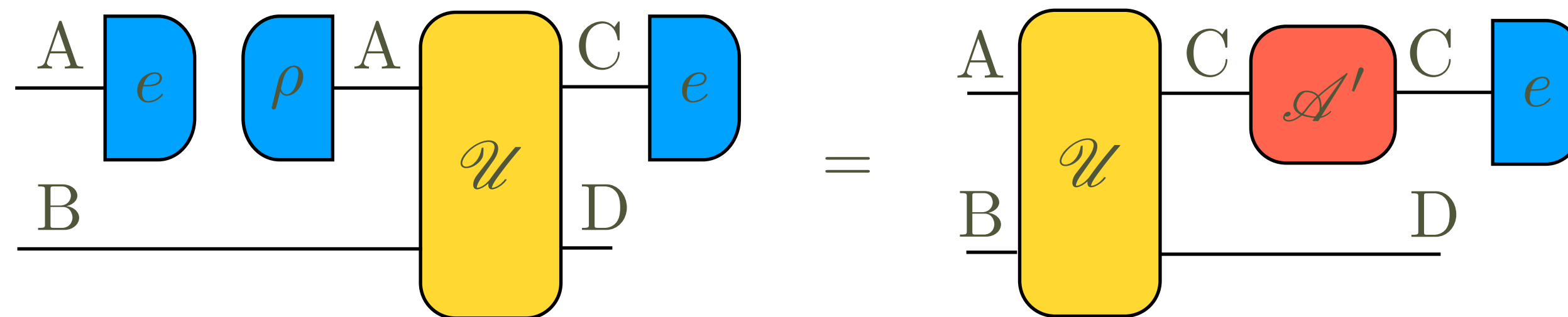


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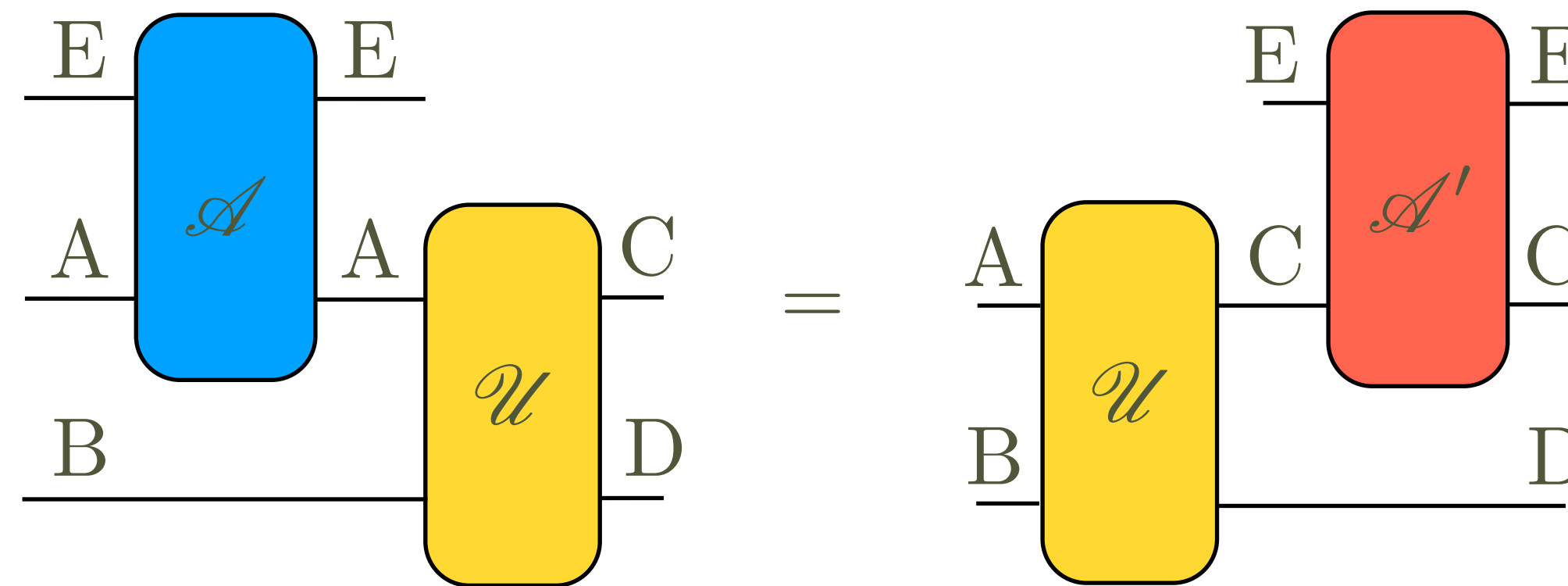


Consider  $E = I$ ,  $\frac{A}{\quad} \mathcal{A} \frac{A}{\quad} = \frac{A}{\quad} e \quad \rho \frac{A}{\quad}$  then discard  $C$

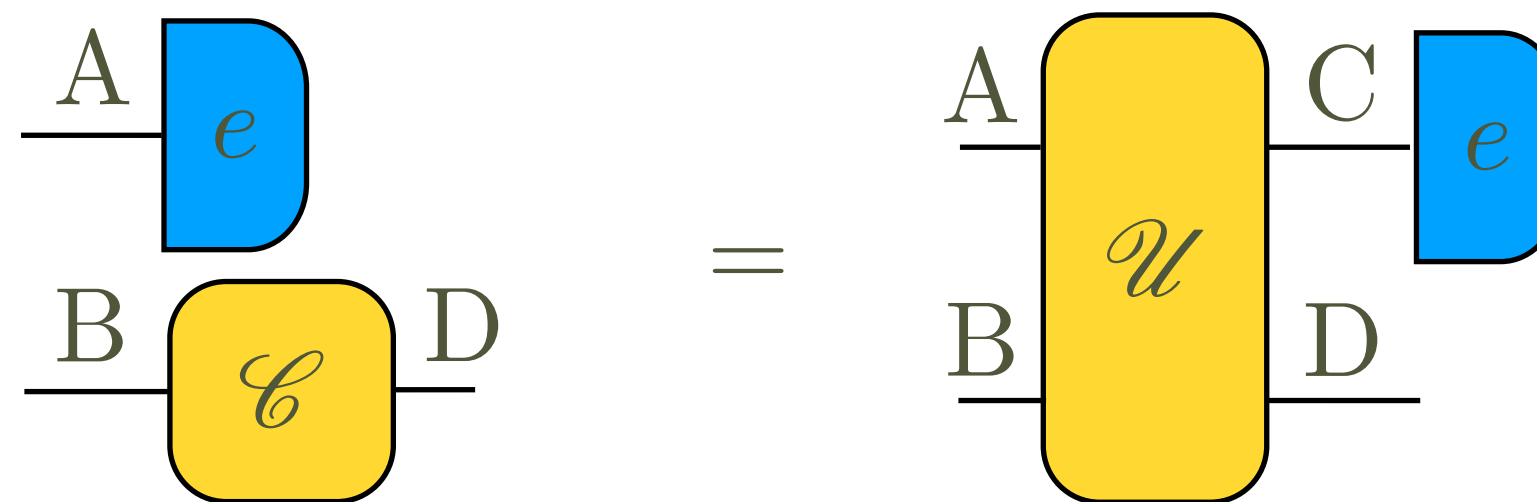


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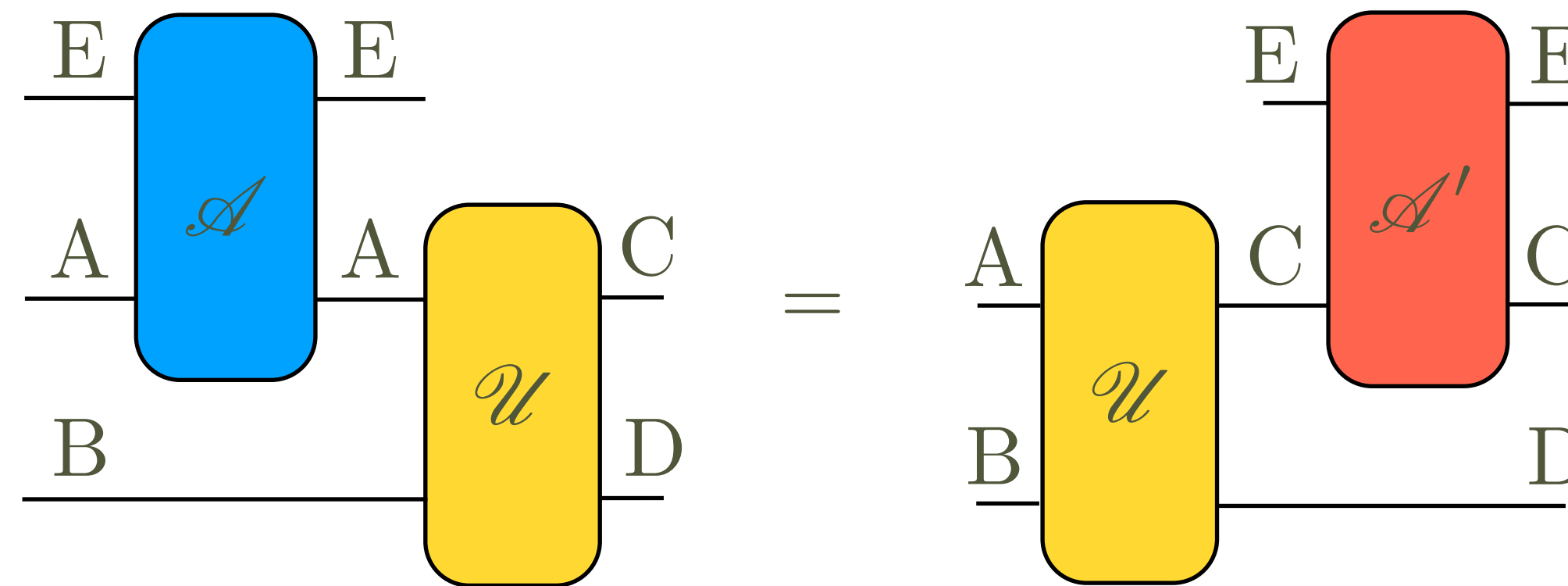


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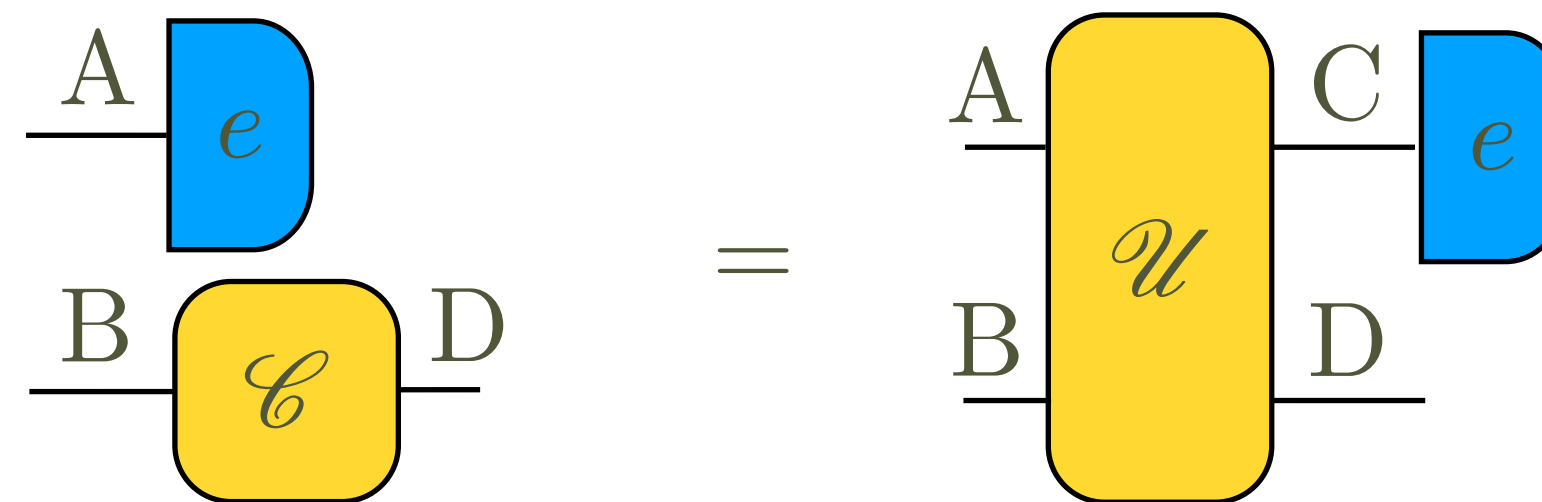


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$\rightarrow A \not\leftrightarrow D$

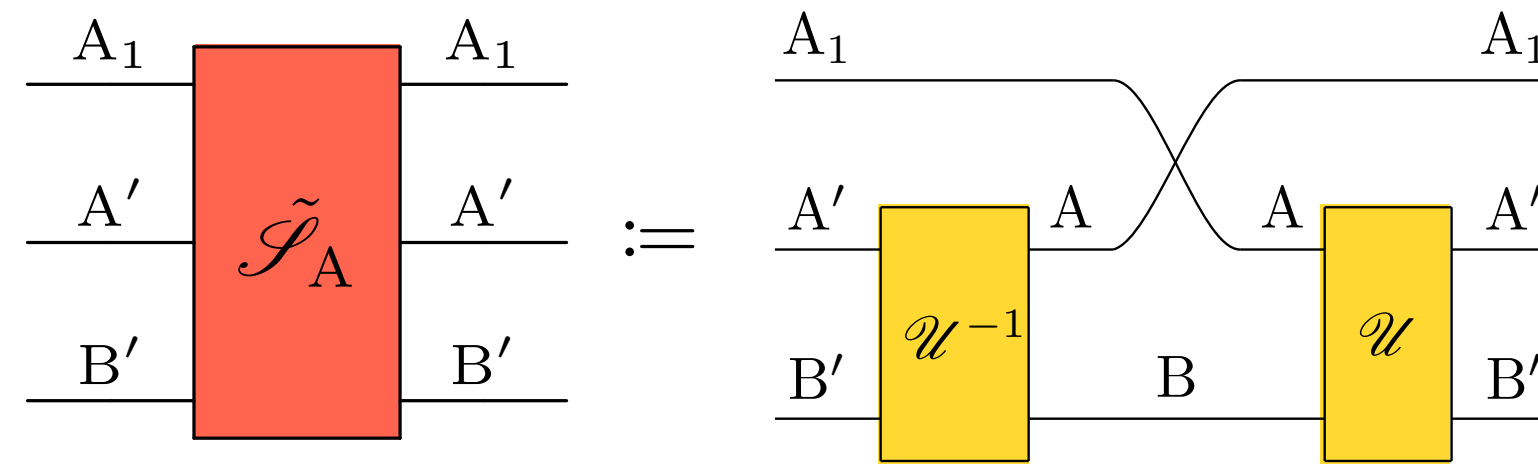
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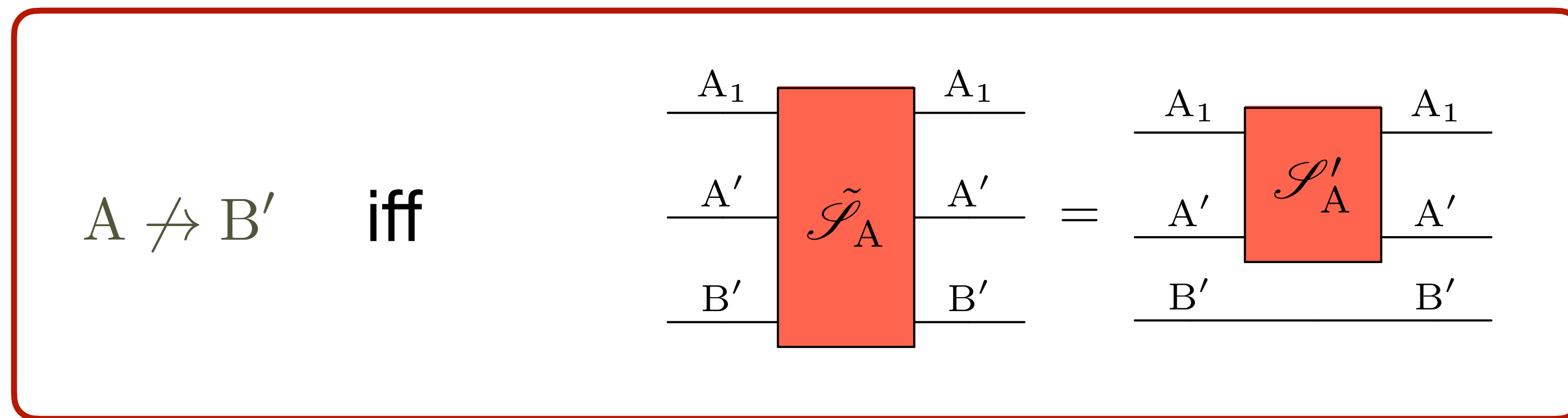
  $A \not\leftrightarrow D$

# Necessary and sufficient condition

- Definition:

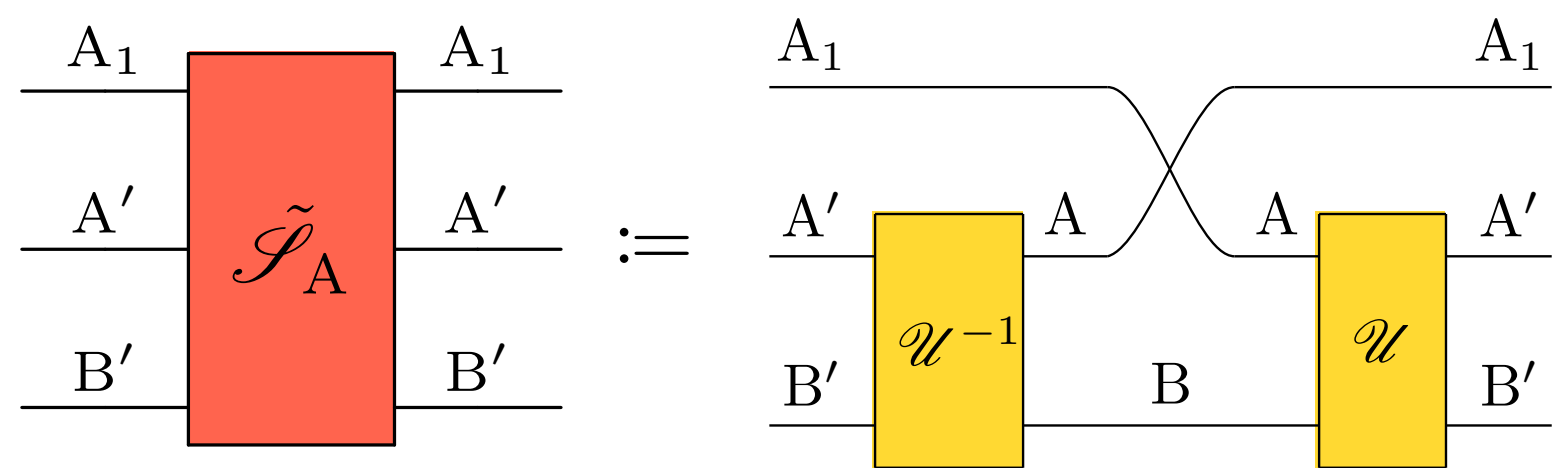


- Condition:



# Necessary and sufficient condition

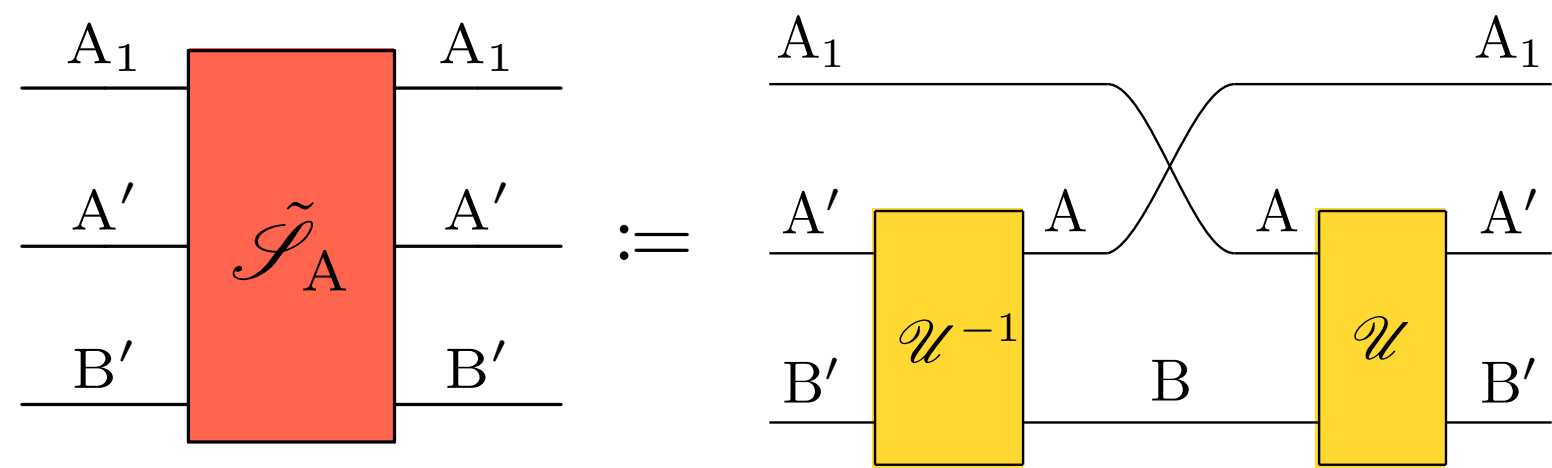
## Proof



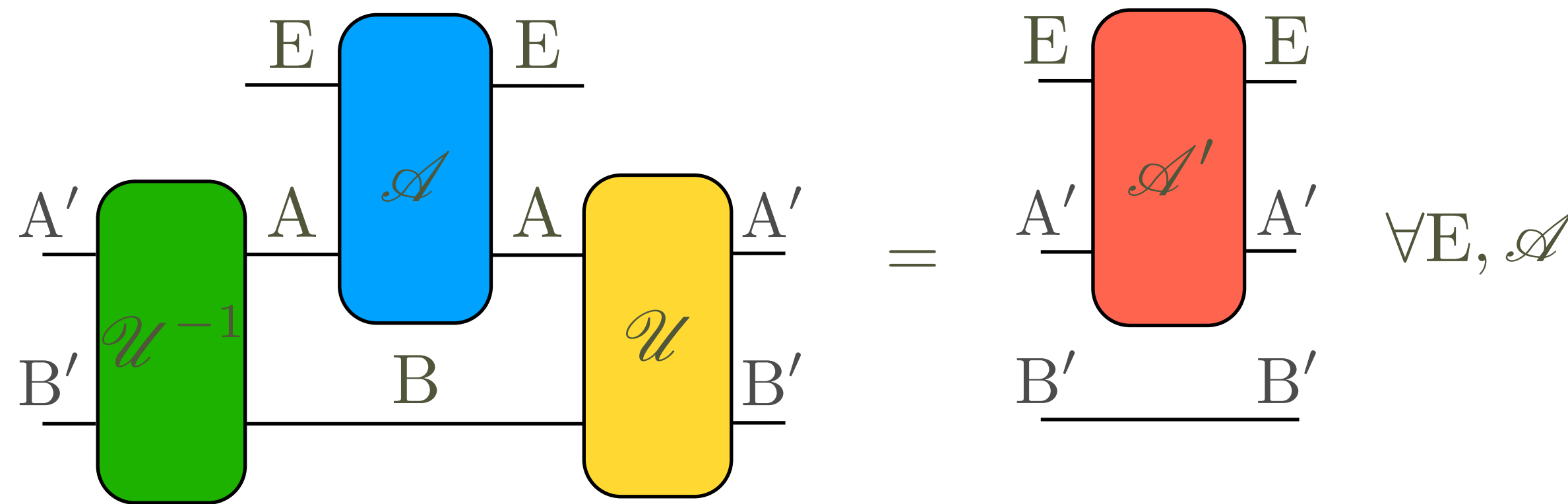


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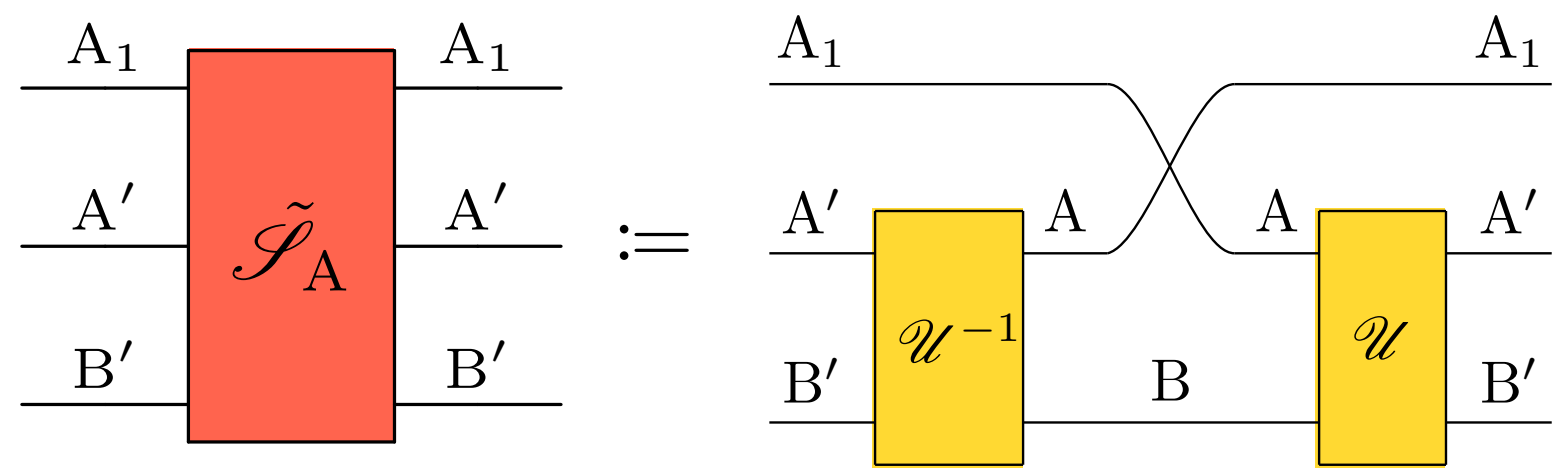


If  $A \not\rightarrow B'$

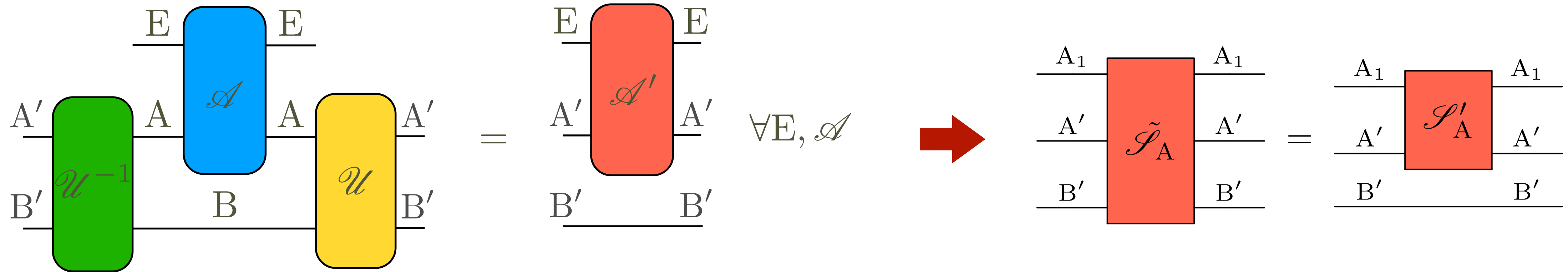


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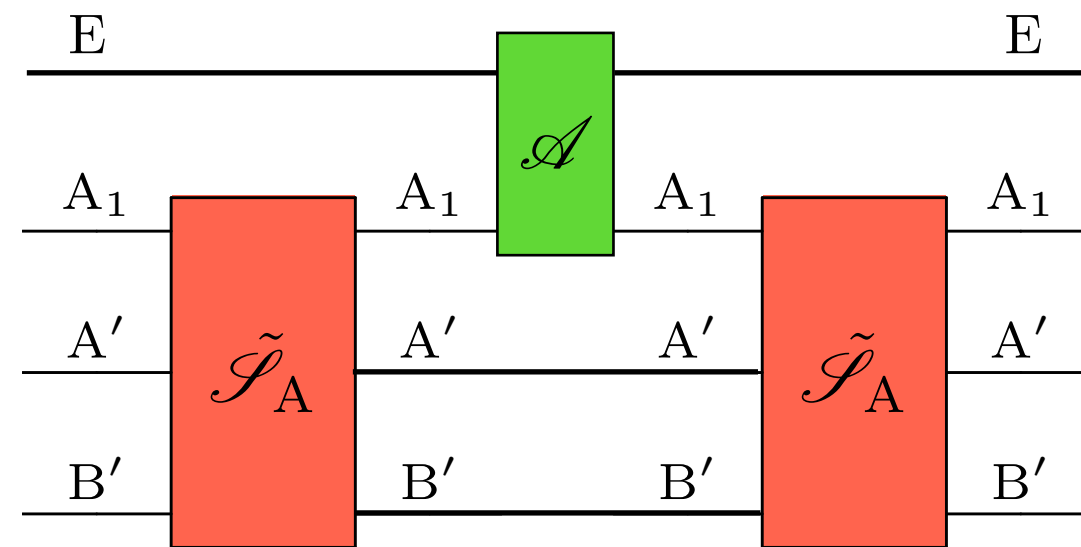


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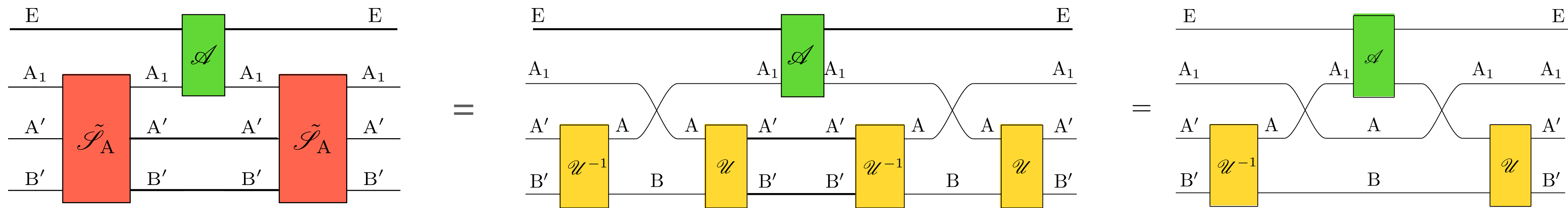
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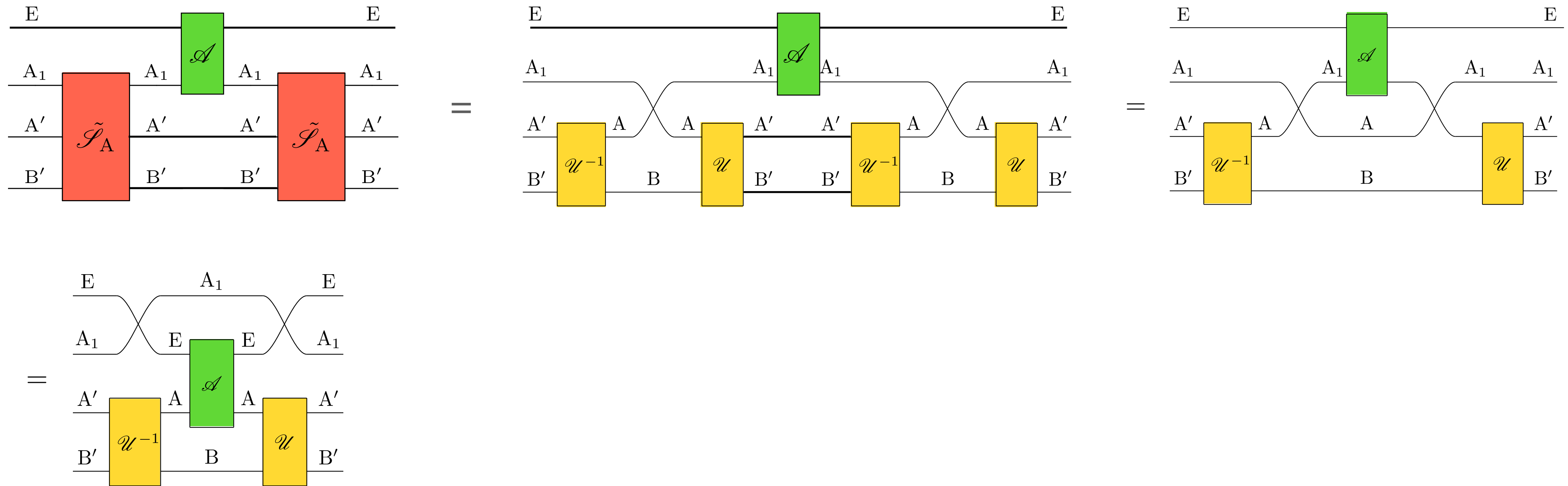
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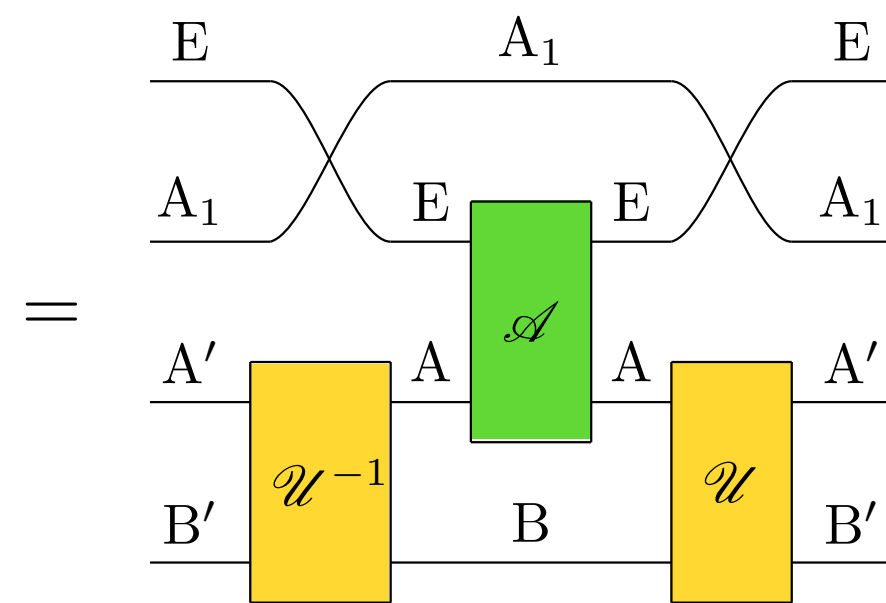
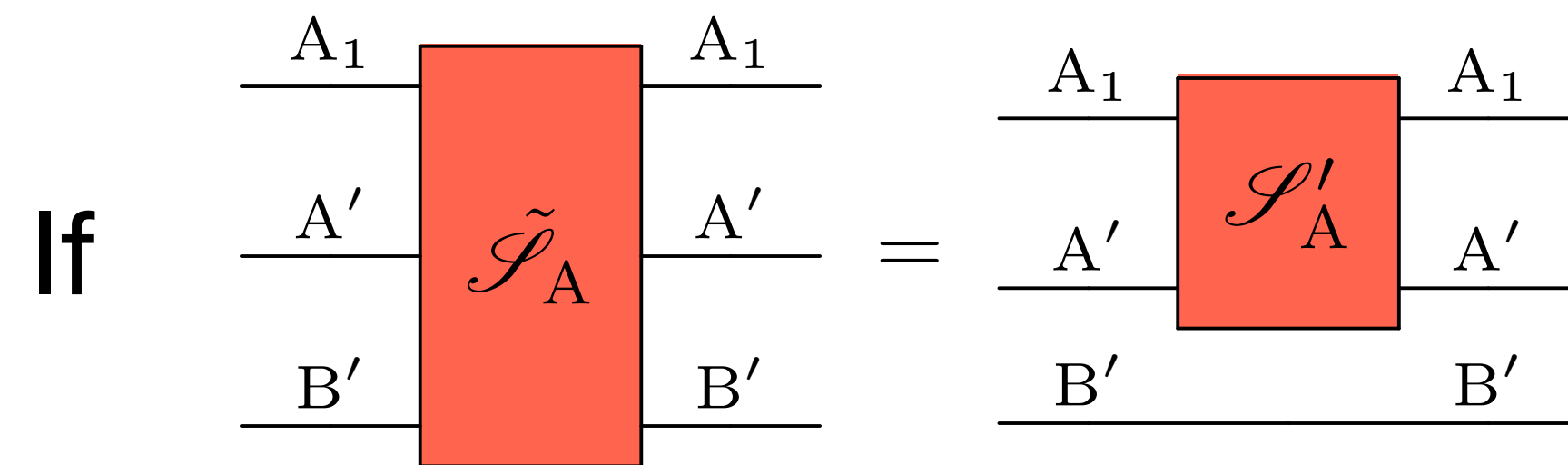
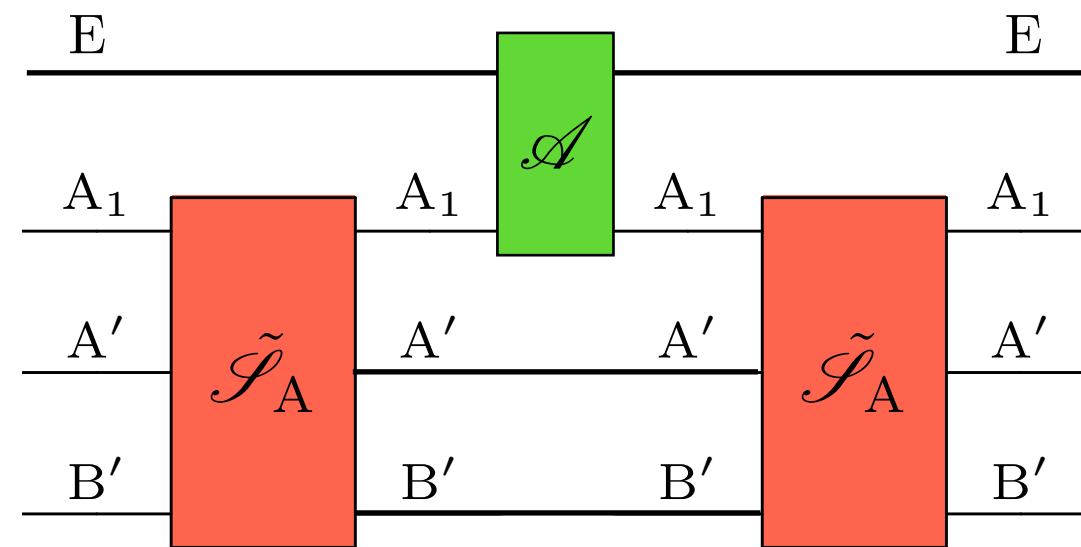
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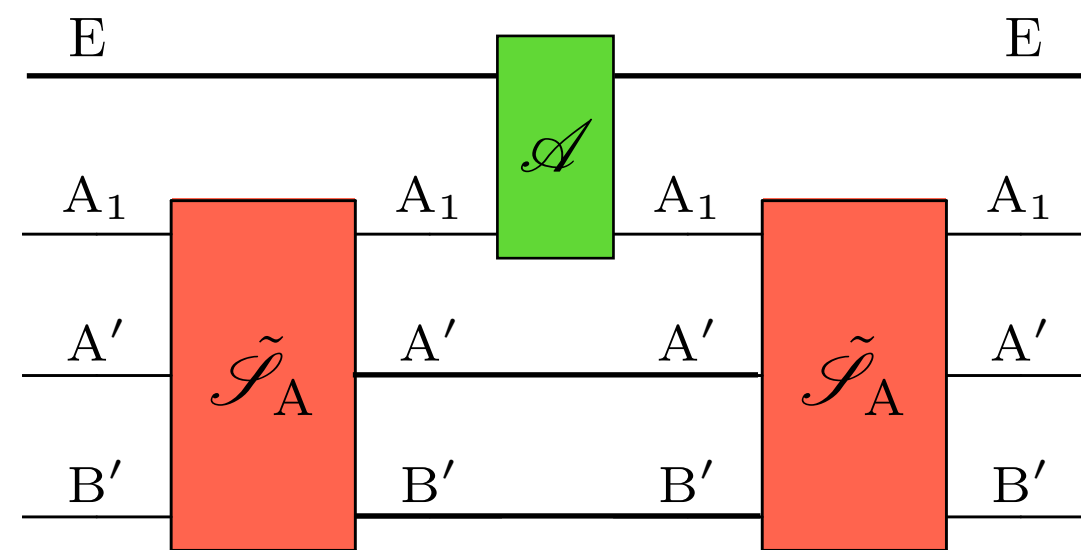
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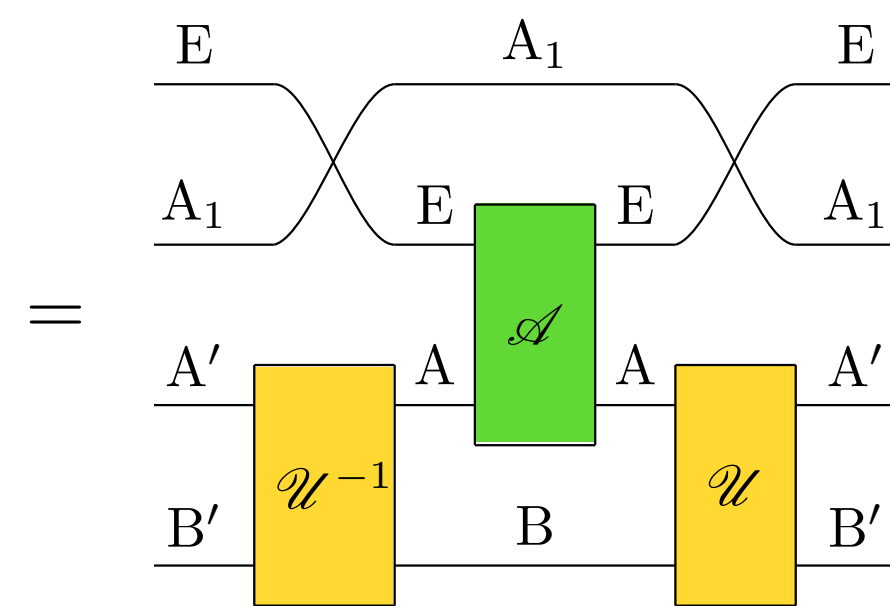
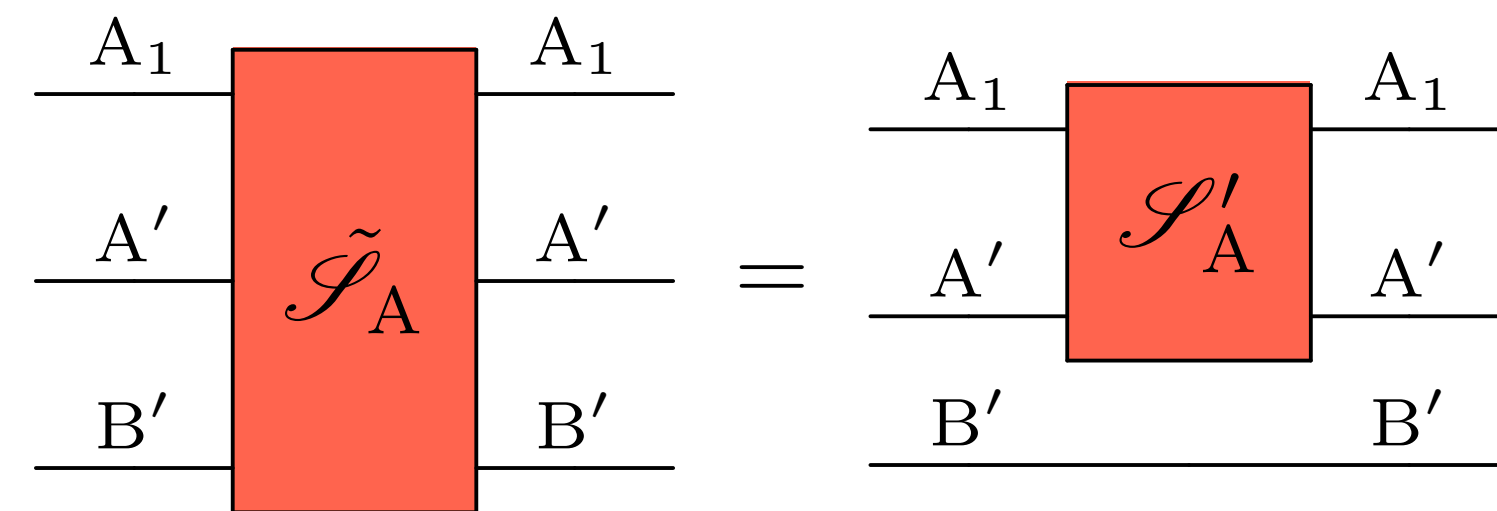


# Necessary and sufficient condition

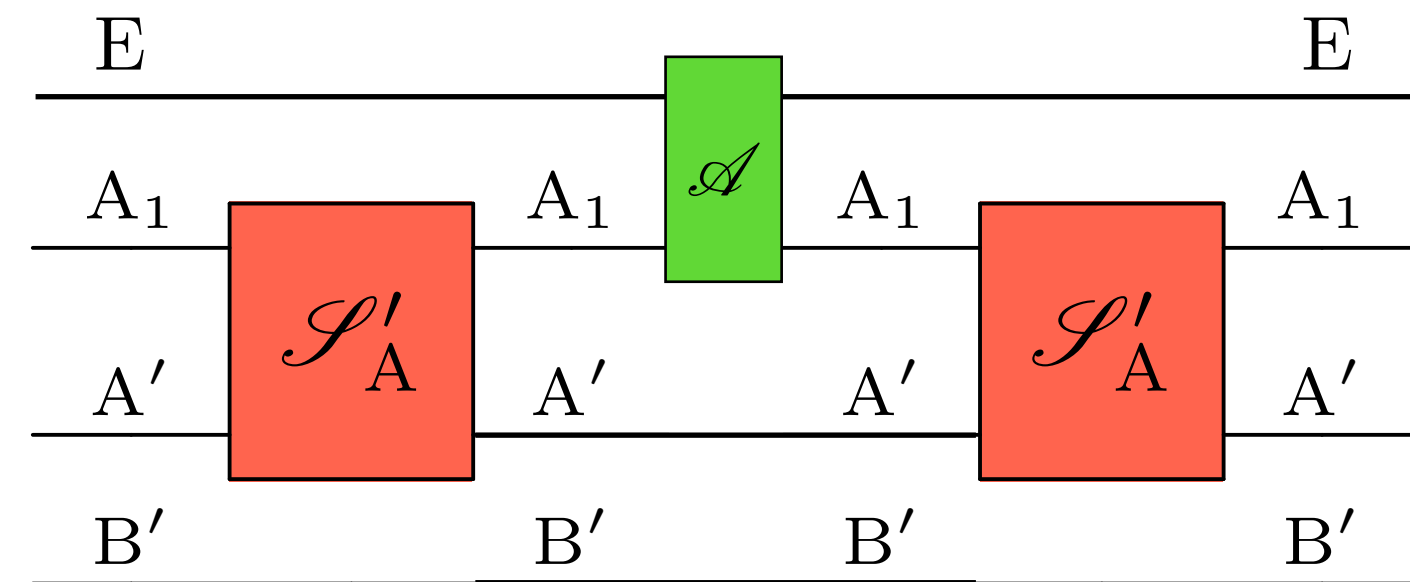
## Proof



If

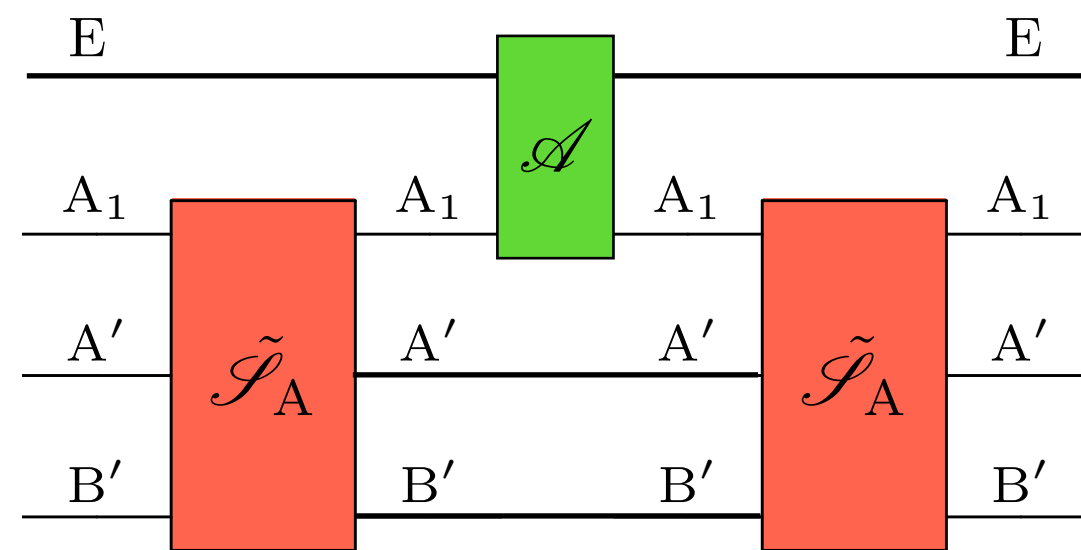


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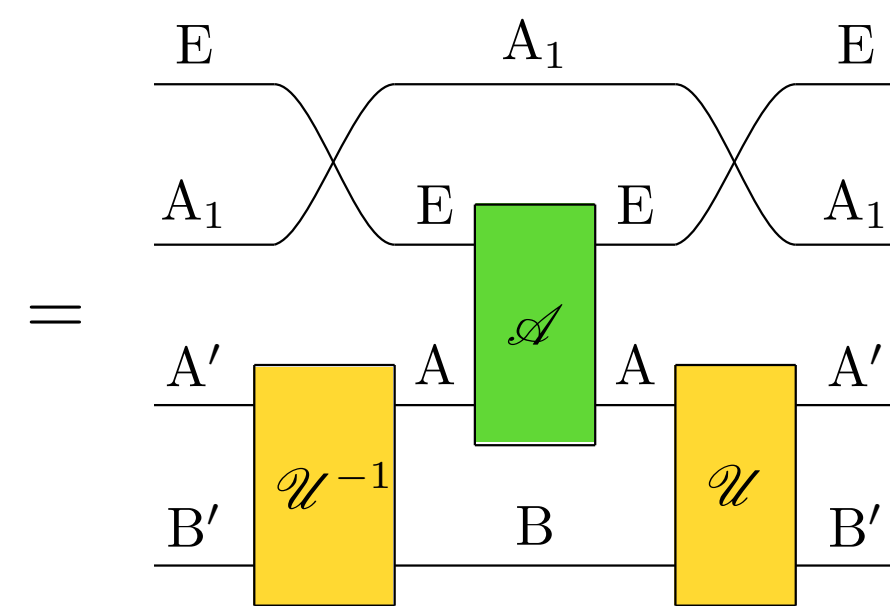
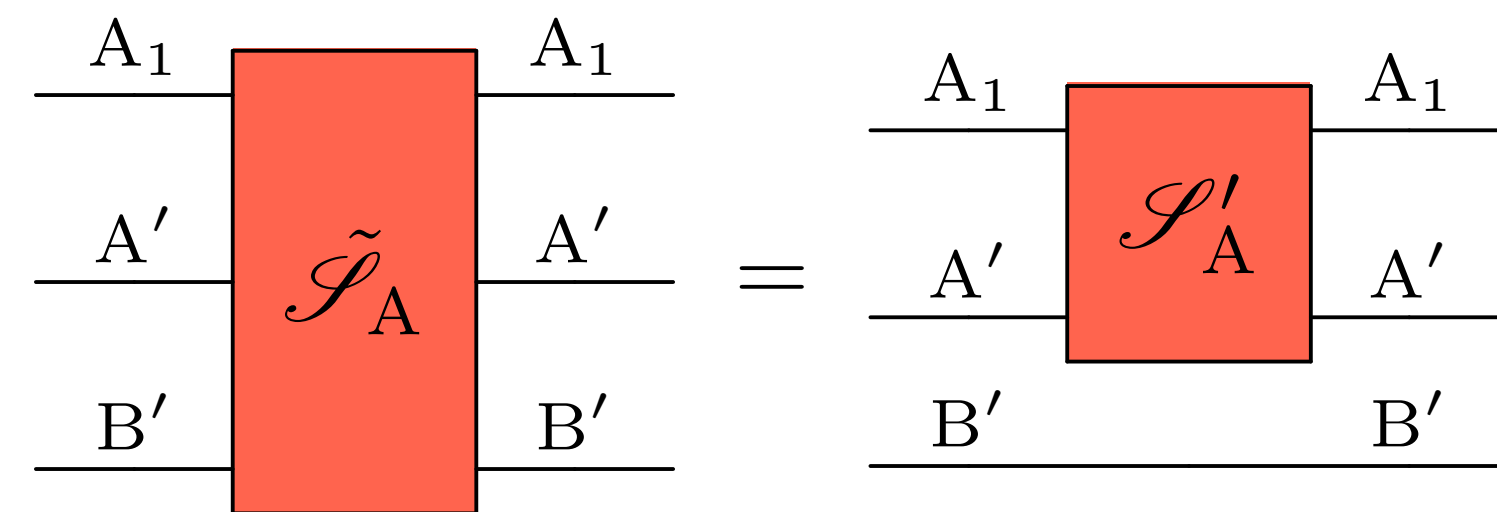


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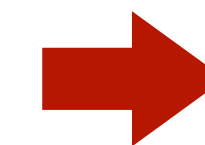
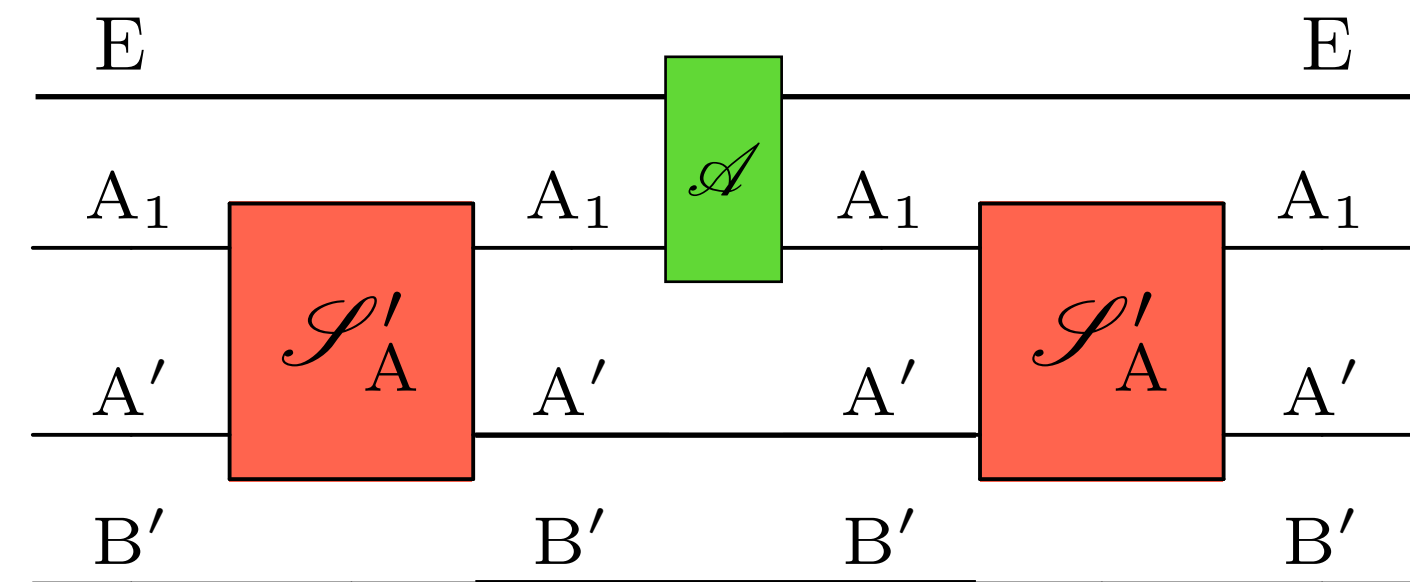
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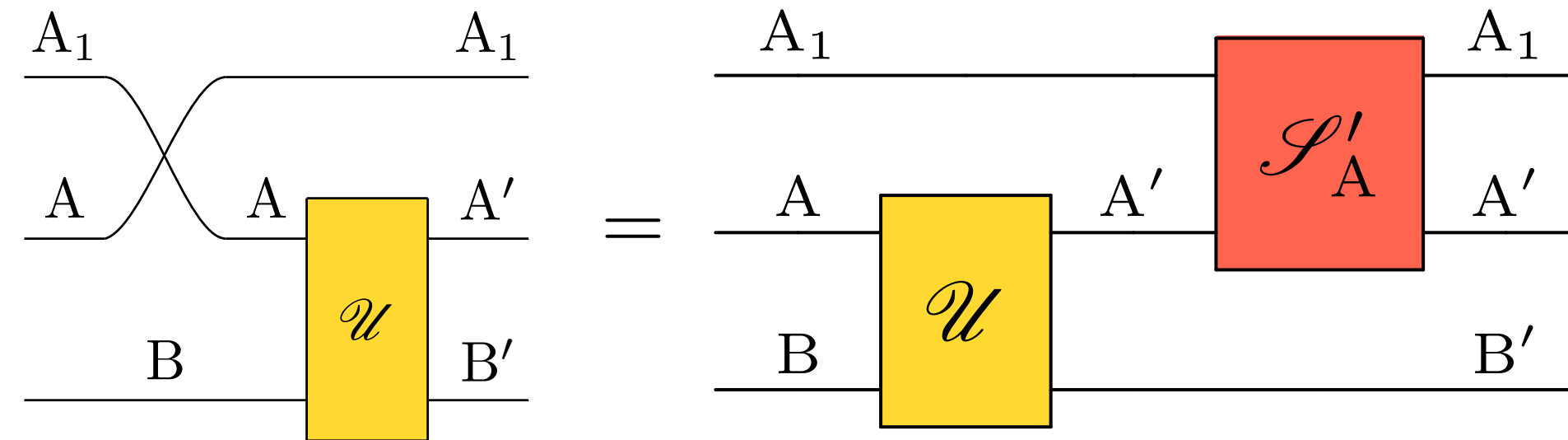


$A \not\leftrightarrow B'$



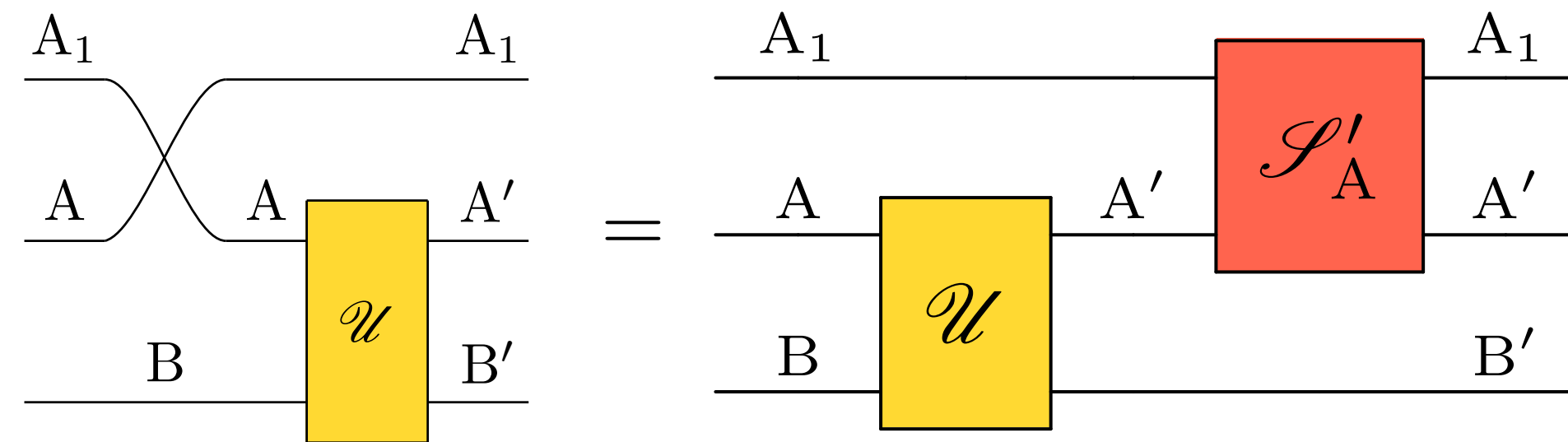
# Necessary condition: comb structure

- Suppose that  $A \not\rightarrow B'$ . Then it must be

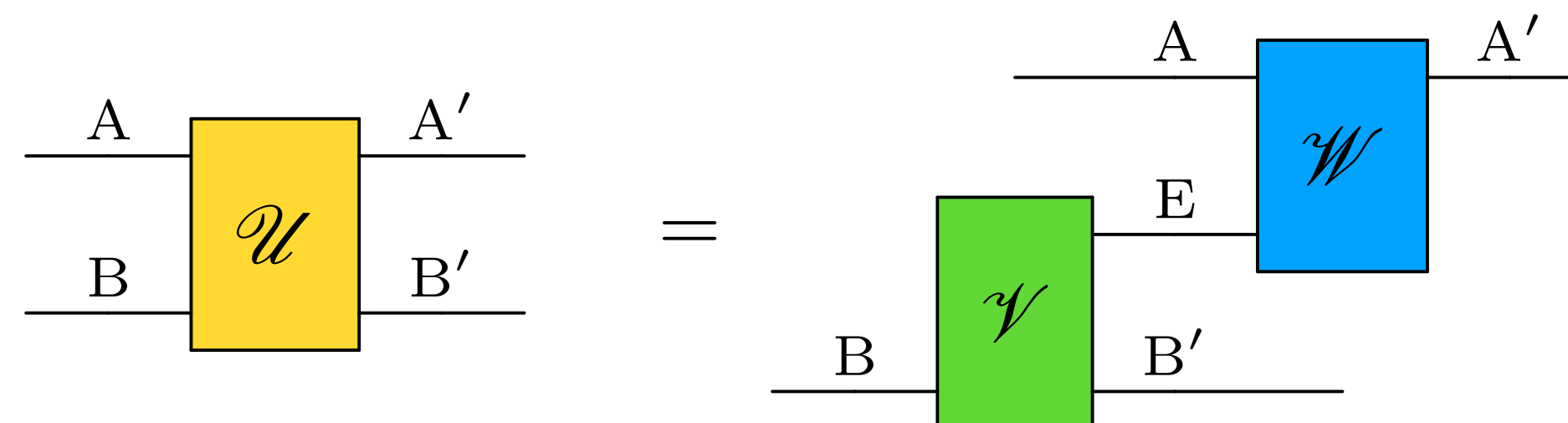


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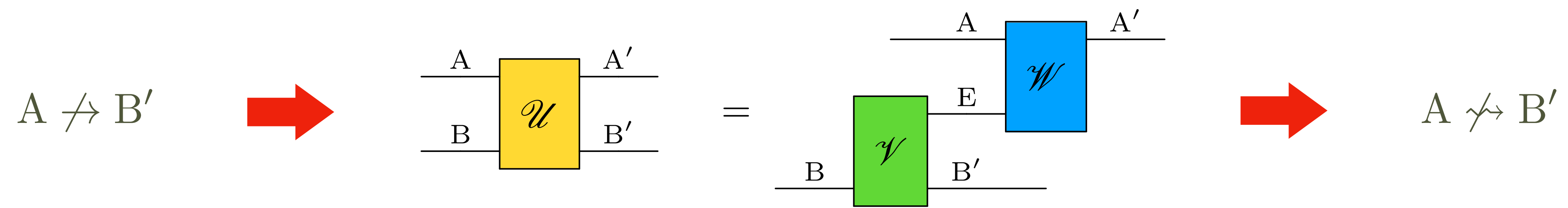
- Suppose that  $A \not\rightarrow B'$ . Then it must be



- Preparing a state of  $A$  and discarding  $A_1$  we obtain that



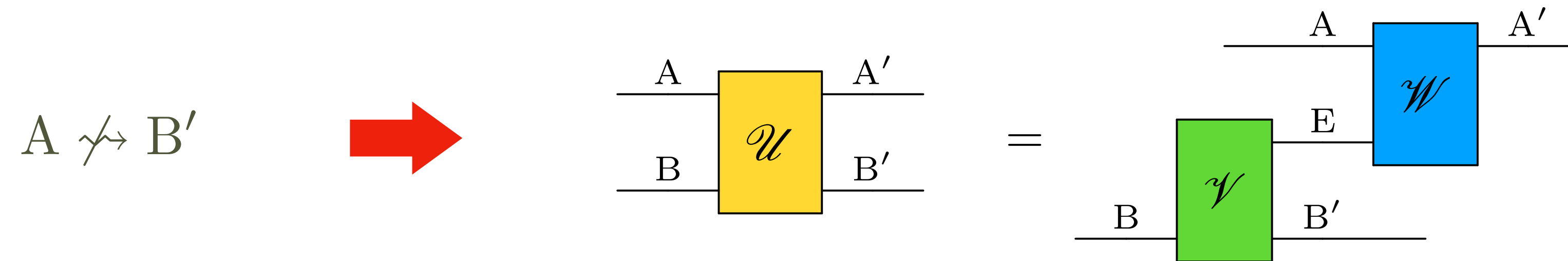
# Chain of conditions



# Classical theory

## Example 1

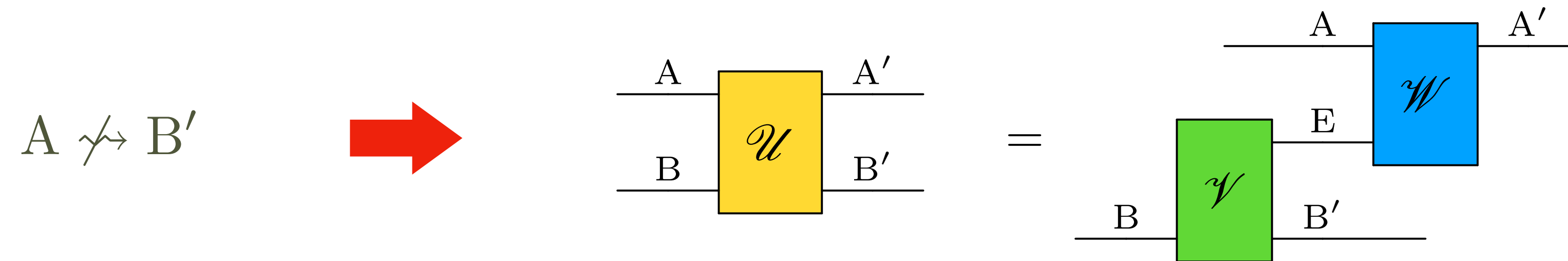
- One can prove that **in classical theory**



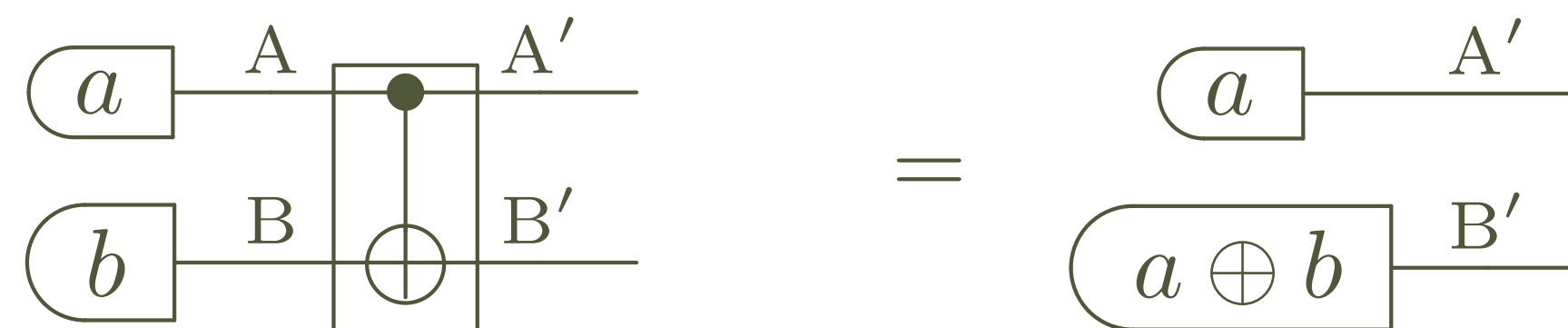
# Classical theory

## Example 1

- One can prove that **in classical theory**



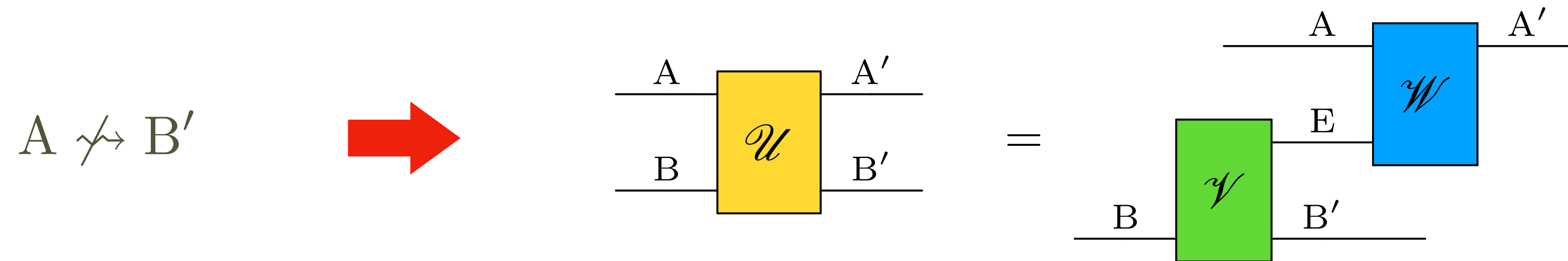
- Classical C-not



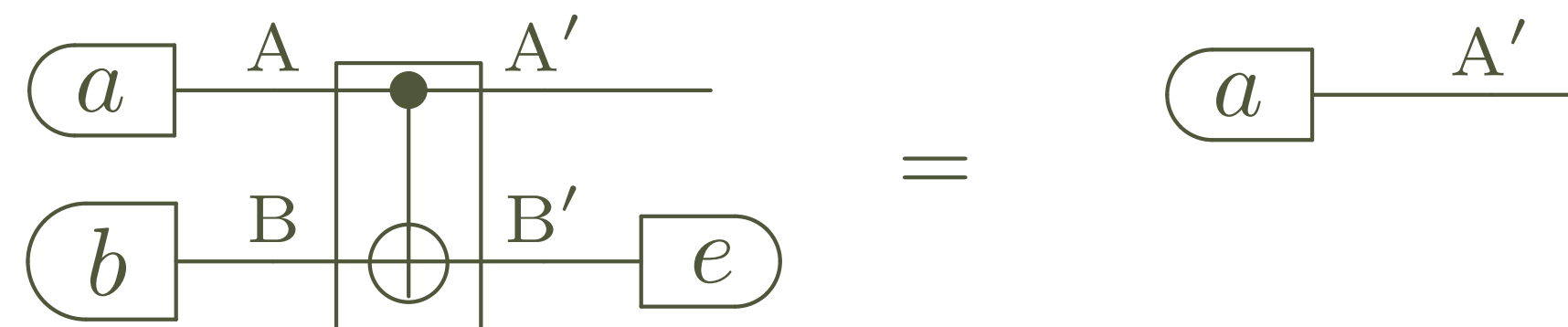
# Classical theory

## Example 1

- One can prove that **in classical theory**



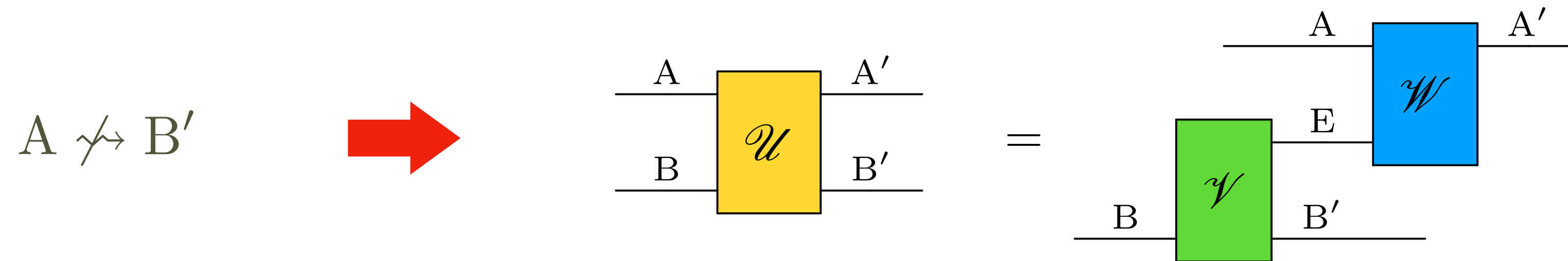
- Classical C-not



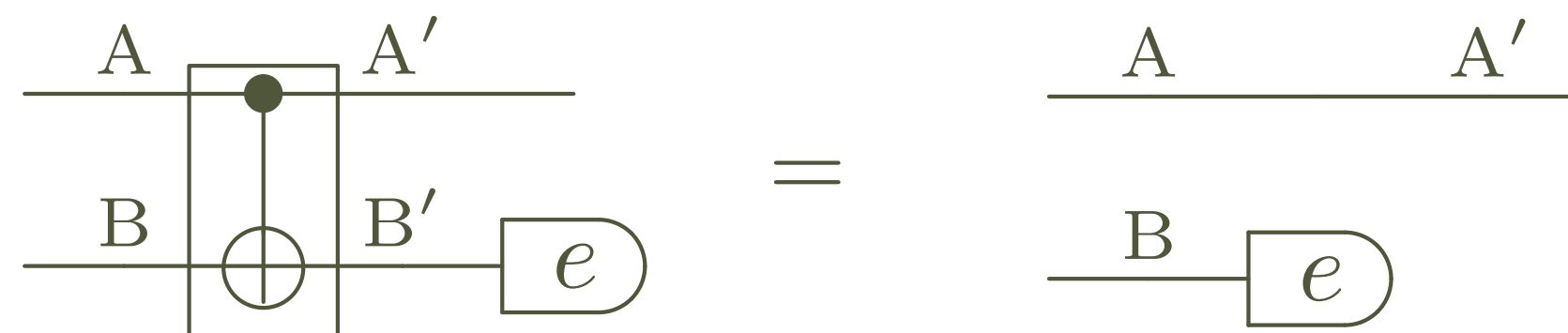
# Classical theory

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- One can prove that **in classical theory**



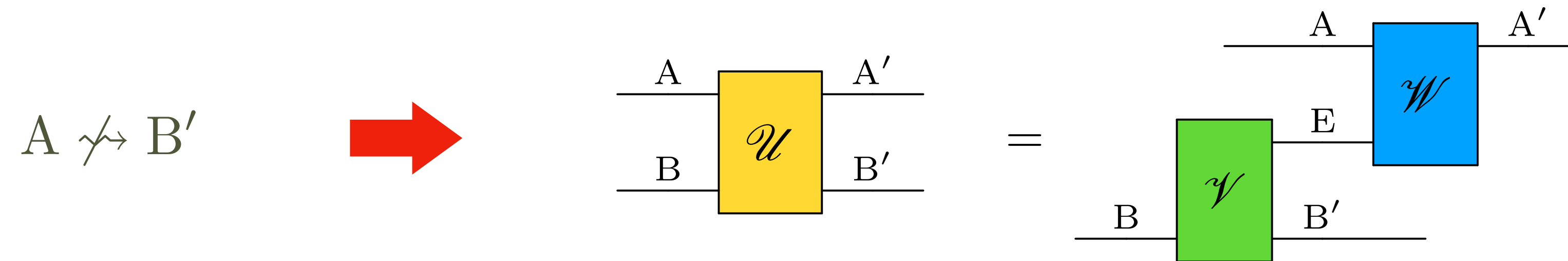
- Classical C-not



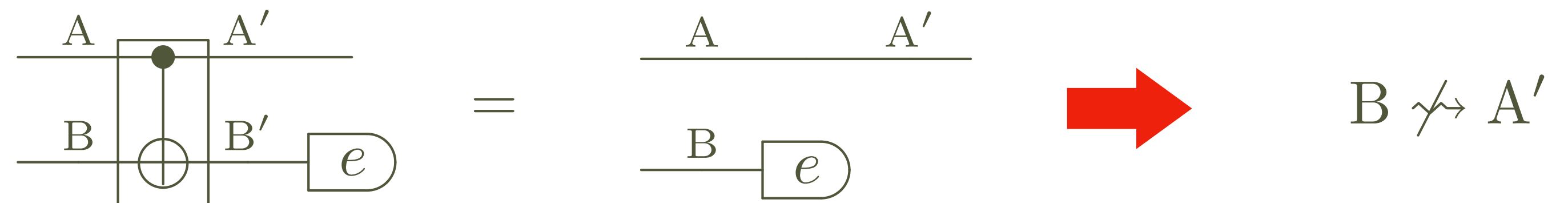
# Classical theory

## Example 1

- One can prove that **in classical theory**



- Classical C-not

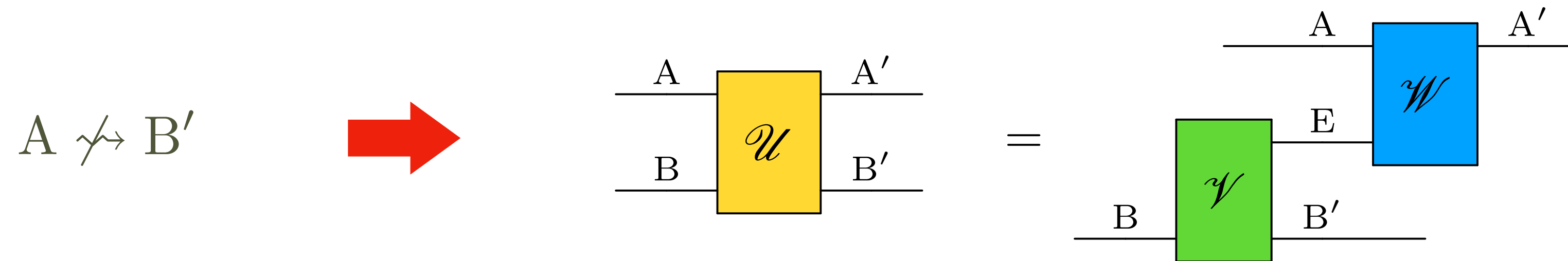




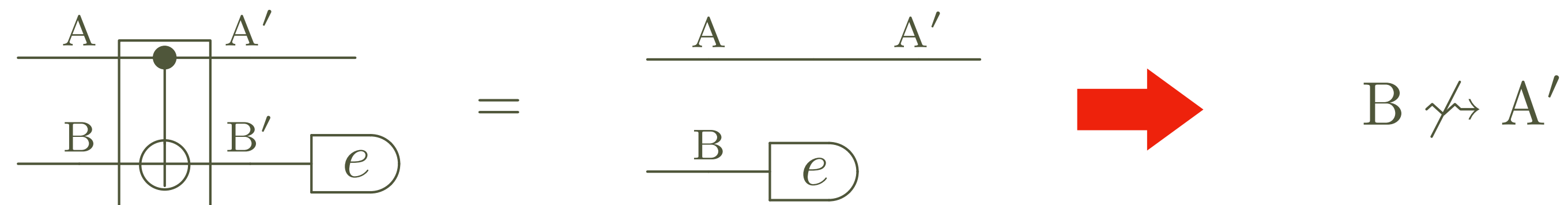
# Classical theory

## Example 1

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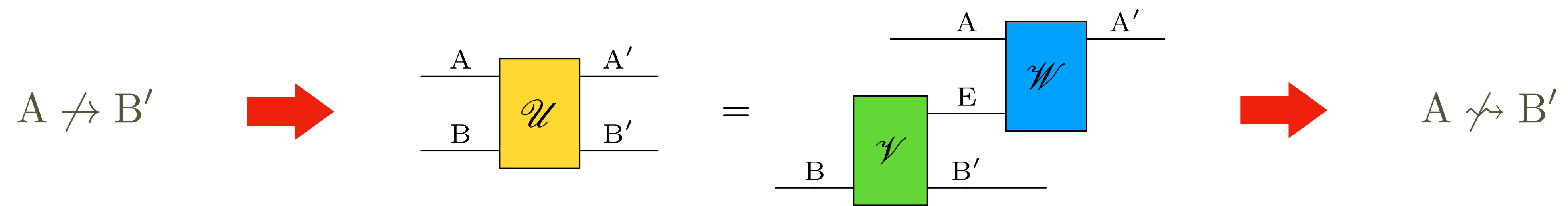
- Classical C-not



- However,  $\tilde{\mathcal{I}}_B \neq \mathcal{I}'_B \otimes \mathcal{I}_{A'}$ , thus  $B \rightarrow A'$

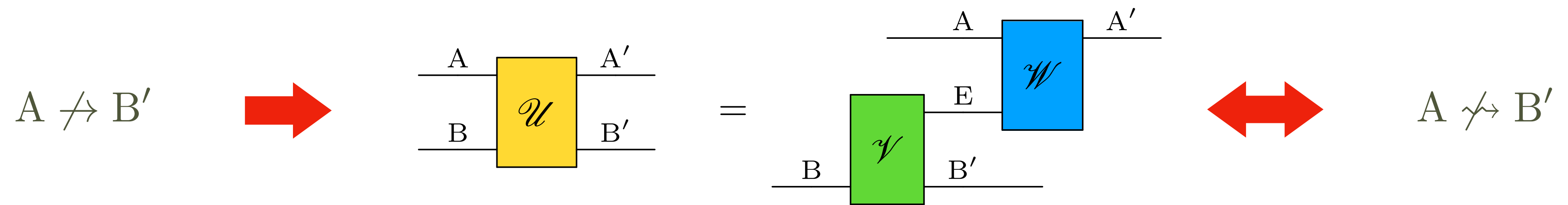
# Chain of conditions

## In classical theory



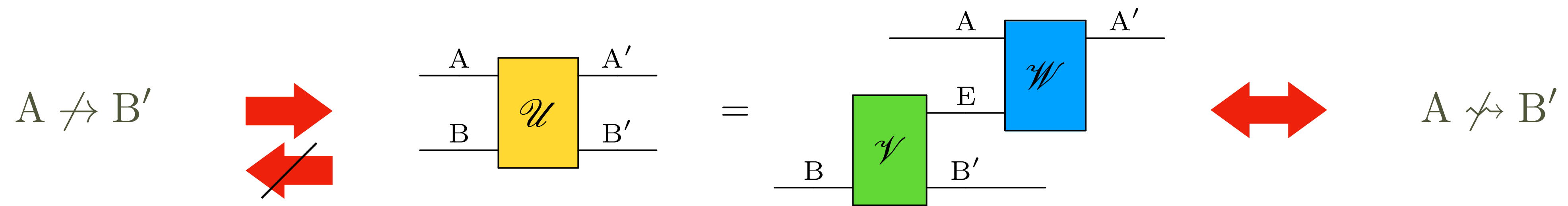
# Chain of conditions

## In classical theory



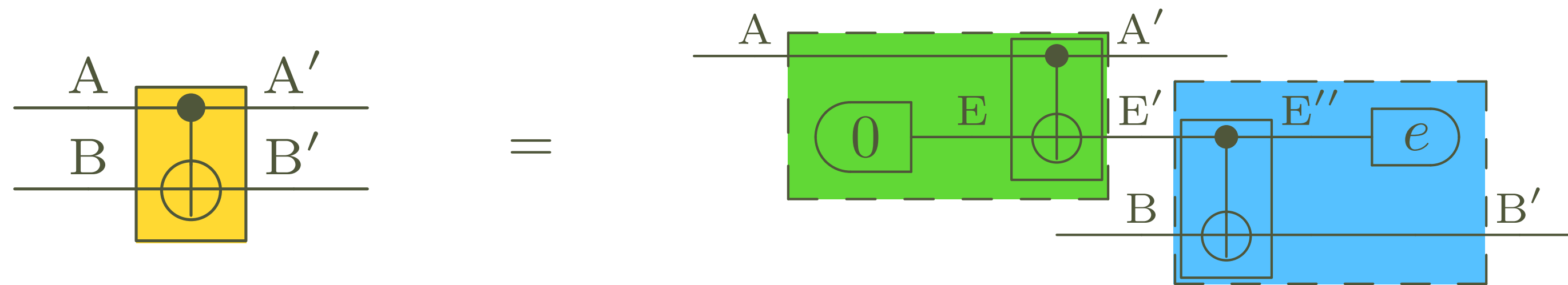
# Chain of conditions

## In classical theory



# Again on the classical C-not

## Comb structure

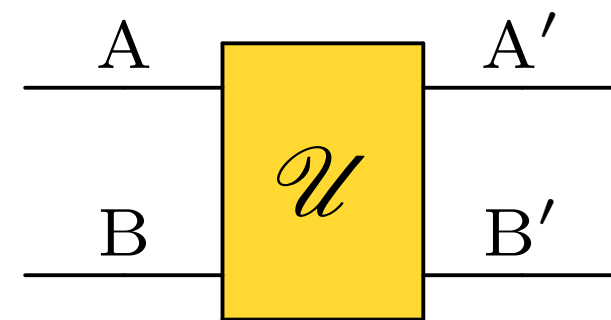
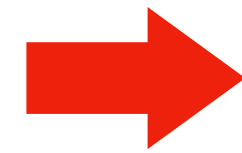


# Quantum theory

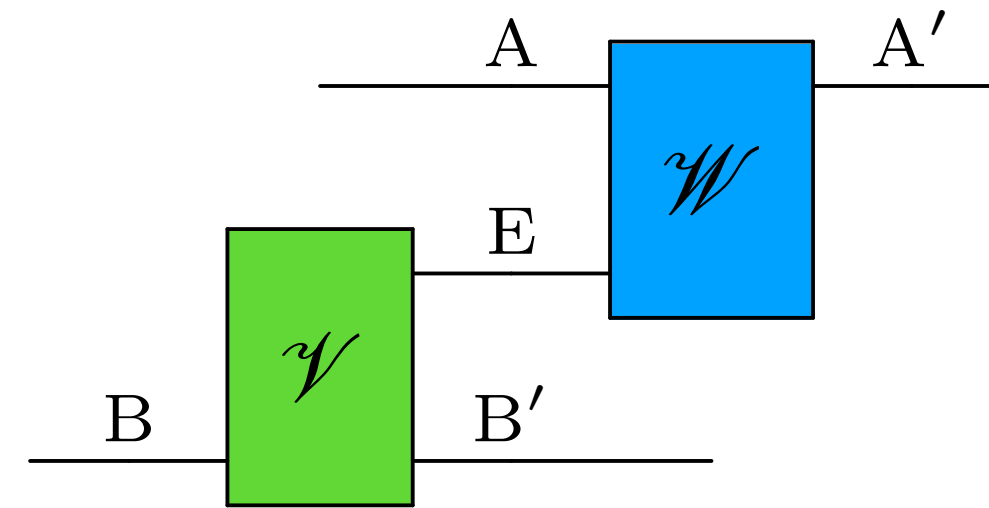
## Example 2

- Also in quantum theory

$A \not\rightarrow B'$



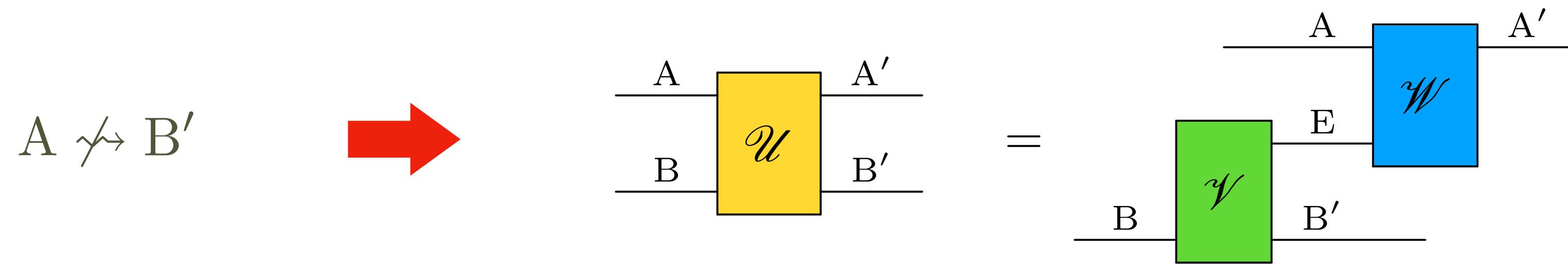
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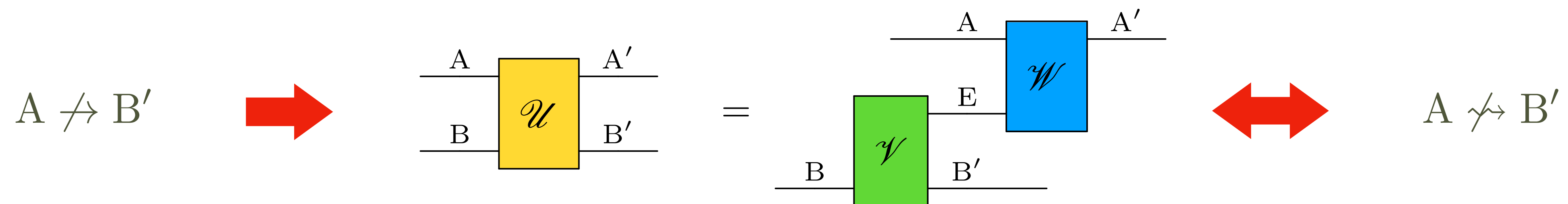
# Quantum theory

## Example 2

- Also in quantum theory



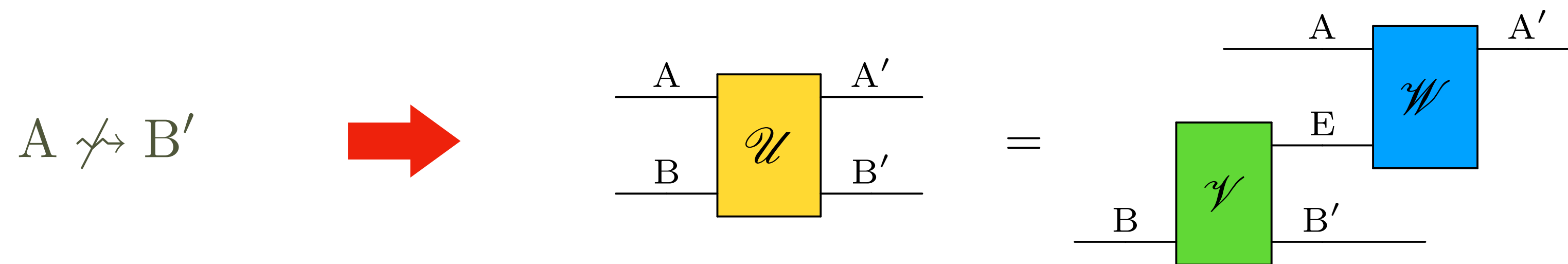
- Thus



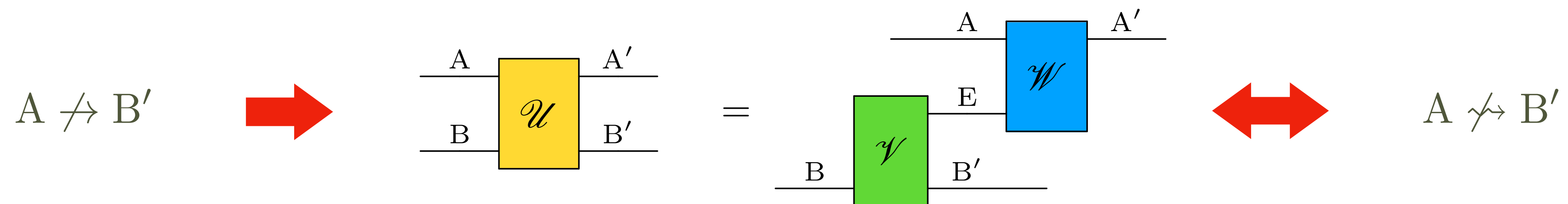
# Quantum theory

## Example 2

- Also in quantum theory



- Thus

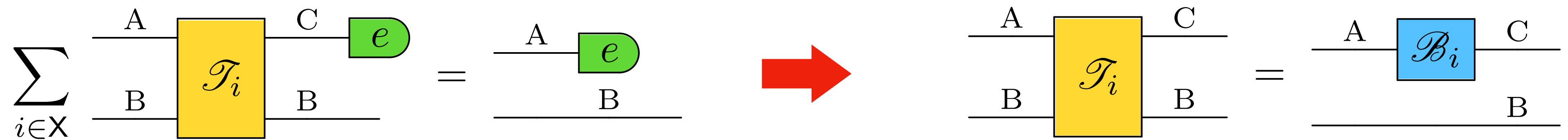


- What about the first implication?



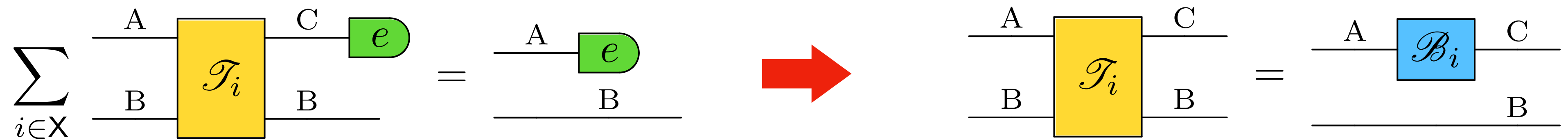
# Quantum theory

- From the characterisation of Kraus decompositions of a given channel



# Quantum theory

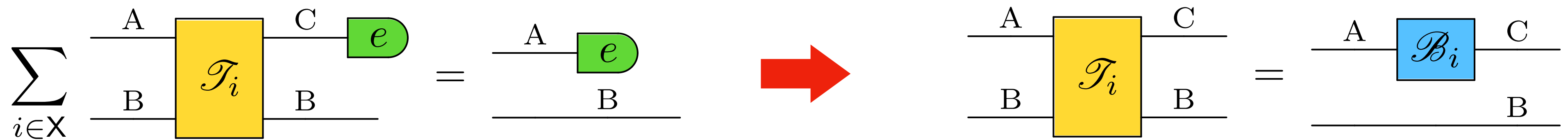
- From the characterisation of Kraus decompositions of a given channel



- Also from **purification**

# Quantum theory

- From the characterisation of Kraus decompositions of a given channel

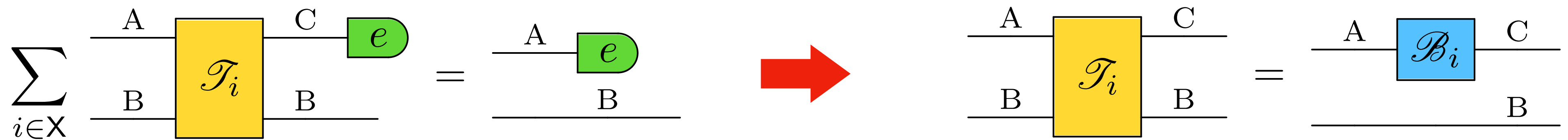


- Also from **purification**

➔ The above result holds also in Fermionic theory and Real Quantum theory

# Quantum theory

- From the characterisation of Kraus decompositions of a given channel



- Also from **purification**

➔ The above result holds also in Fermionic theory and Real Quantum theory

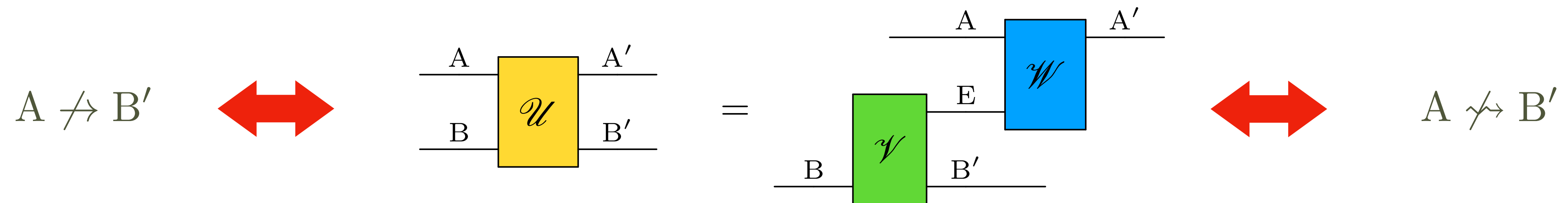
**No interaction without disturbance**

# Quantum theory

- From no interaction without disturbance one has  $A \not\leftrightarrow B \rightarrow A \leftrightarrow B$

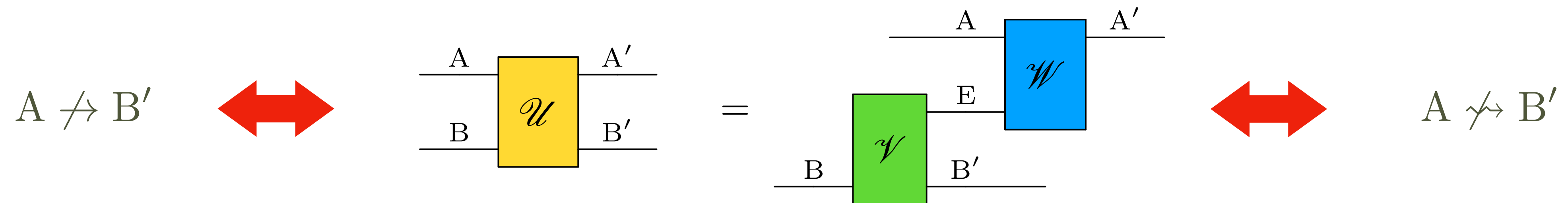
# Quantum theory

- From no interaction without disturbance one has  $A \not\leftrightarrow B \implies A \not\leftrightarrow B$
- Thus



# Quantum theory

- From no interaction without disturbance one has  $A \not\leftrightarrow B \implies A \not\leftrightarrow B'$
- Thus



- True in every theory with purification or just **no interaction without disturbance**

# Proof

## Depending on time

$$\underbrace{\begin{array}{c} A \\ \hline \underbrace{B \mid U}_{B'} \end{array}}^{A'} = \underbrace{\begin{array}{c} A \\ \hline \underbrace{B \mid e}_{B'} \end{array}}^{A'} \quad (A \not\sim B) \Rightarrow \underbrace{\begin{array}{c} A' \\ \hline \underbrace{B' \mid \underbrace{A \mid e}_{B'}}_{B'} \end{array}}^{A'} = \underbrace{\begin{array}{c} A' \\ \hline \underbrace{e}_{B'} \end{array}}^{A'}$$

$$\underbrace{\begin{array}{c} A_1 \\ \hline \underbrace{A' \mid \underbrace{e \mid \underbrace{A \mid e}_{B'}}_{B'}}_{B'} \end{array}}^{A_1} = \underbrace{\begin{array}{c} A_1 \\ \hline \underbrace{A' \mid \underbrace{B \mid \underbrace{A \mid e}_{B'}}_{B'}}_{B'} \end{array}}^{A_1} = \underbrace{\begin{array}{c} A_1 \\ \hline \underbrace{A' \mid \underbrace{A \mid e}_{B'}}_{B'} \end{array}}^{A_1} = \underbrace{\begin{array}{c} A_1 \\ \hline \underbrace{e}_{B'} \end{array}}^{A_1}$$

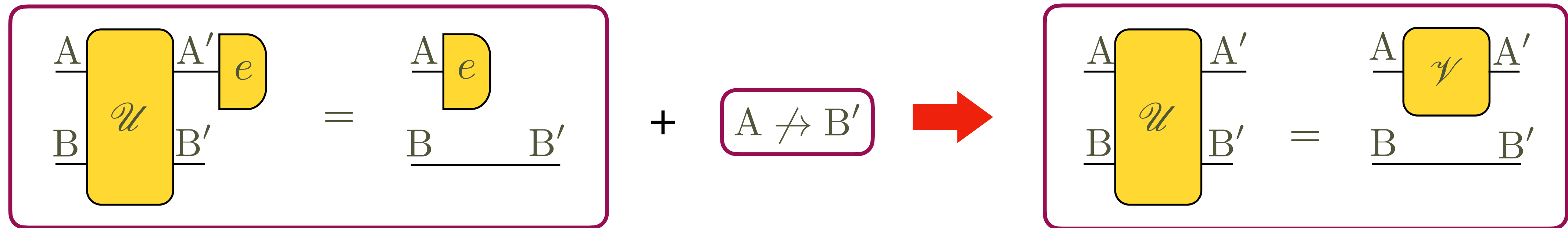
(no int. w.o. dist)

$$\Rightarrow \underbrace{\begin{array}{c} A_1 \\ \hline \underbrace{A' \mid \underbrace{e \mid \underbrace{A \mid e}_{B'}}_{B'}}_{B'} \end{array}}^{A_1} = \underbrace{\begin{array}{c} A_1 \\ \hline \underbrace{A' \mid \underbrace{B \mid \underbrace{A \mid e}_{B'}}_{B'}}_{B'} \end{array}}^{A_1} \Rightarrow A \not\sim B'$$



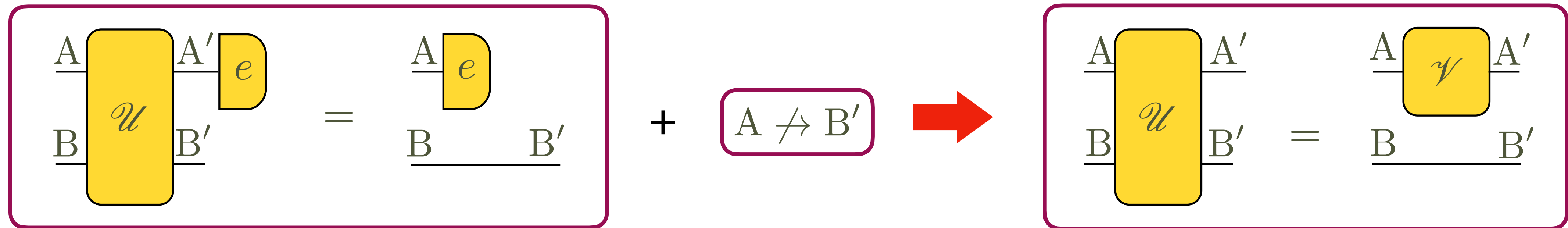
# Interaction without disturbance

- What about a theory featuring **interactions without disturbance**?

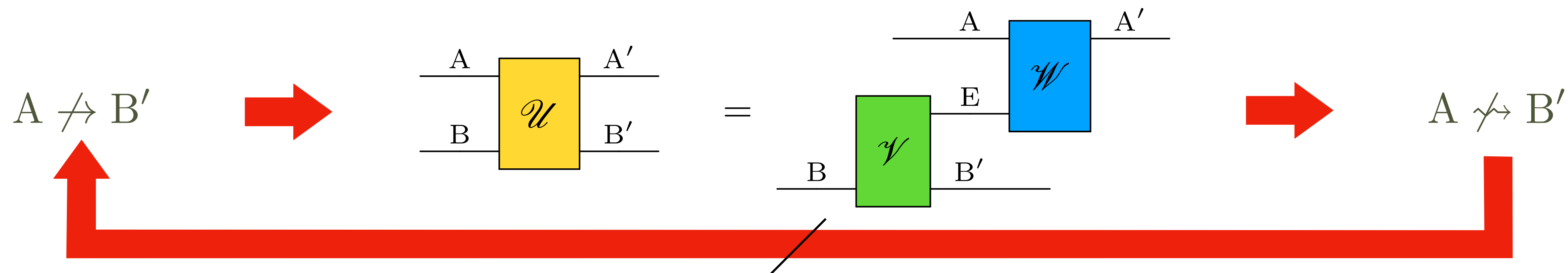


# Interaction without disturbance

- What about a theory featuring **interactions without disturbance**?



- Thus, if the special interaction without disturbance is reversible, one has

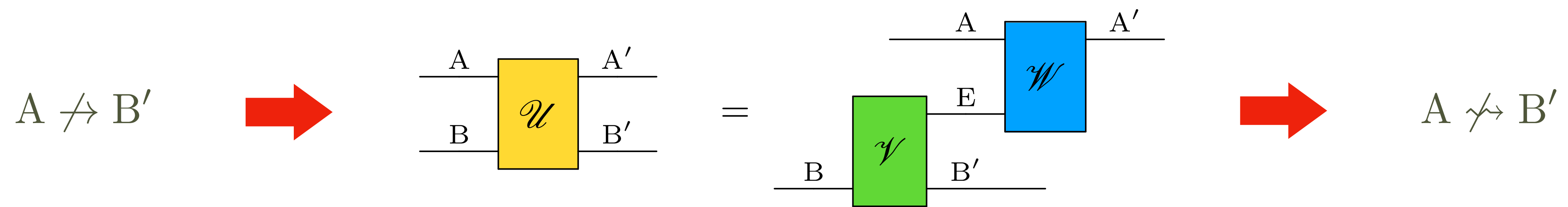


# Open question

**The quest for counterexamples**

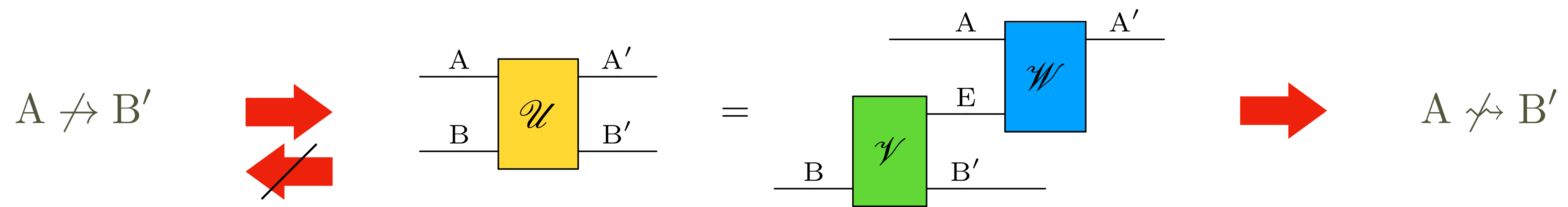
# Open question

## The quest for counterexamples



# Open question

## The quest for counterexamples



# Open question

## The quest for counterexamples

