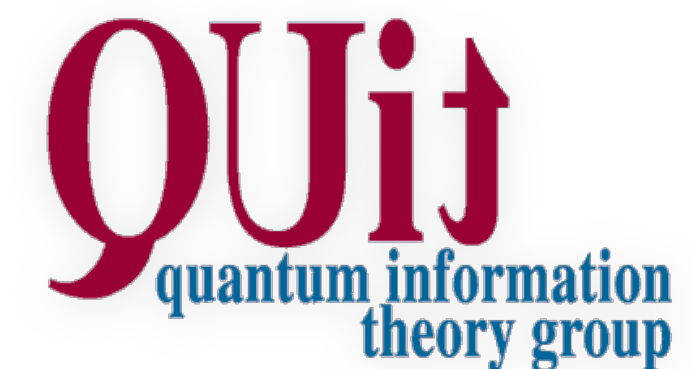


Operational probabilistic theories and cellular automata: how I learned to stop worrying and love C^* algebras

School on Advanced Topics in Quantum Information and Foundations

Quantum Information Unit and the Yukawa Institute for Theoretical Physics, Kyoto University



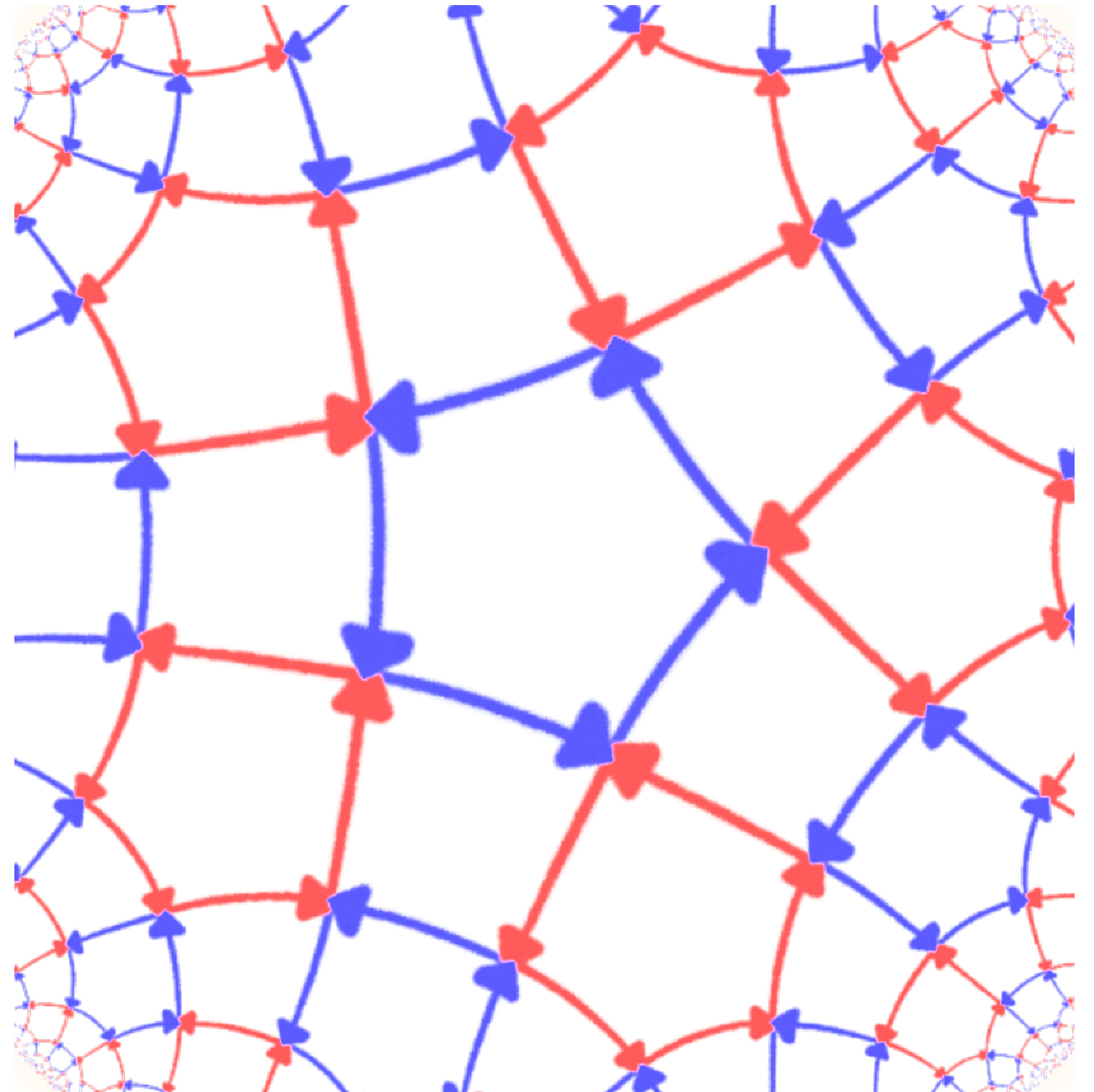
Paolo Perinotti - February 8-12 2021

Lecture 5

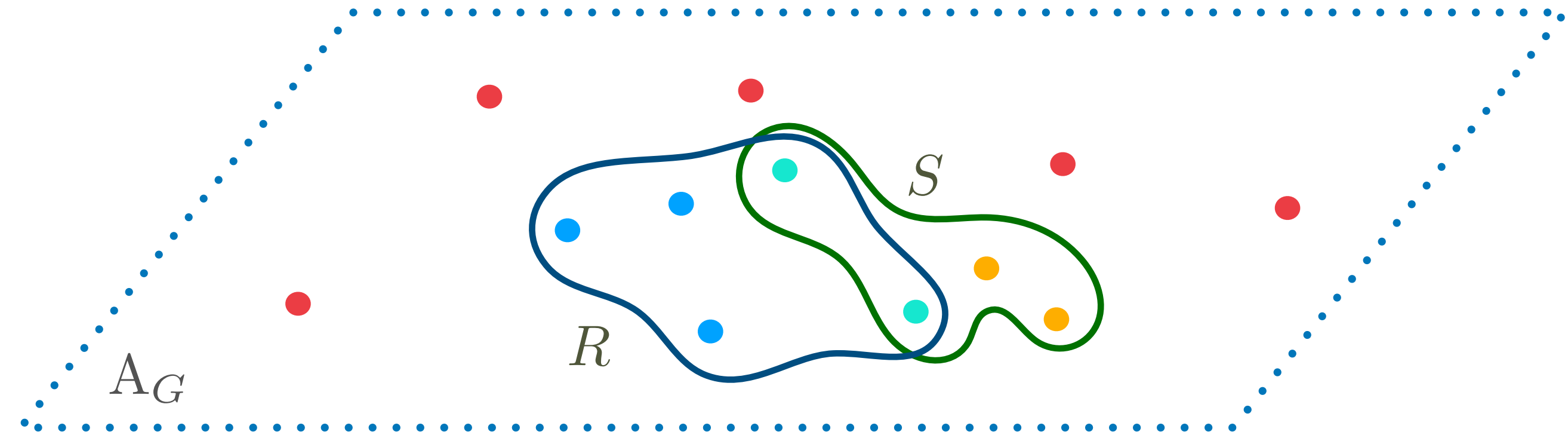
Homogeneity and Cayley graphs

Summary

- Local processes in one-to one correspondence with CA
- Homogeneity
 - Homologous regions
- Cayley graphs and translations



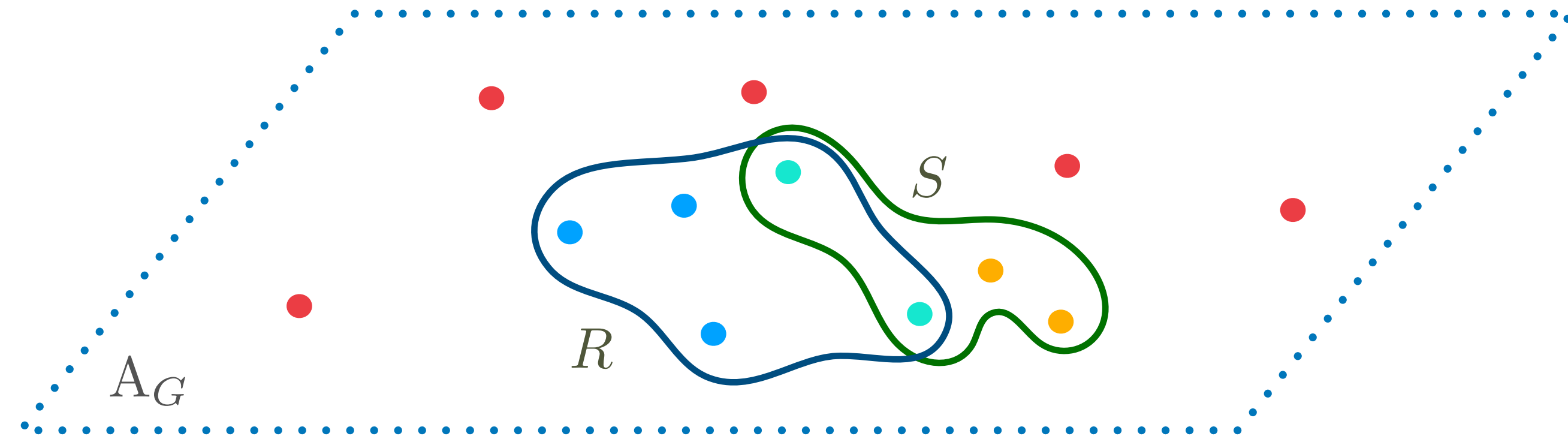
Local action



Local action

- \mathcal{V}^\dagger acts locally on R if for all C, S

$$\mathcal{V}_S^\dagger = \mathcal{V}_{R \cap S}^\dagger \otimes \mathcal{I}_{S \setminus R}$$

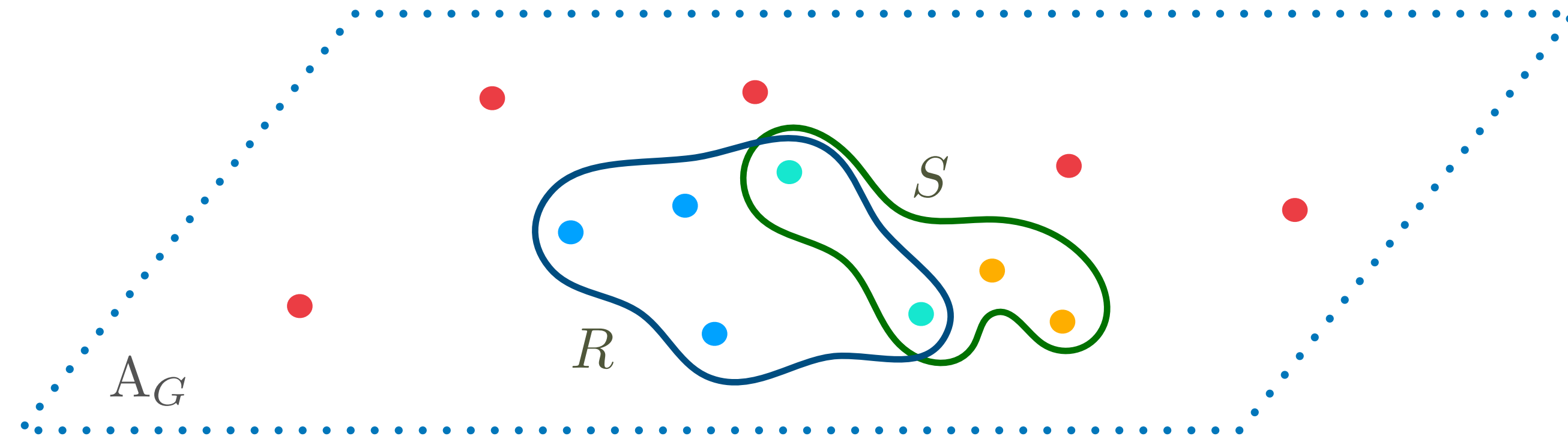


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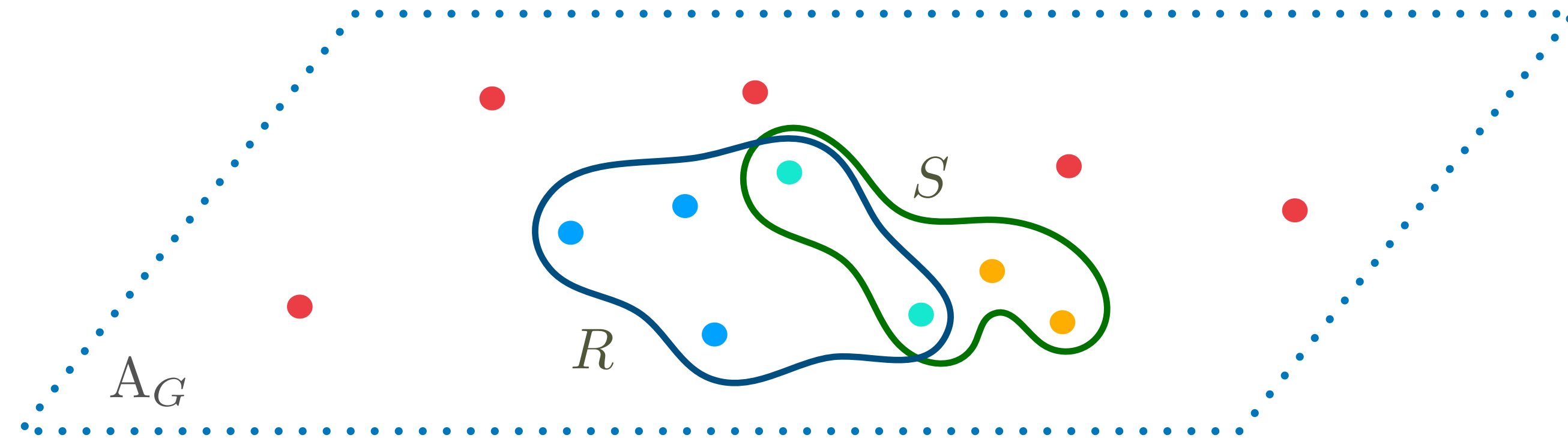
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- If \mathcal{V}^\dagger acts locally on R and \mathcal{W}^\dagger acts locally on $G \setminus R$

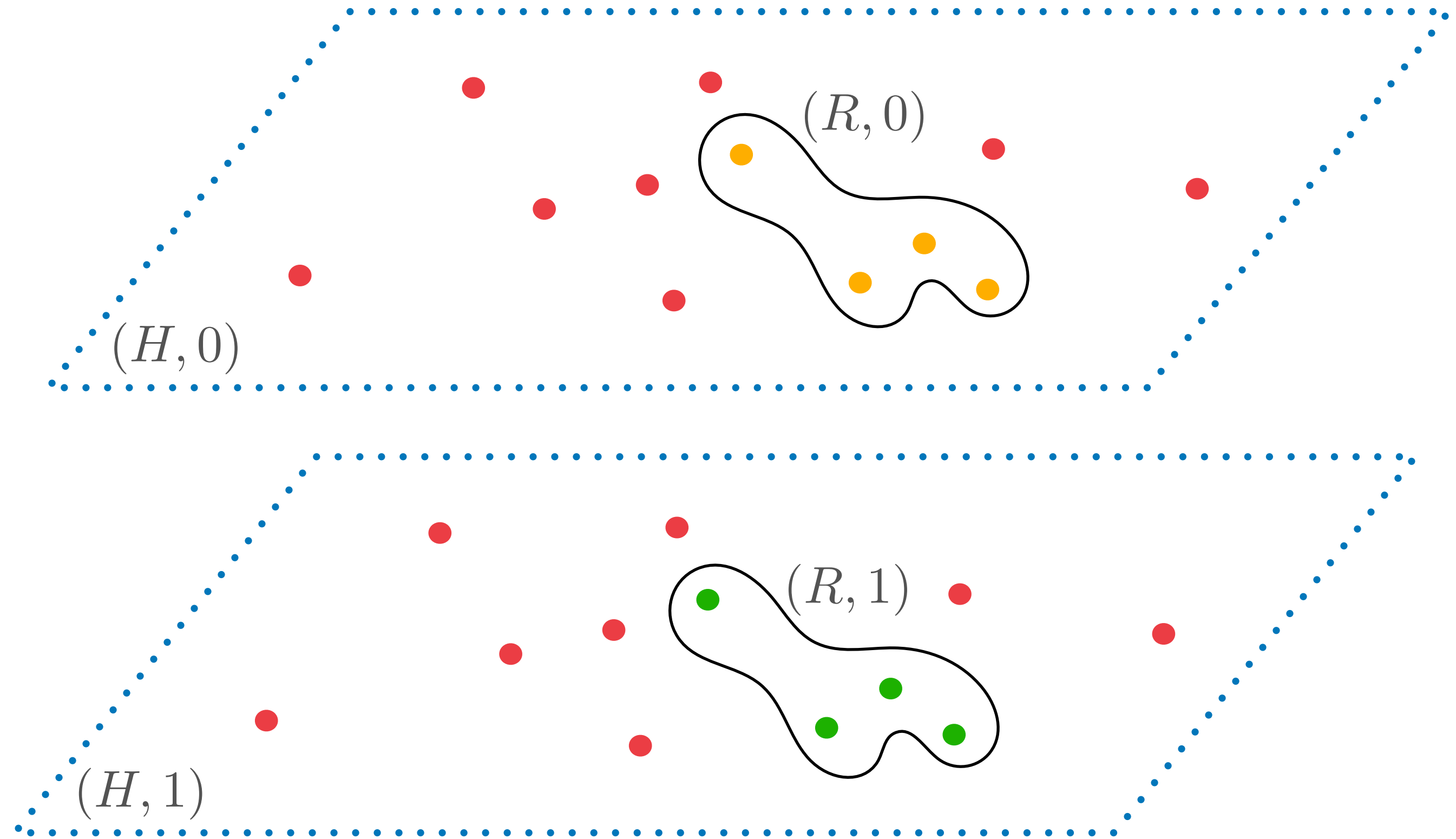
$$\mathcal{V}'^\dagger \otimes \mathcal{W}'^\dagger := \mathcal{V}^\dagger \mathcal{W}^\dagger = \mathcal{W}^\dagger \mathcal{V}^\dagger$$



Example

The swap GUR

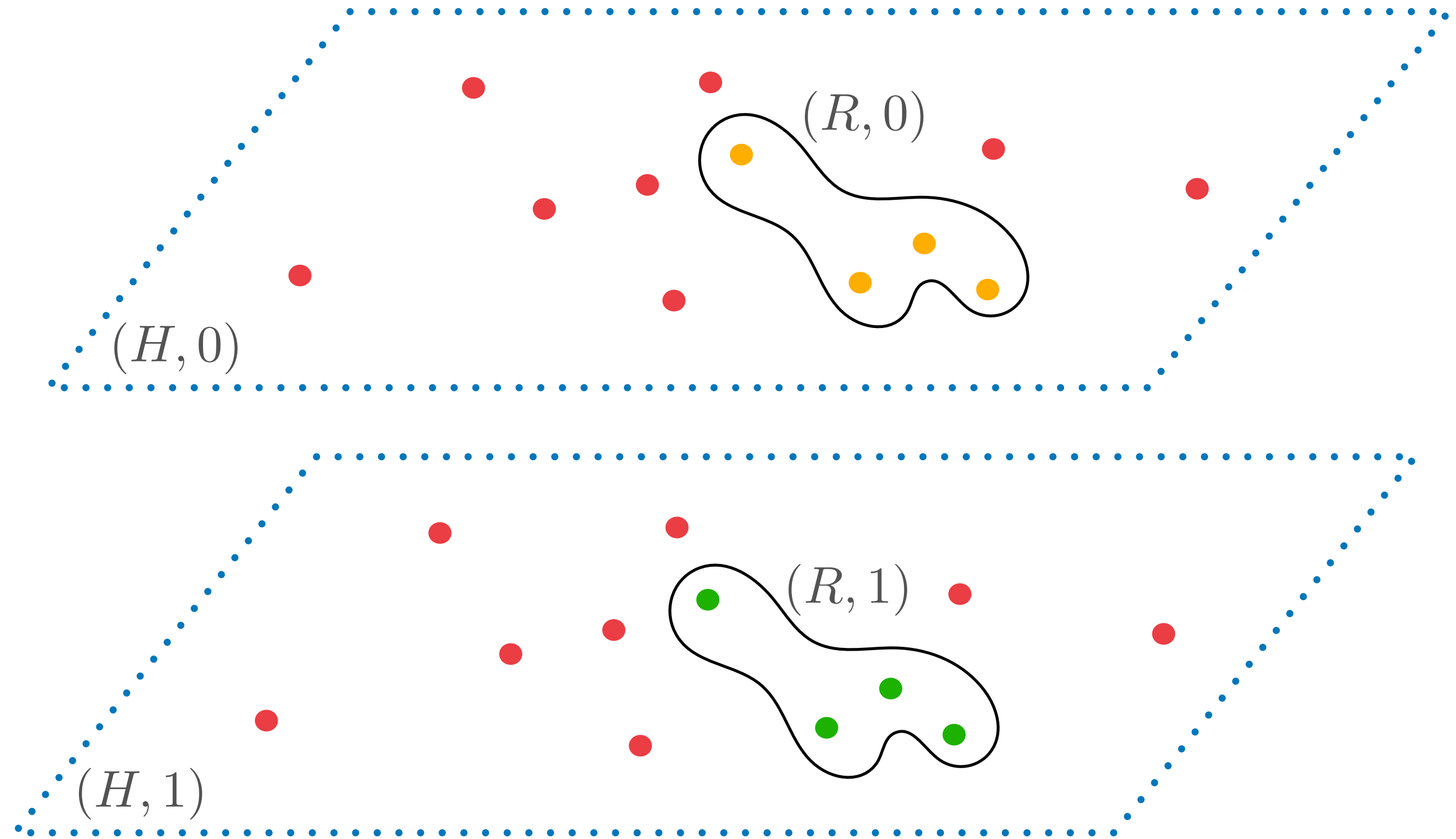
- Let $G = H \times \{0, 1\}$



Example

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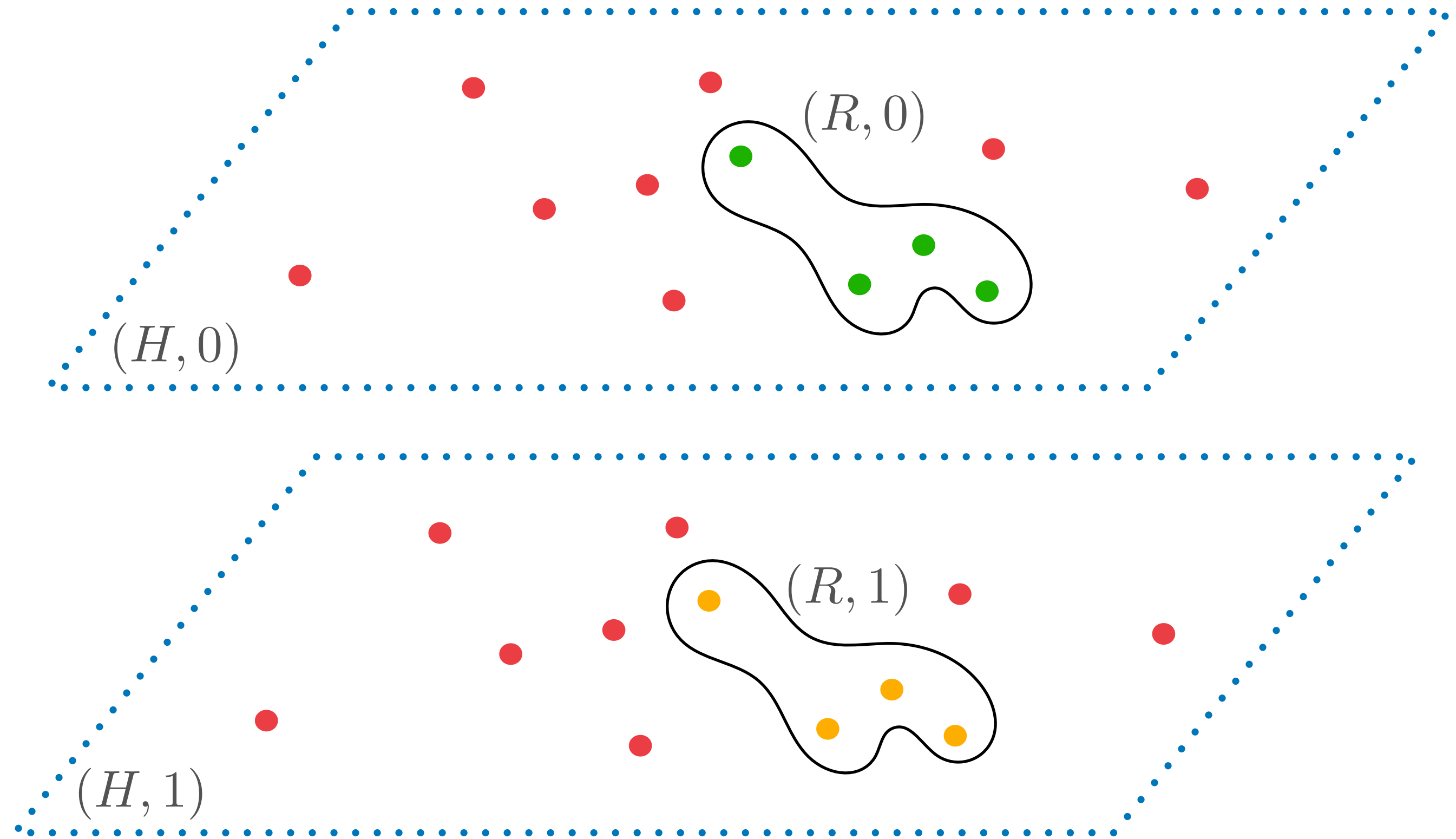
- Let $G = H \times \{0, 1\}$
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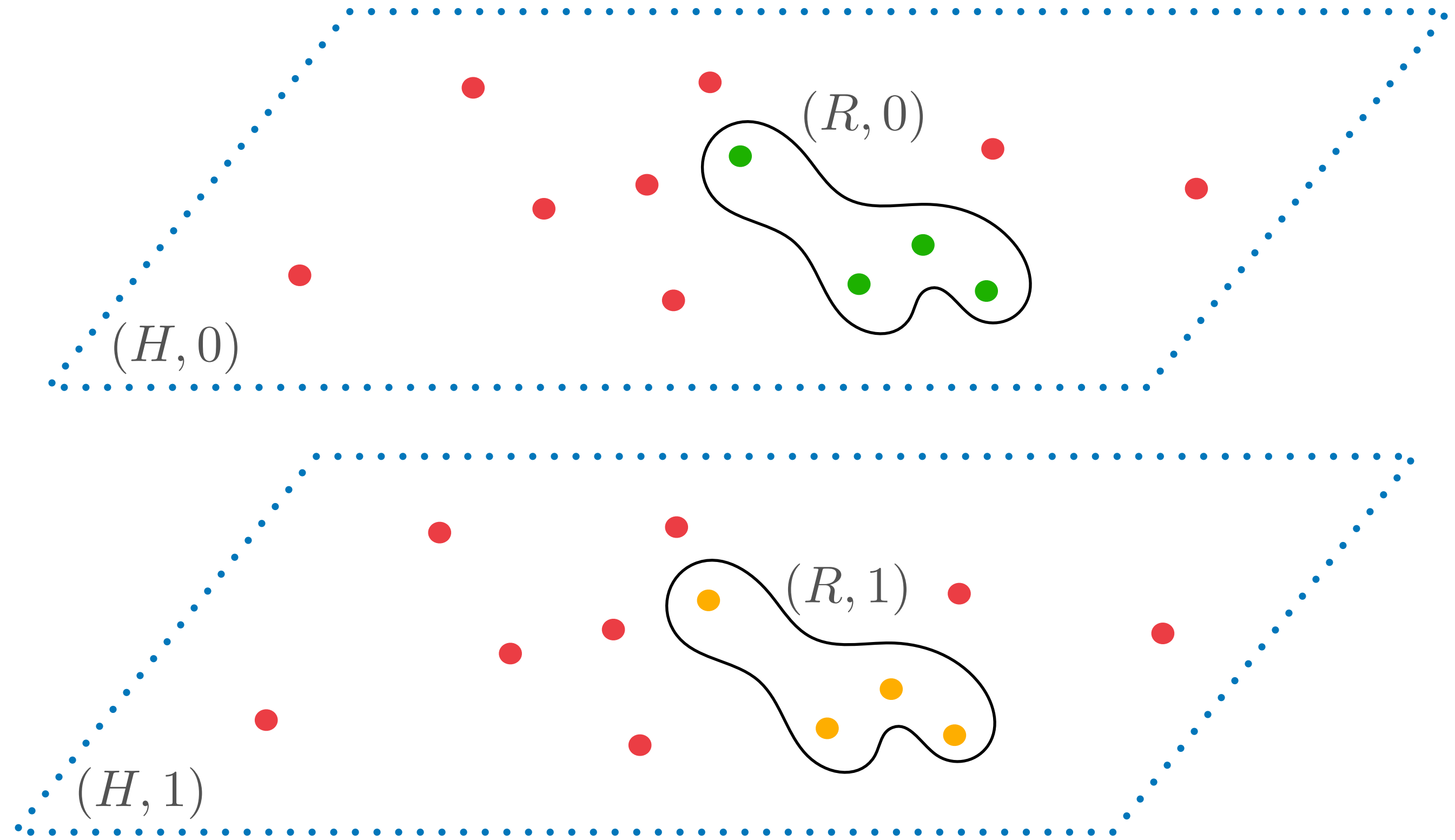
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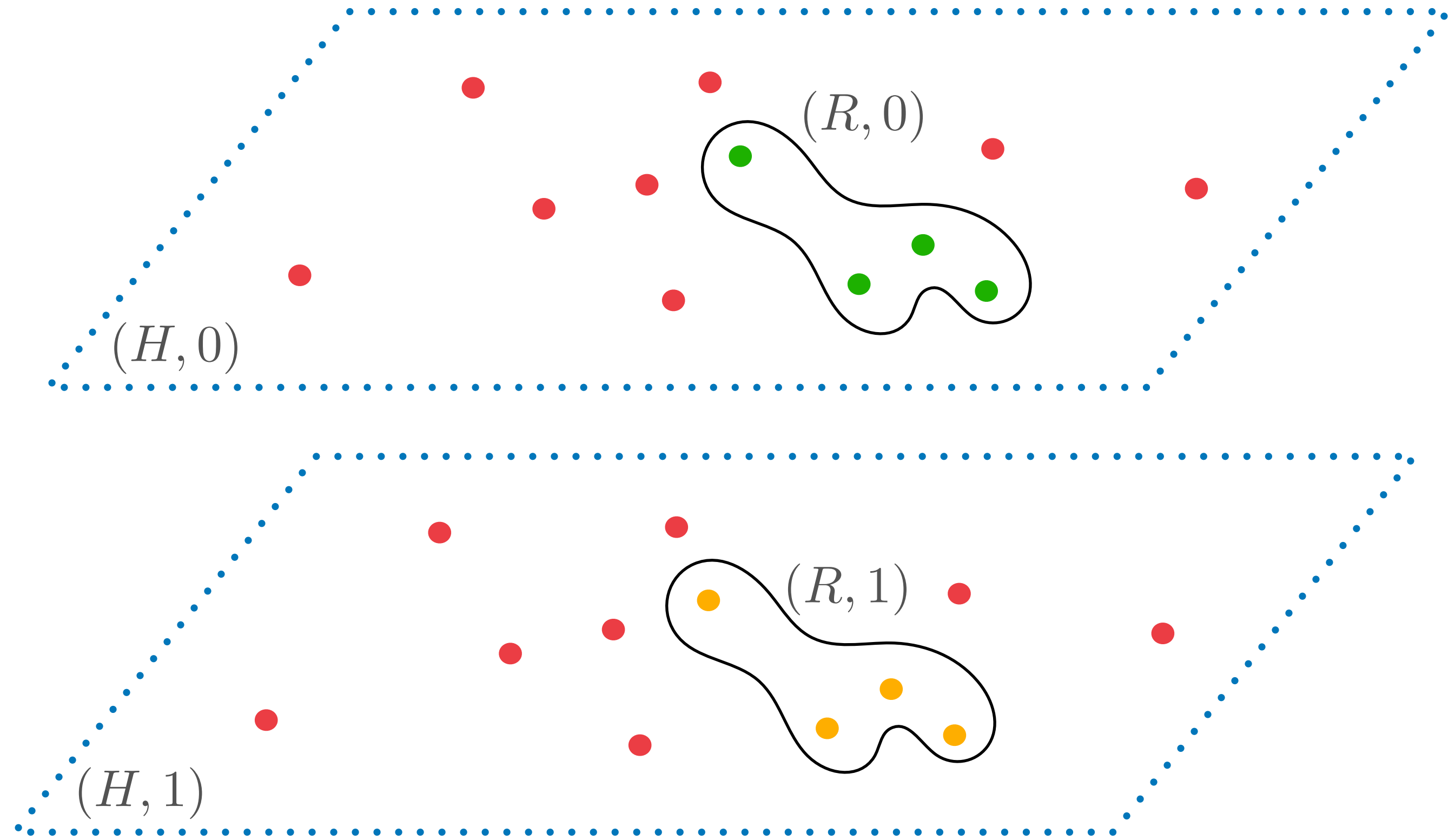
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 $\mathcal{V}^\dagger \mathcal{S}_{R \cup S}^\dagger \mathcal{V}^{-1\dagger} = \mathcal{V}^\dagger \mathcal{S}_R^\dagger \mathcal{V}^{-1\dagger} \mathcal{V}^\dagger \mathcal{S}_S^\dagger \mathcal{V}^{-1\dagger}$



The “T” process

Imprinting a GUR in local transformation

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Imprinting a GUR in local transformation

- Consider $\mathcal{V}^\dagger \otimes \mathcal{V}^{-1\dagger}$
it is easily shown that

$$\begin{aligned}\mathcal{V}^\dagger \otimes \mathcal{V}^{-1\dagger} &= (\mathcal{V}^\dagger \otimes \mathcal{I}_G^\dagger) \mathcal{S}^\dagger (\mathcal{V}^{-1\dagger} \otimes \mathcal{I}_G^\dagger) \mathcal{S}^\dagger \\ &= \left[\prod_{g \in G} (\mathcal{V}^\dagger \otimes \mathcal{I}_G^\dagger) \mathcal{S}_g^\dagger (\mathcal{V}^{-1\dagger} \otimes \mathcal{I}_G^\dagger) \right] \mathcal{S}^\dagger\end{aligned}$$

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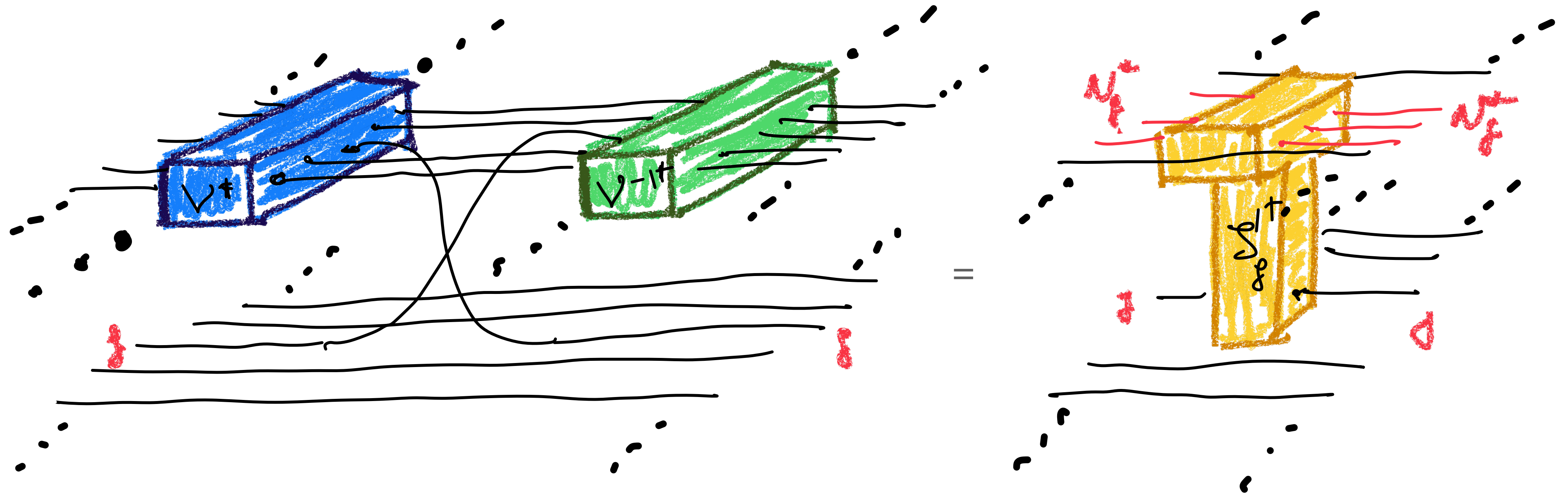
The “T” process

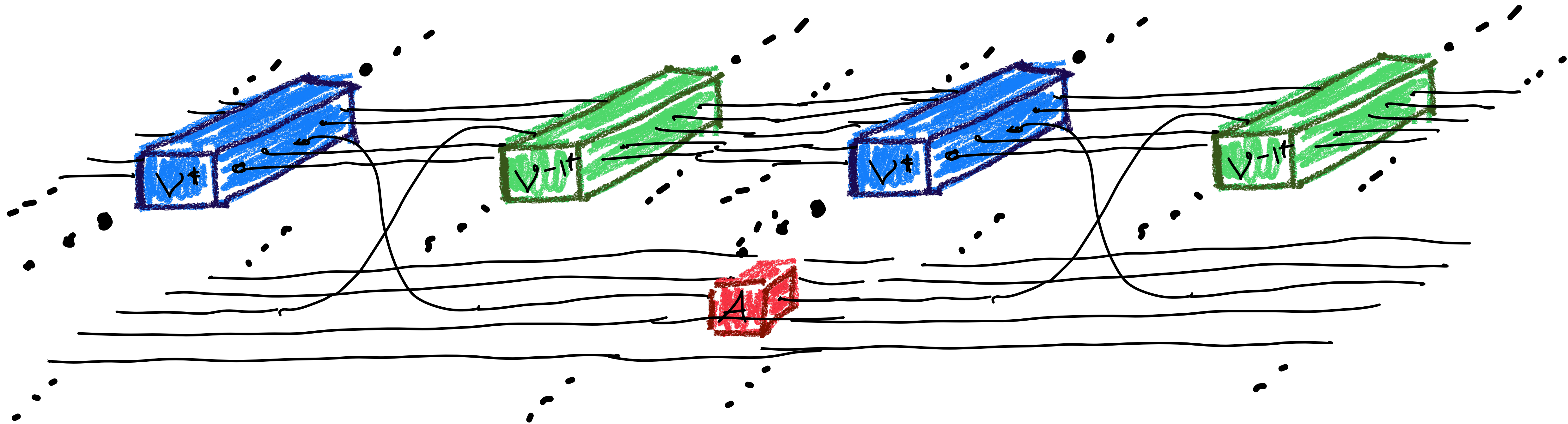
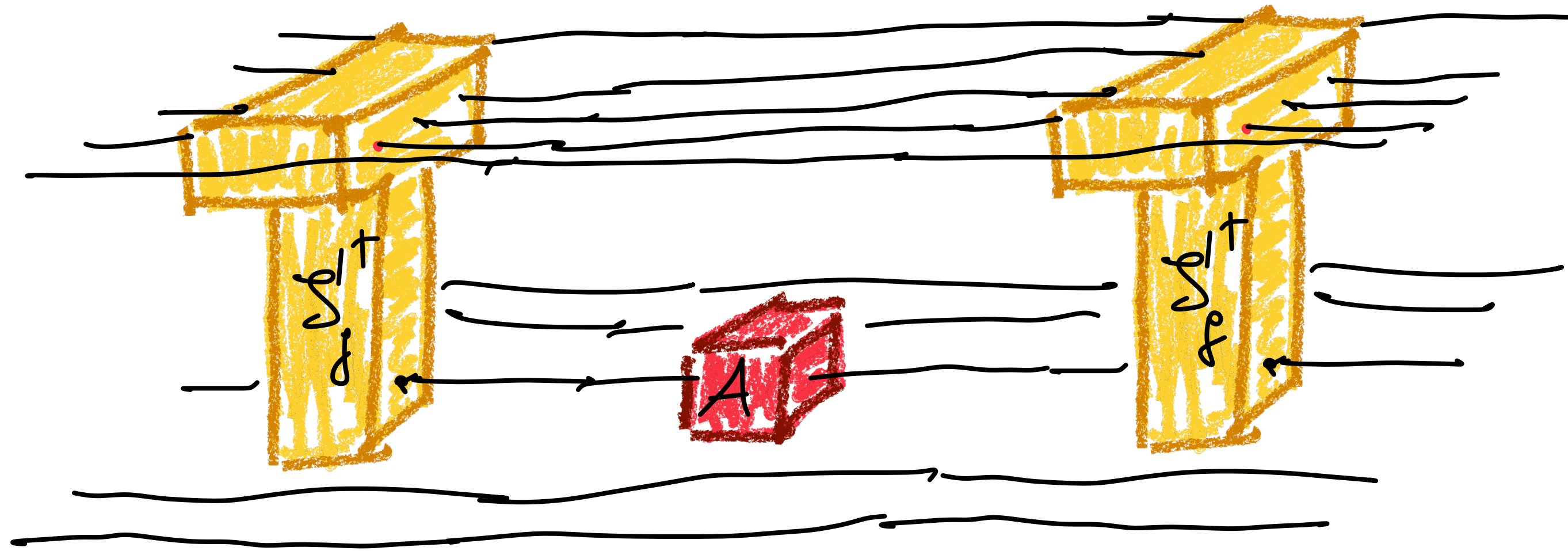
Imprinting a GUR in local transformation

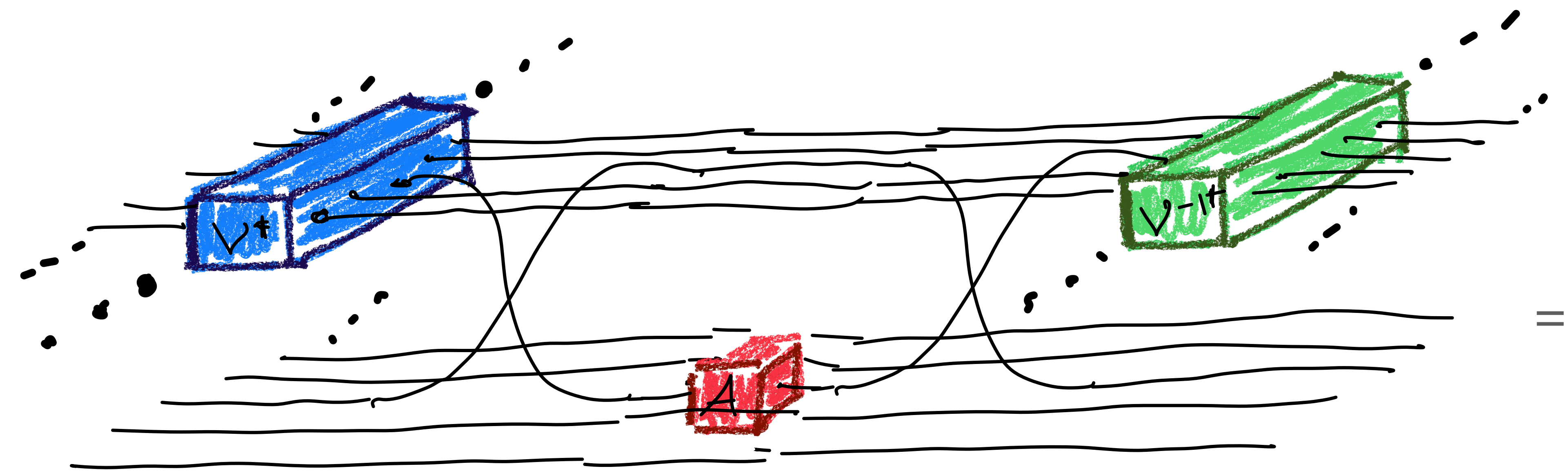
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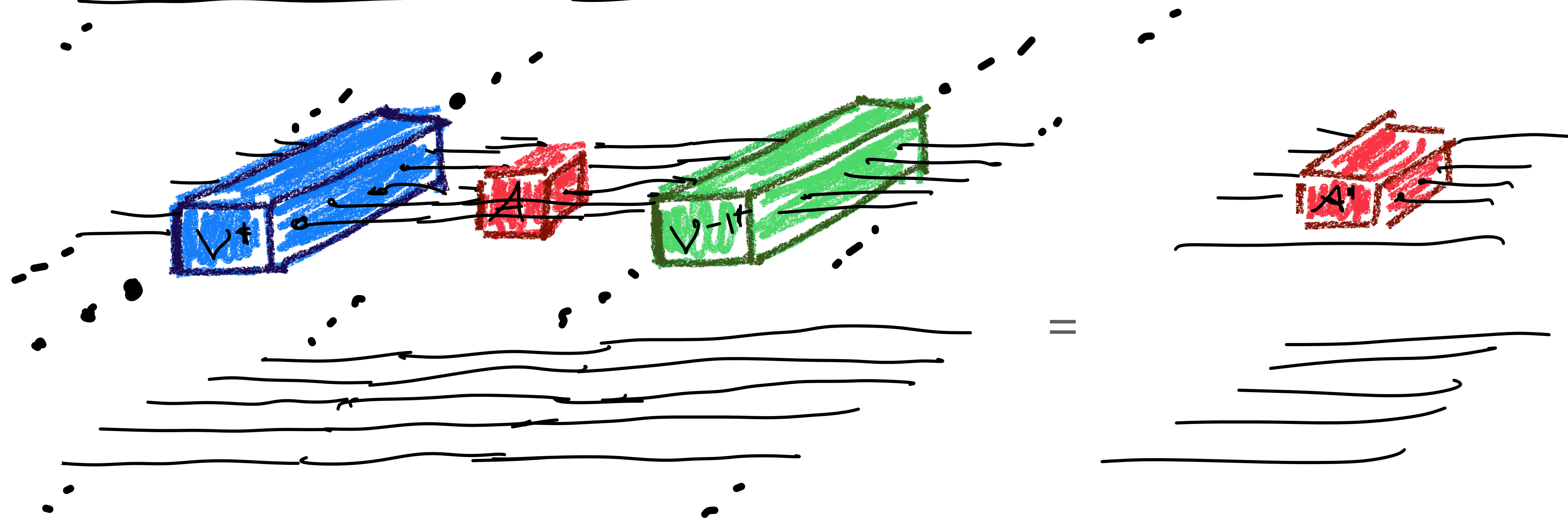
- Let $\mathcal{S}'_g^\dagger := (\mathcal{V}^\dagger \otimes \mathcal{I}_G^\dagger) \mathcal{S}_g^\dagger (\mathcal{V}^{-1\dagger} \otimes \mathcal{I}_G^\dagger)$
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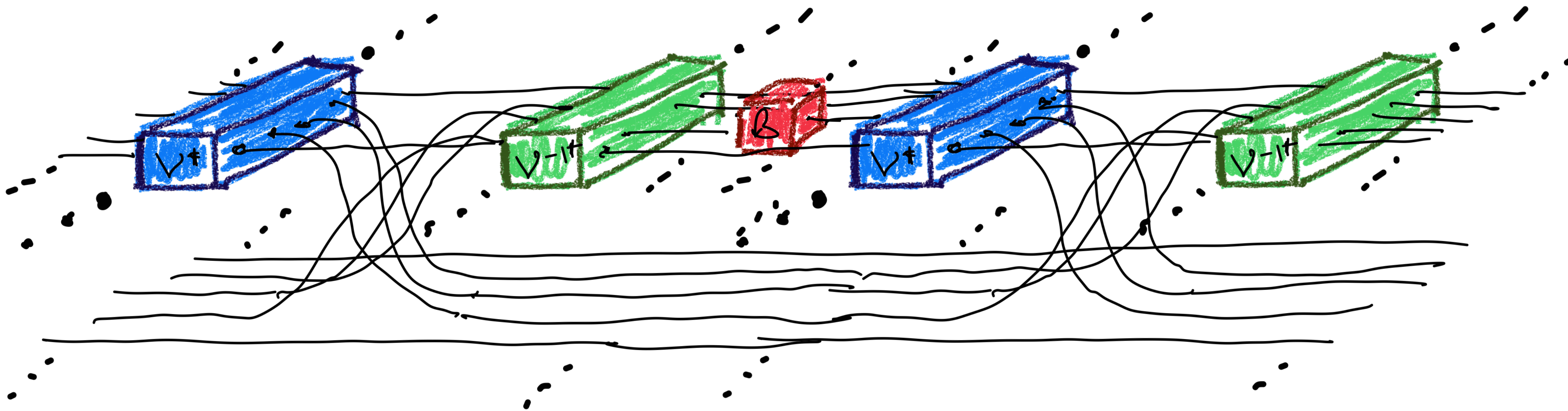
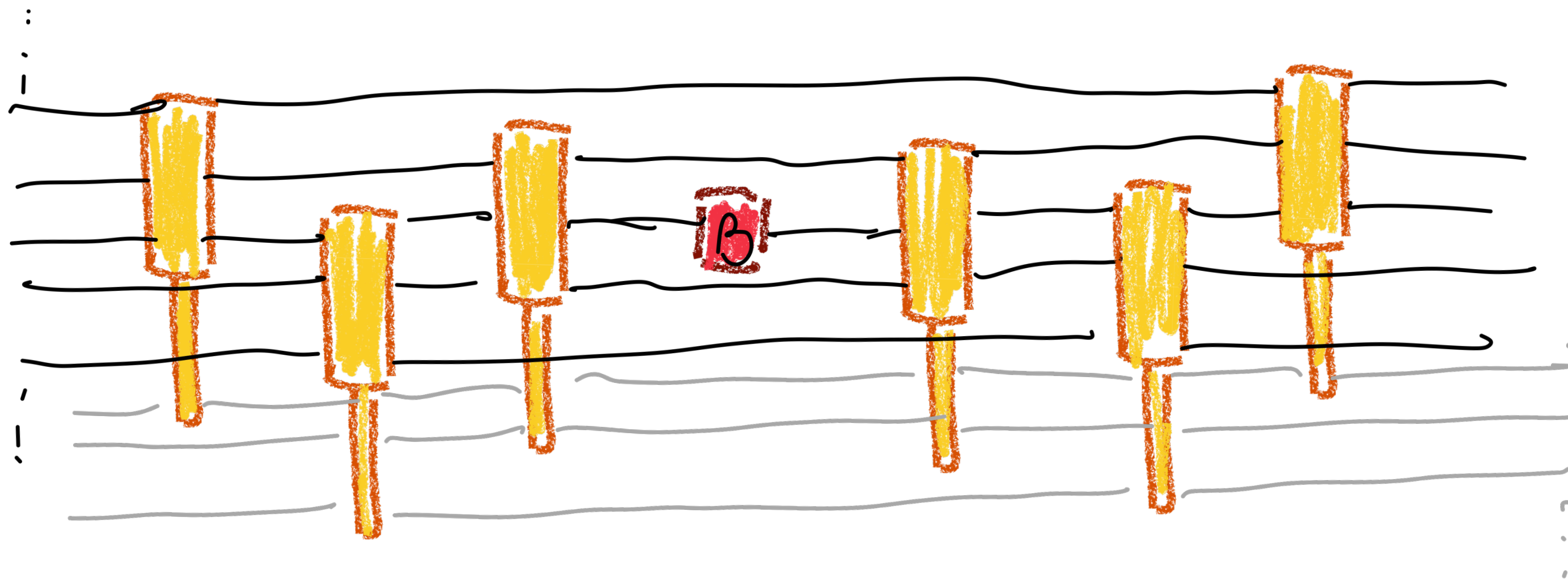


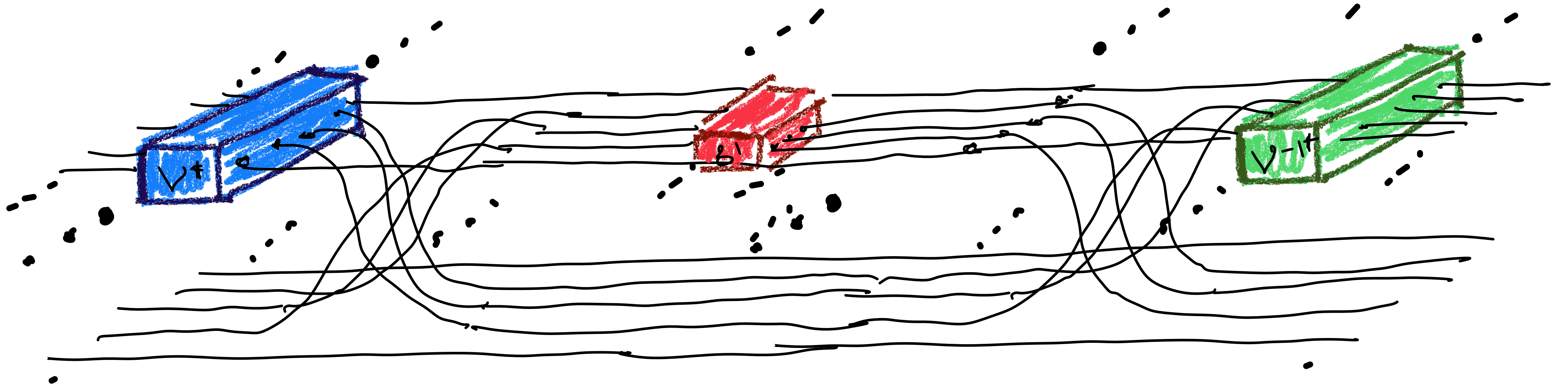


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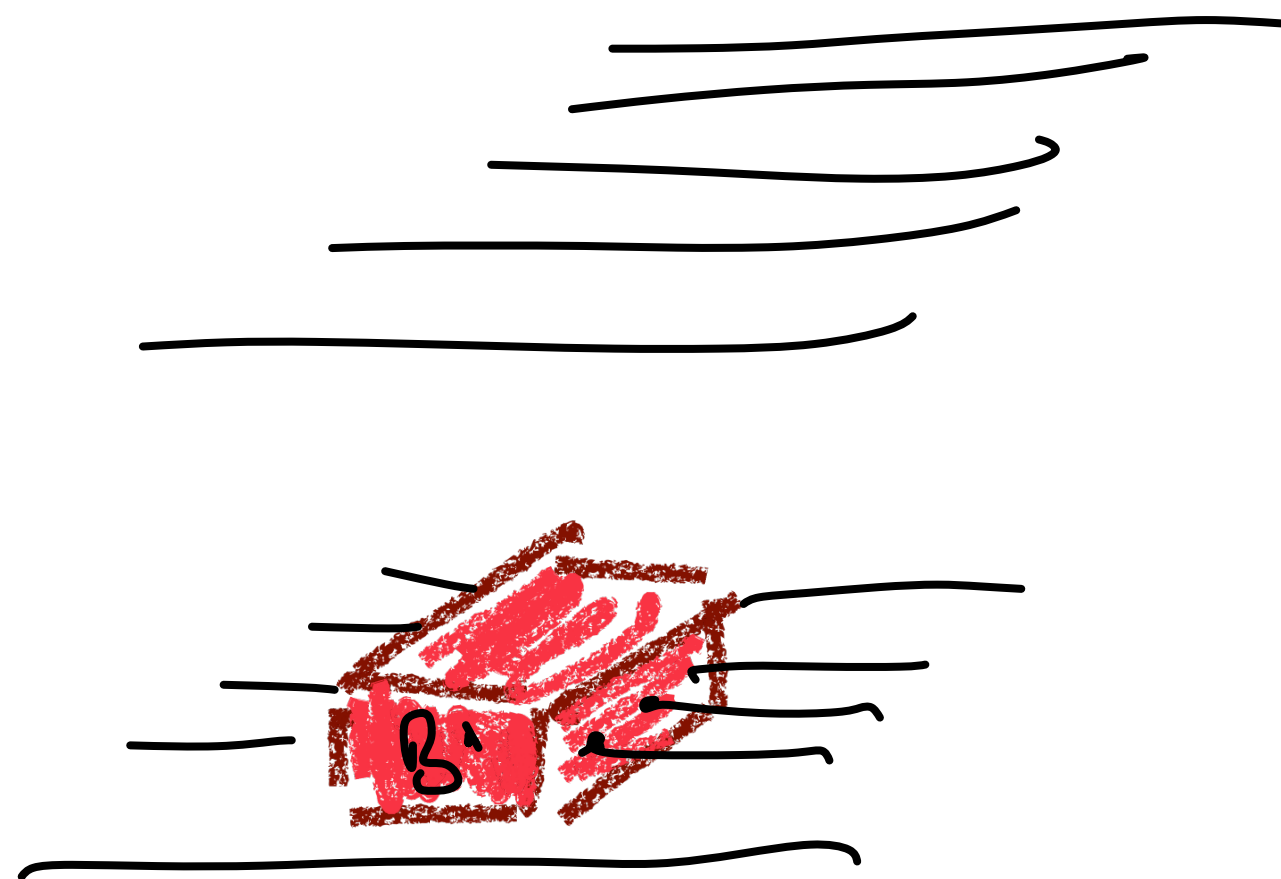


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Reversing the correspondence

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- defines a GUR through

$$\mathcal{V}' \otimes \mathcal{V}'^{-1} = \left[\prod_{g \in G} (\mathcal{V}' \otimes \mathcal{S}'_G) \mathcal{S}'_g (\mathcal{V}'^{-1} \otimes \mathcal{S}'_G) \right] \mathcal{S}'$$

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Translation invariance

- For QCA the local rule at any site summarises all the information needed about the evolution thanks to translation invariance

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- Cellular automata are defined as homogeneous global update rules
- Intuitively speaking: homogeneity consists in “treating” every cell equally

Homogeneity

Requirements for a precise definition

- The only operational criterion to establish equality of consists in “running the same test” in two cells and comparing the statistics of outcomes

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Homogeneity

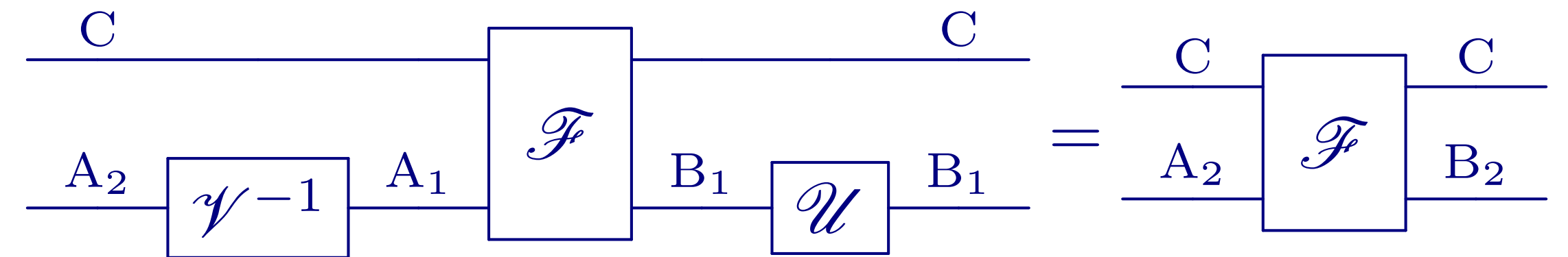
Requirements for a precise definition

- The only operational criterion to establish equality of consists in “running the same test” in two cells and comparing the statistics of outcomes
- The definition of homogeneity requires first a precise notion of
 - “running the same test” on different cells
 - exchanging the role of two cells

Operationally equivalent regions

Running the same test

- Two systems A and B are operationally equivalent ($A \cong B$) if there is a reversible transformation between them

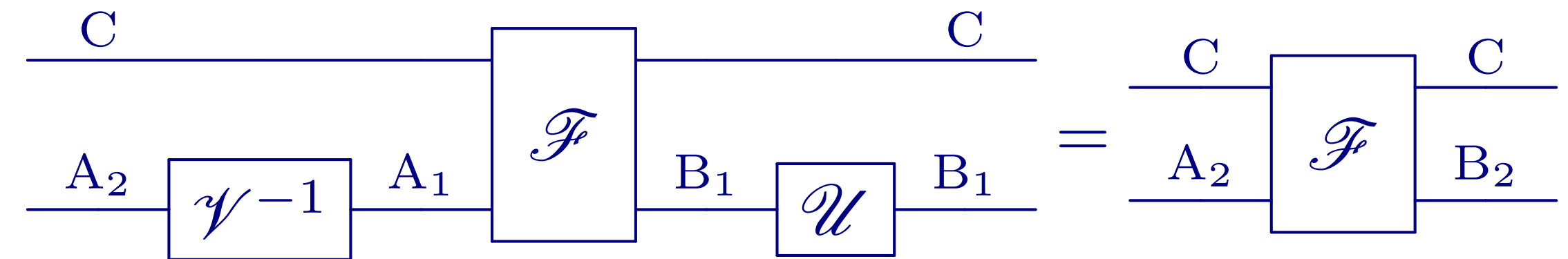


$$\psi \in [A_1 \rightarrow A_2], \quad \mathcal{U} \in [B_1 \rightarrow B_2]$$

Operationally equivalent regions

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- The reversible transformation defines the notion of “performing the same test”

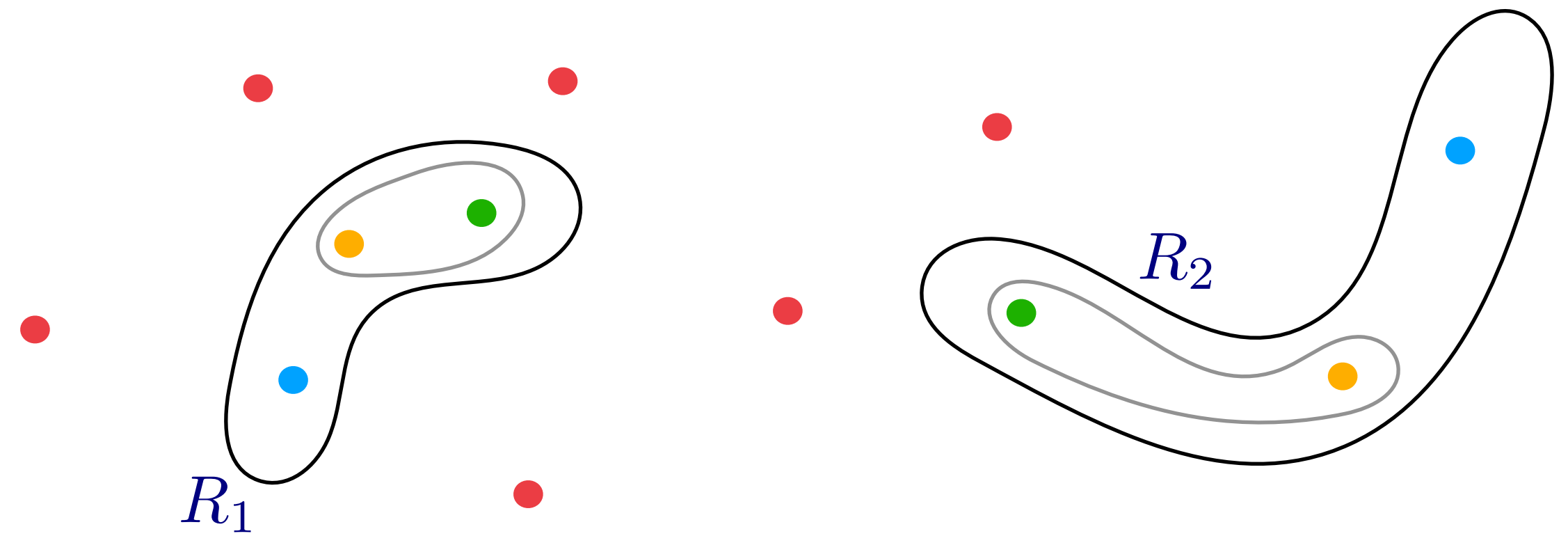


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Operationally equivalent regions

Accounting for internal structure

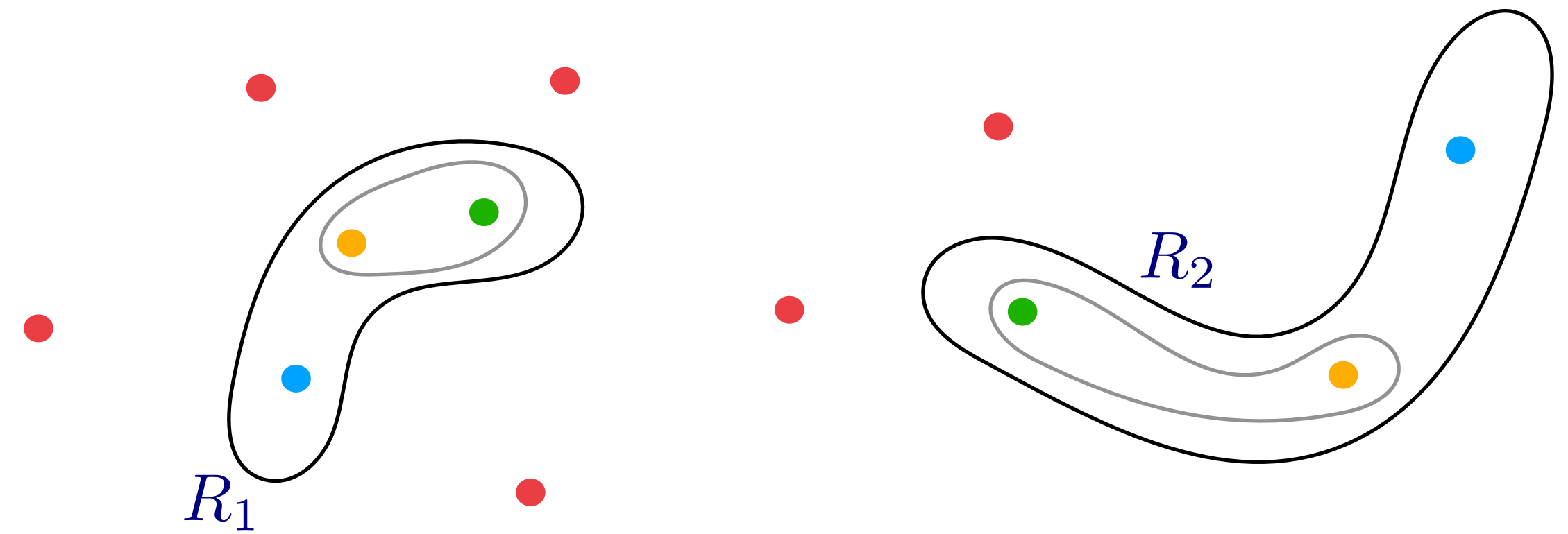
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Operationally equivalent regions

Accounting for internal structure

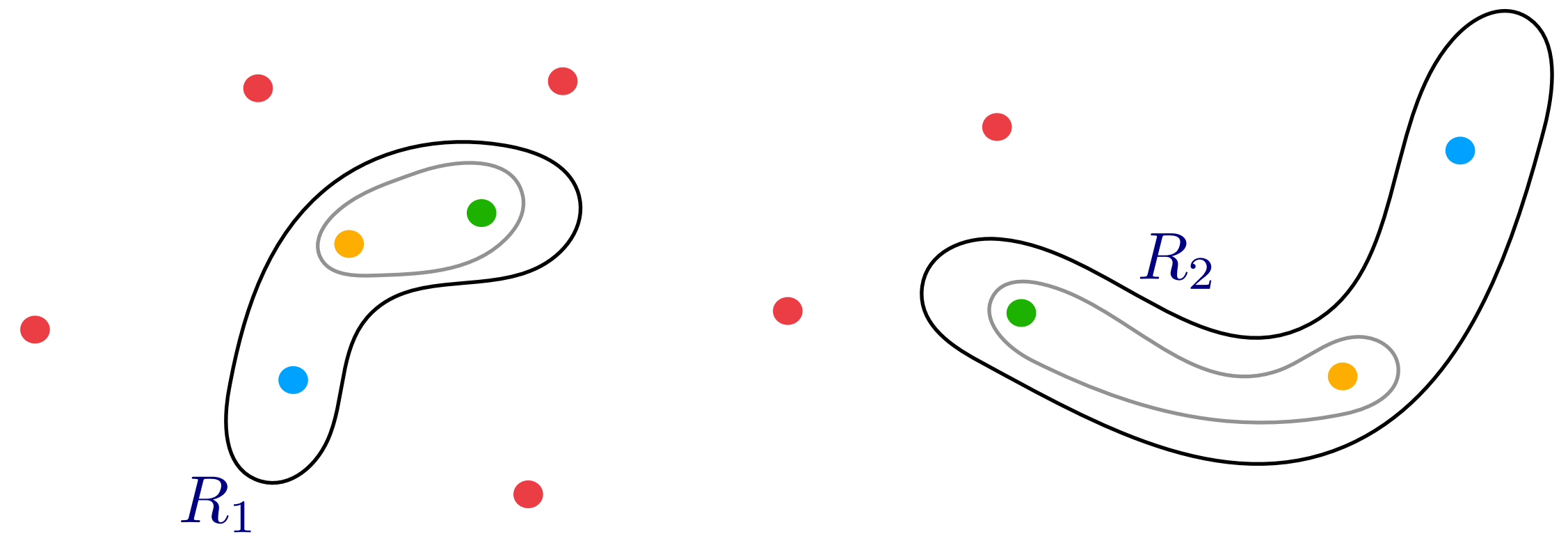
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Operationally equivalent regions

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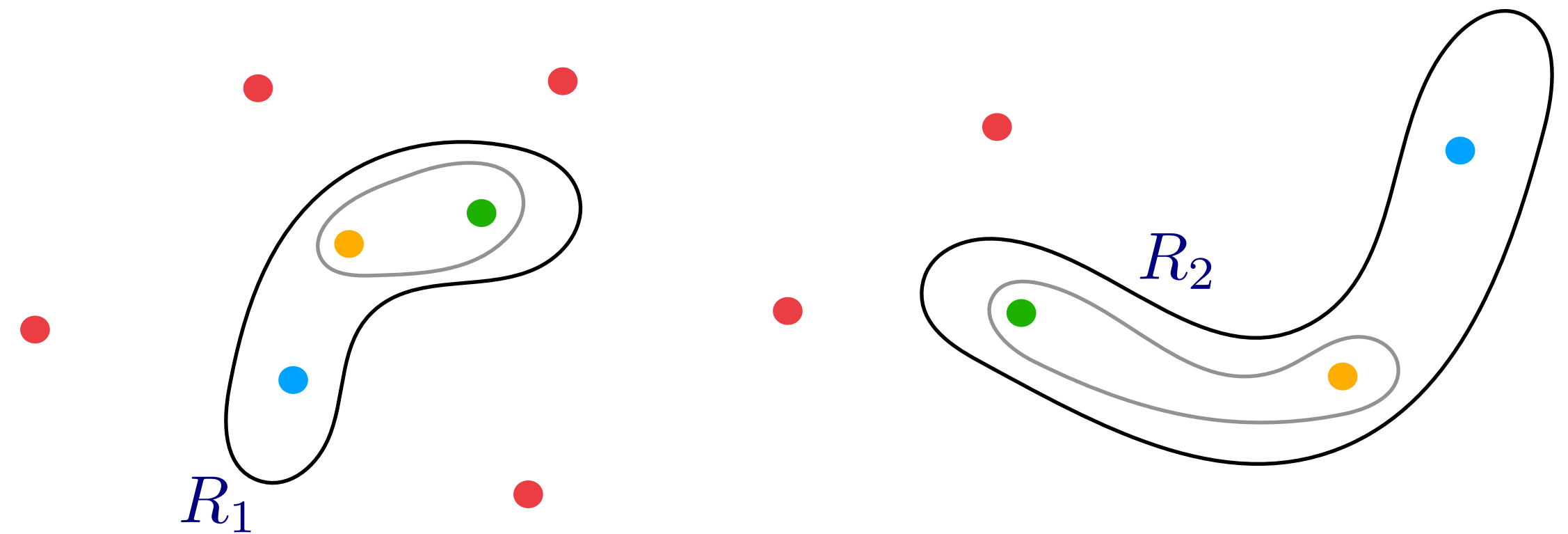
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Operationally equivalent regions

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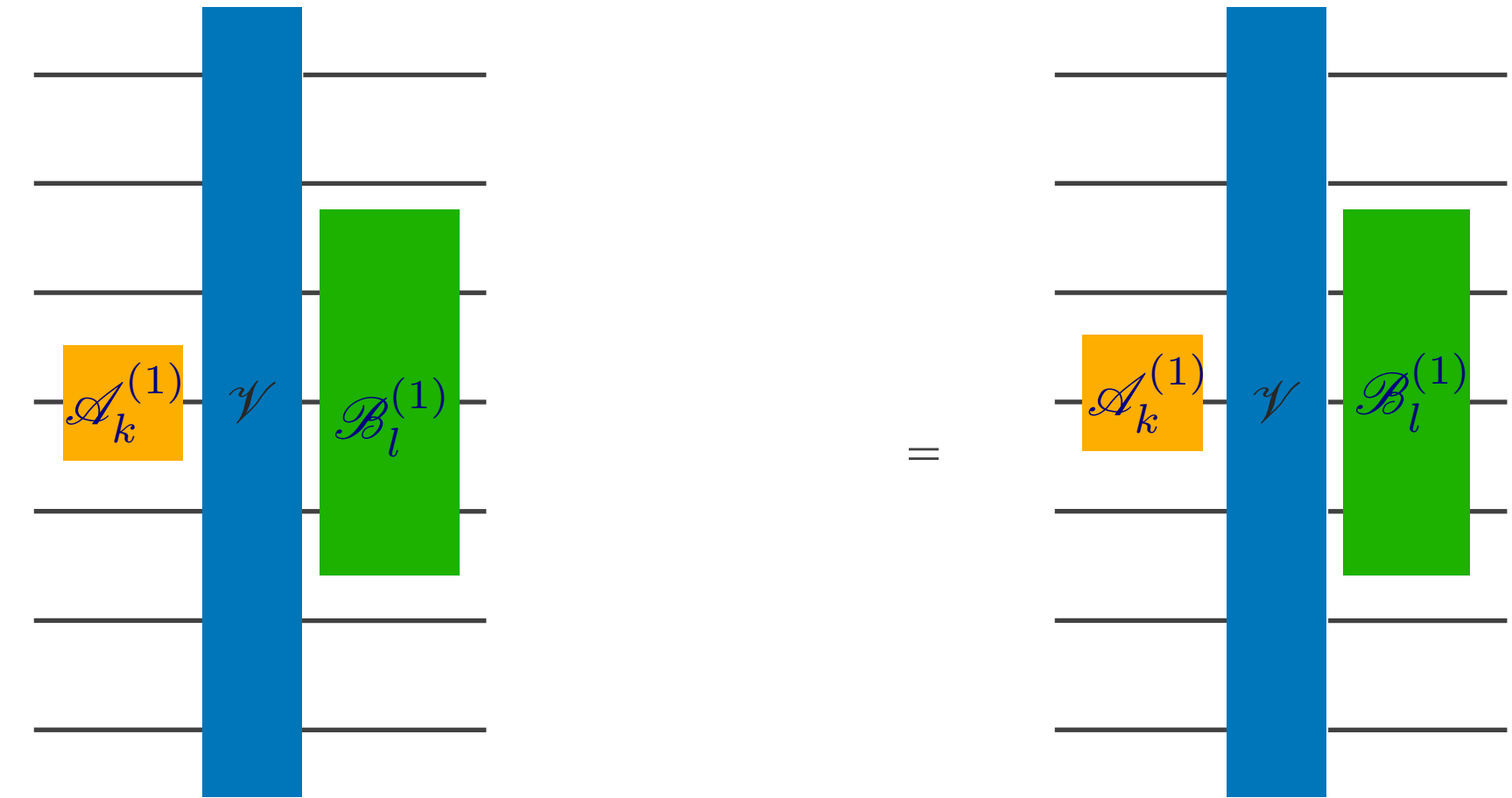
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 $R_1 = \{g_1, \dots, g_l\}$, $R_2 = \{h_1, \dots, h_l\}$, $A_{g_i} \cong A_{h_i}$
- In this case $\mathcal{U}_{R_1, R_2} = \bigotimes_{g_i \in R_1} \mathcal{U}_{g_i, h_i}$



Homologous regions

“Treated in the same way”

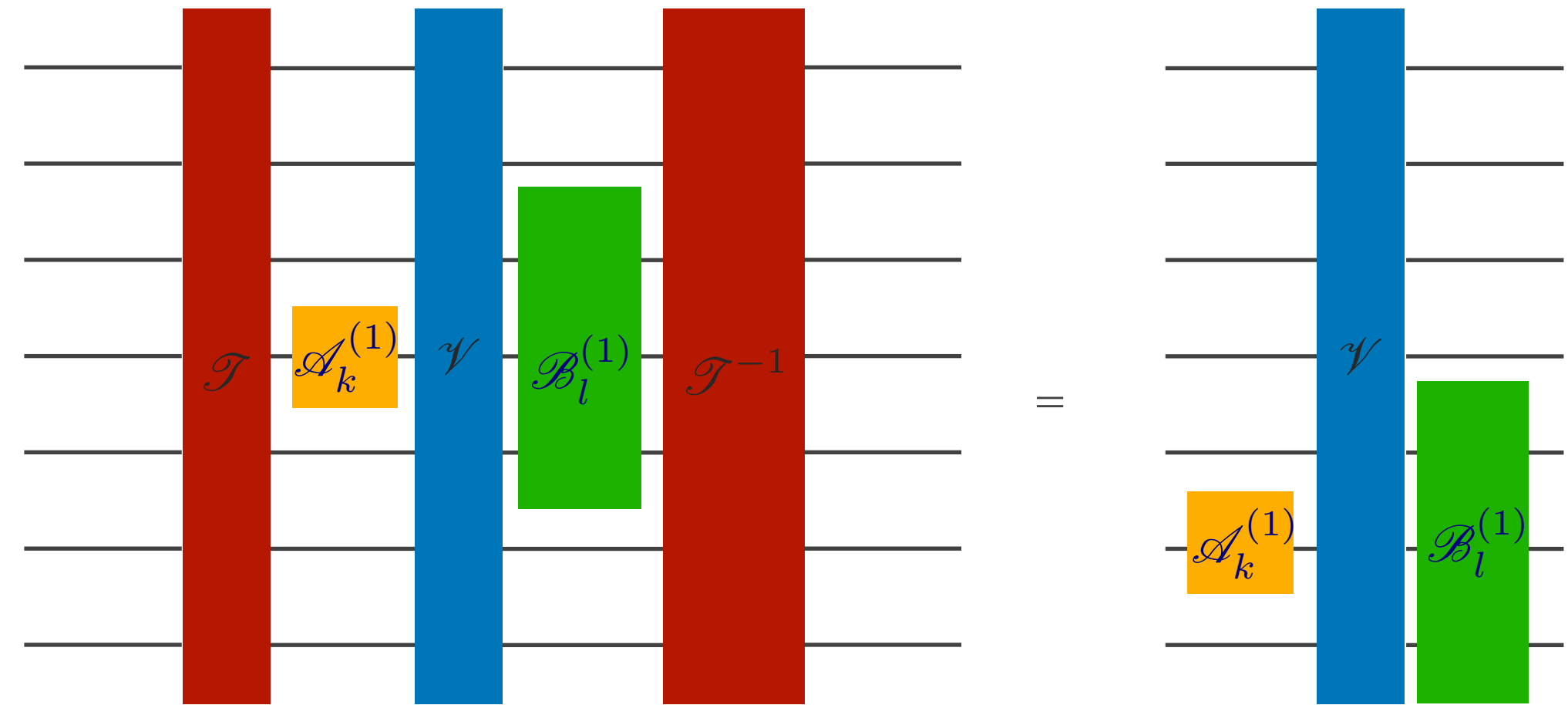
- Given a GUR $(G, A, \mathcal{V}^\dagger)$, the region R_1 is **homologous** to an o.e. region R_2 if $N_{R_1}^+$ is o.e. to $N_{R_2}^+$, and there exists a GUR $(G, A, \mathcal{T}^\dagger)$ such that $\mathcal{B}_l^{(1)} \mathcal{V} \mathcal{A}_k^{(1)} = \mathcal{T}^{-1} \mathcal{B}_l^{(2)} \mathcal{V} \mathcal{A}_k^{(2)} \mathcal{T}$. We write $R_1 \bowtie_{\mathcal{T}} R_2$



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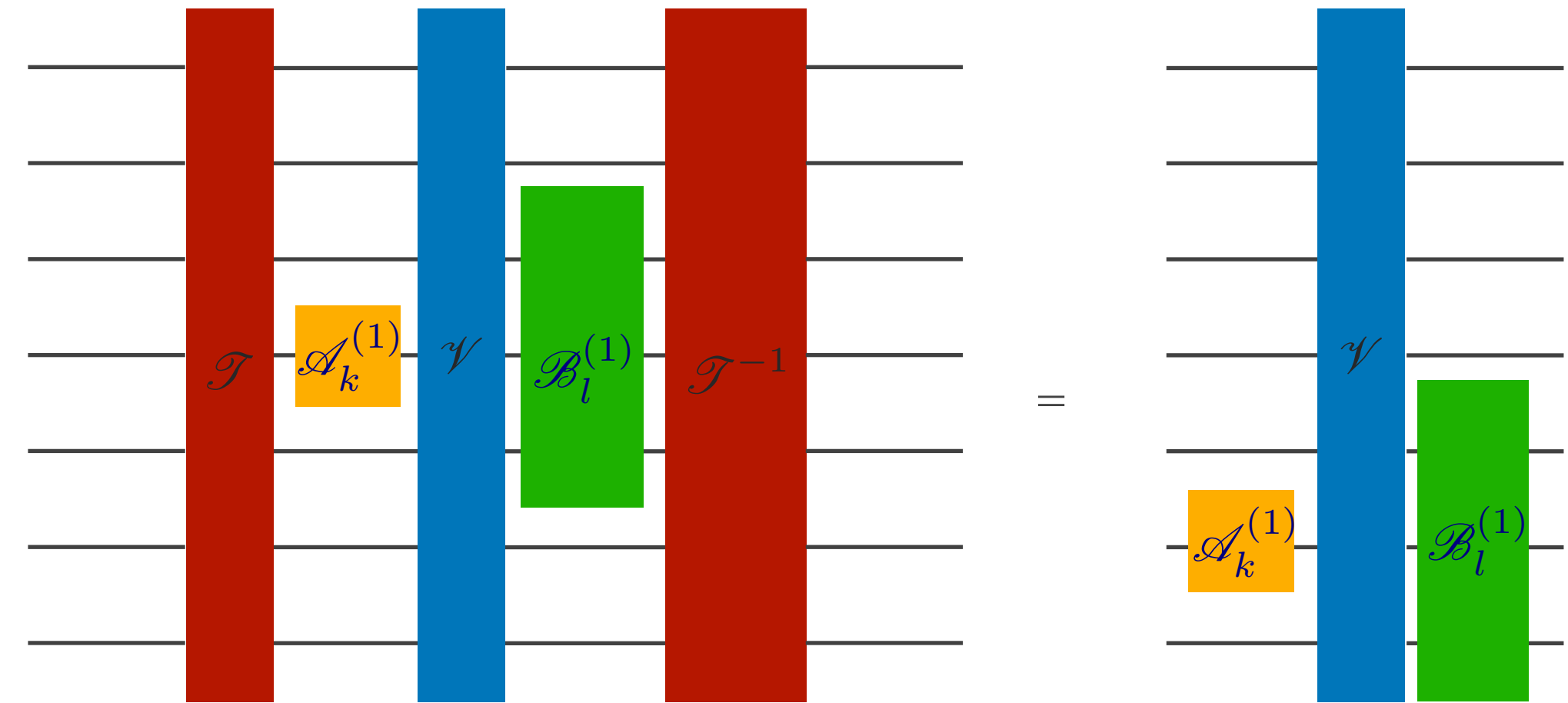
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- If $R_1 \bowtie_{\mathcal{T}} R_2$, one has $\mathcal{T}^{-1} \psi \mathcal{T} = \psi$



Absolute discrimination

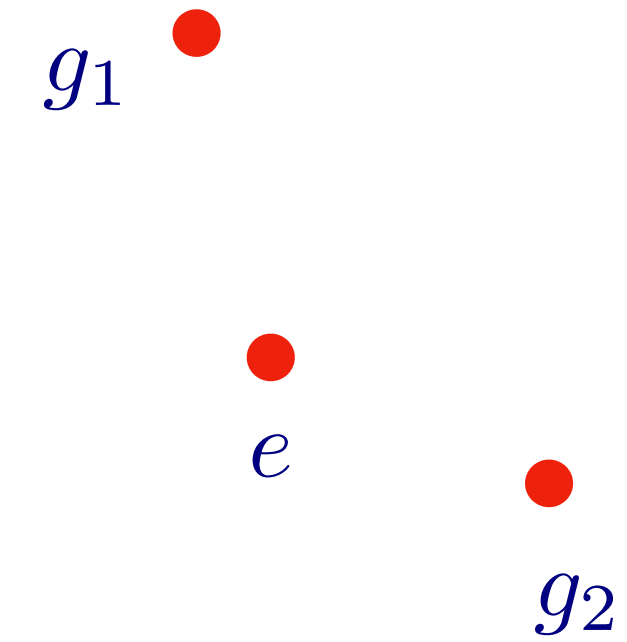
“Treated differently”

- Given a GUR $(G, A, \mathcal{V}^\dagger)$, the cells g_1 and g_2 are discriminated by \mathcal{V}^\dagger if for every GUR $(G, A, \mathcal{T}^\dagger)$ there exists a region $R_1 \ni g_1$ that is not homologous to any $R_2 \ni g_2$ through \mathcal{T}

Relative discrimination

“Seen differently from e ”

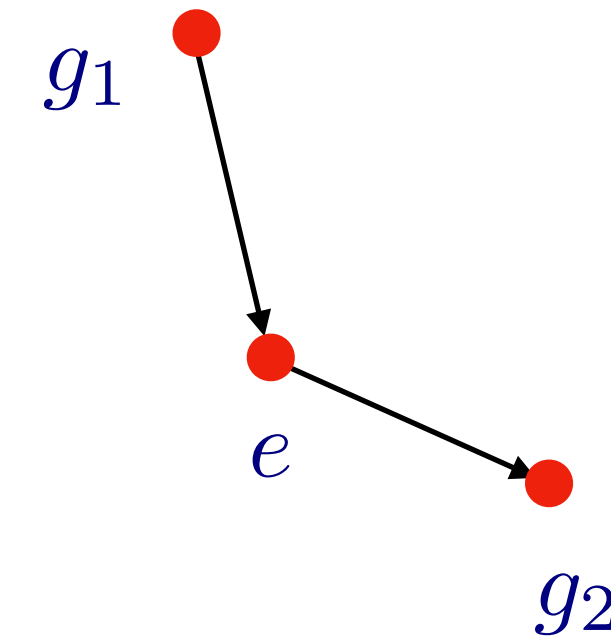
- Given a GUR $(G, A, \mathcal{V}^\dagger)$, the cells g_1 and g_2 are discriminated by \mathcal{V}^\dagger with respect to $e \in G$ if for every GUR $(G, A, \mathcal{T}^\dagger)$ there exists a region $R_1 \ni \{g_1, e\}$ that is not homologous to any $R_2 \ni \{g_2, e\}$, through \mathcal{T} , with $e \bowtie_{\mathcal{T}} e$ and $g_1 \bowtie_{\mathcal{T}} g_2$



Relative discrimination

“Seen differently from e ”

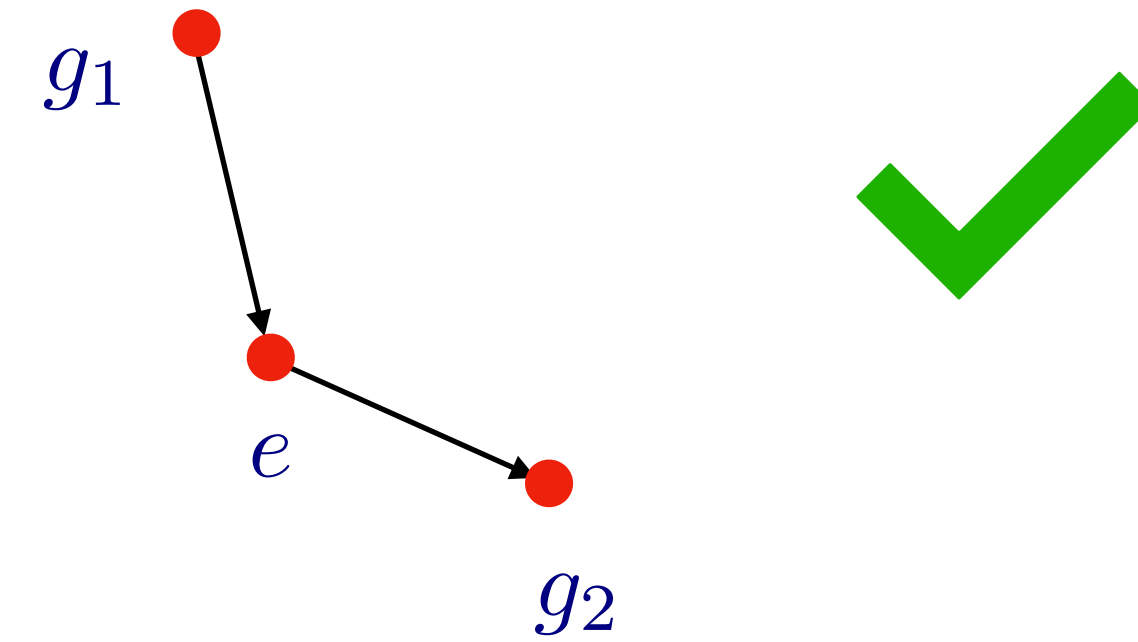
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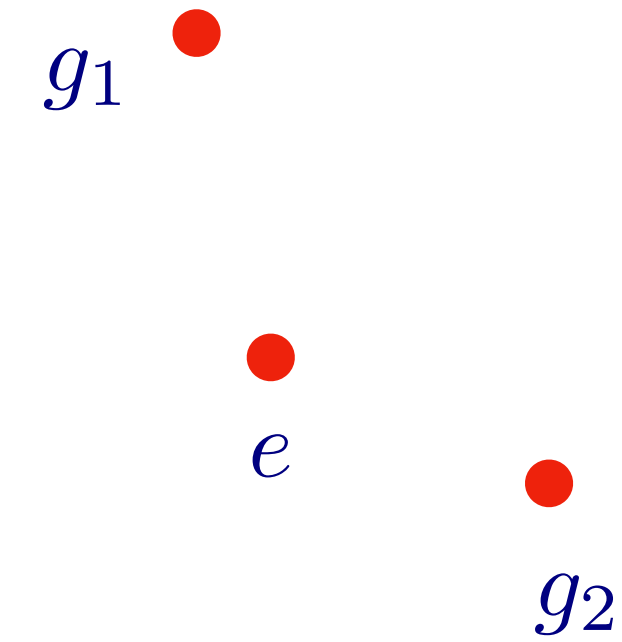
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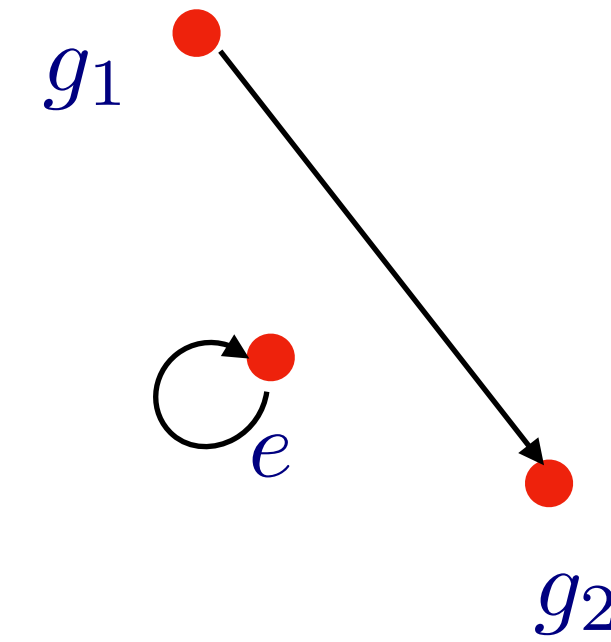
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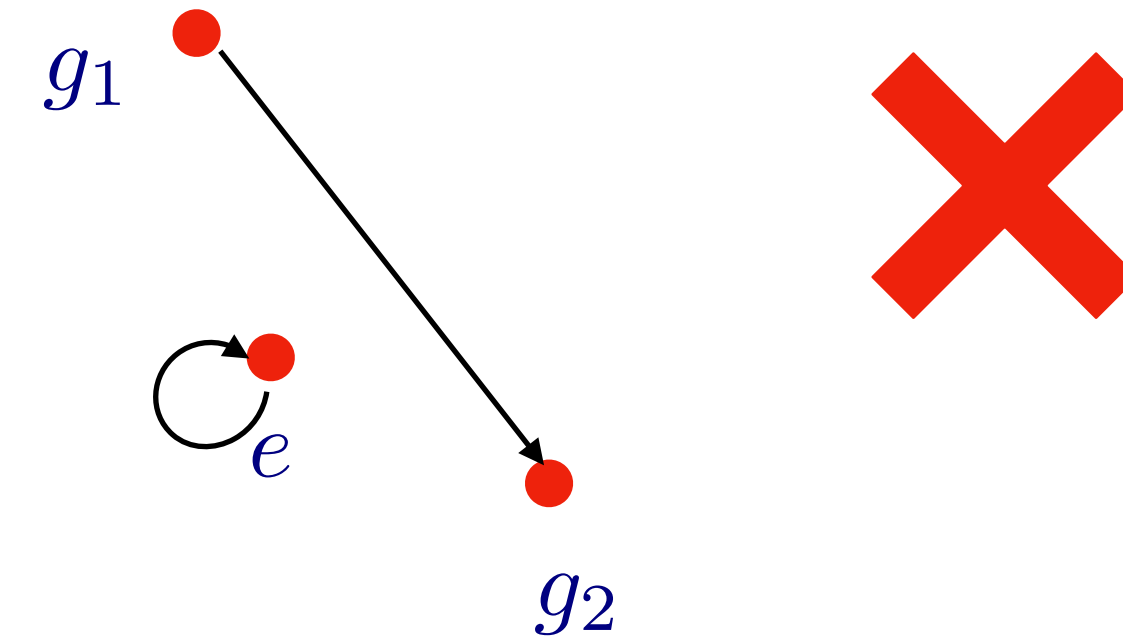
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Homogeneity

No two cells are discriminated

- Given a GUR $(G, A, \mathcal{V}^\dagger)$, we say that it is homogeneous if for every two cells $g_1, g_2 \in \mathcal{V}^\dagger$ does not discriminate them, but it does relatively to a third cell e

Homogeneity

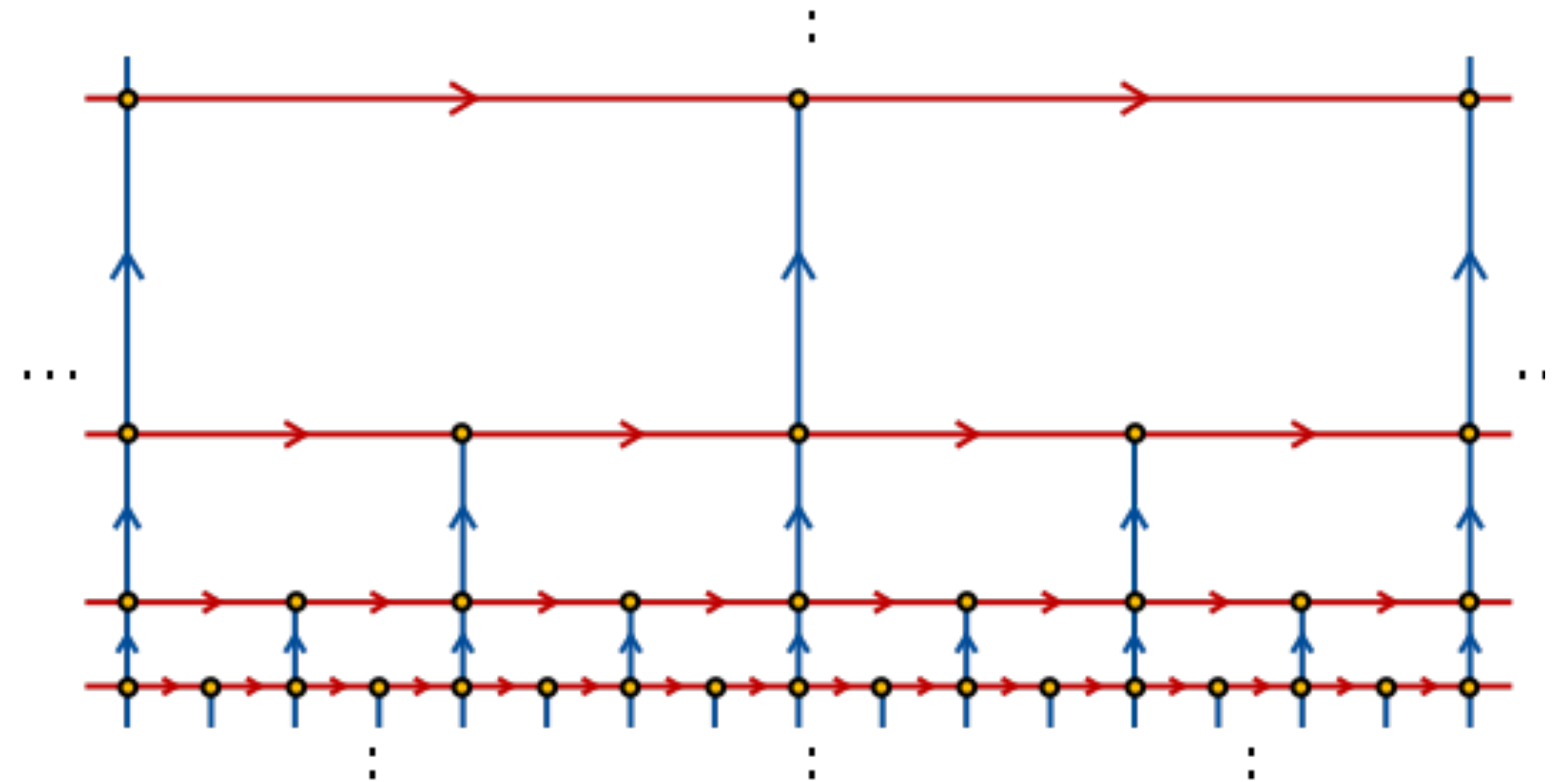
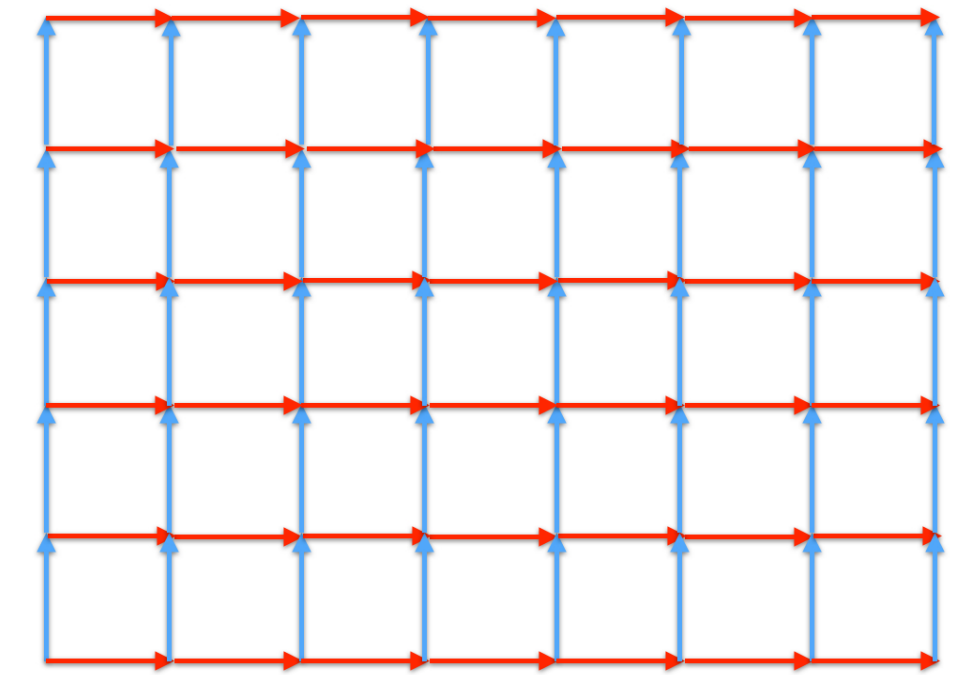
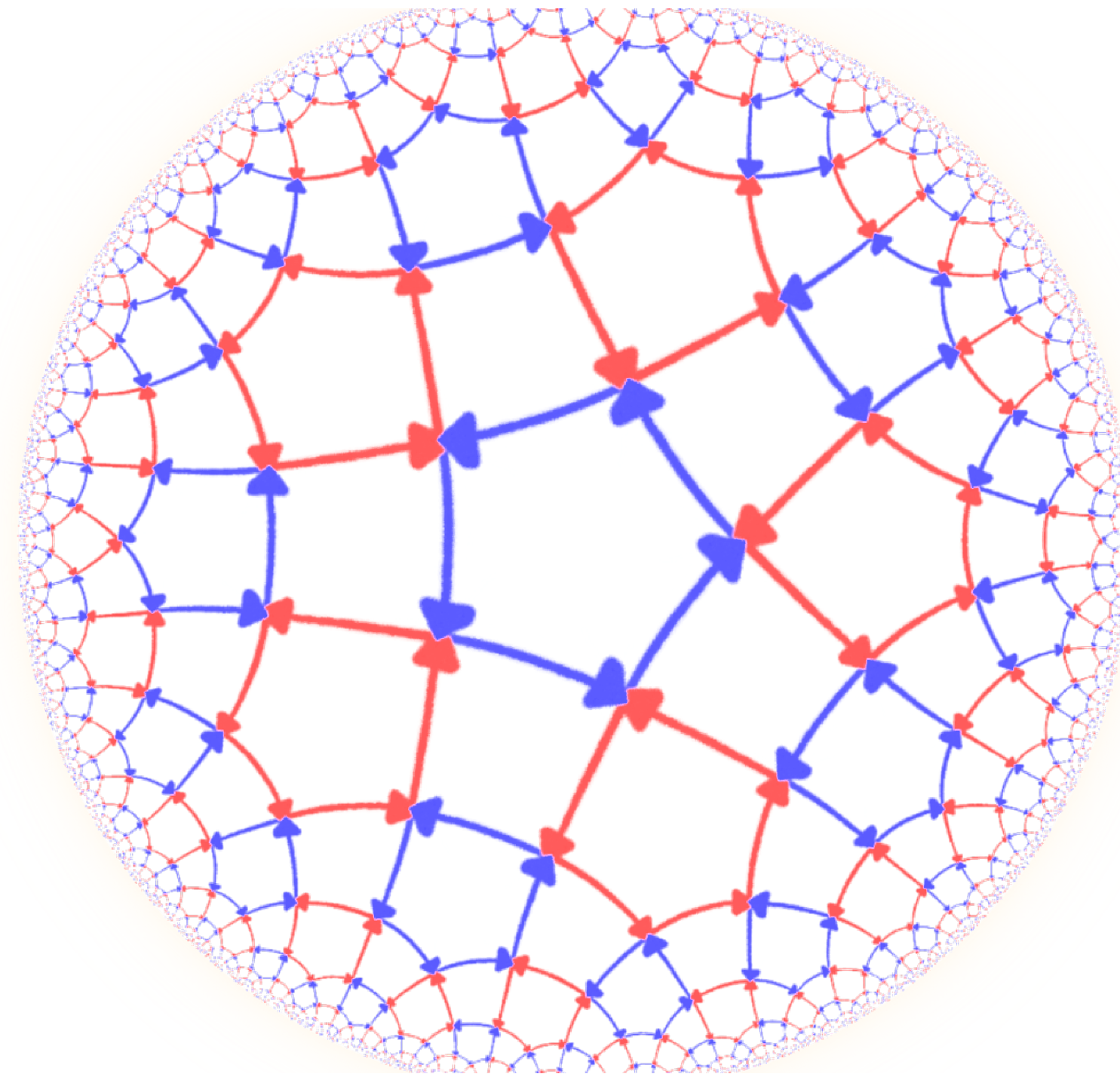
No two cells are discriminated

- Given a GUR $(G, A, \mathcal{V}^\dagger)$, we say that it is homogeneous if for every two cells $g_1, g_2 \in \mathcal{V}^\dagger$ does not discriminate them, but it does relatively to a third cell e
- For every pair (g_1, g_2) , there is \mathcal{I} such that for every $R_1 \ni g_1$ one finds $R_2 \ni g_2$ with $R_1 \bowtie_{\mathcal{I}} R_2$

Translation group

From homogeneity

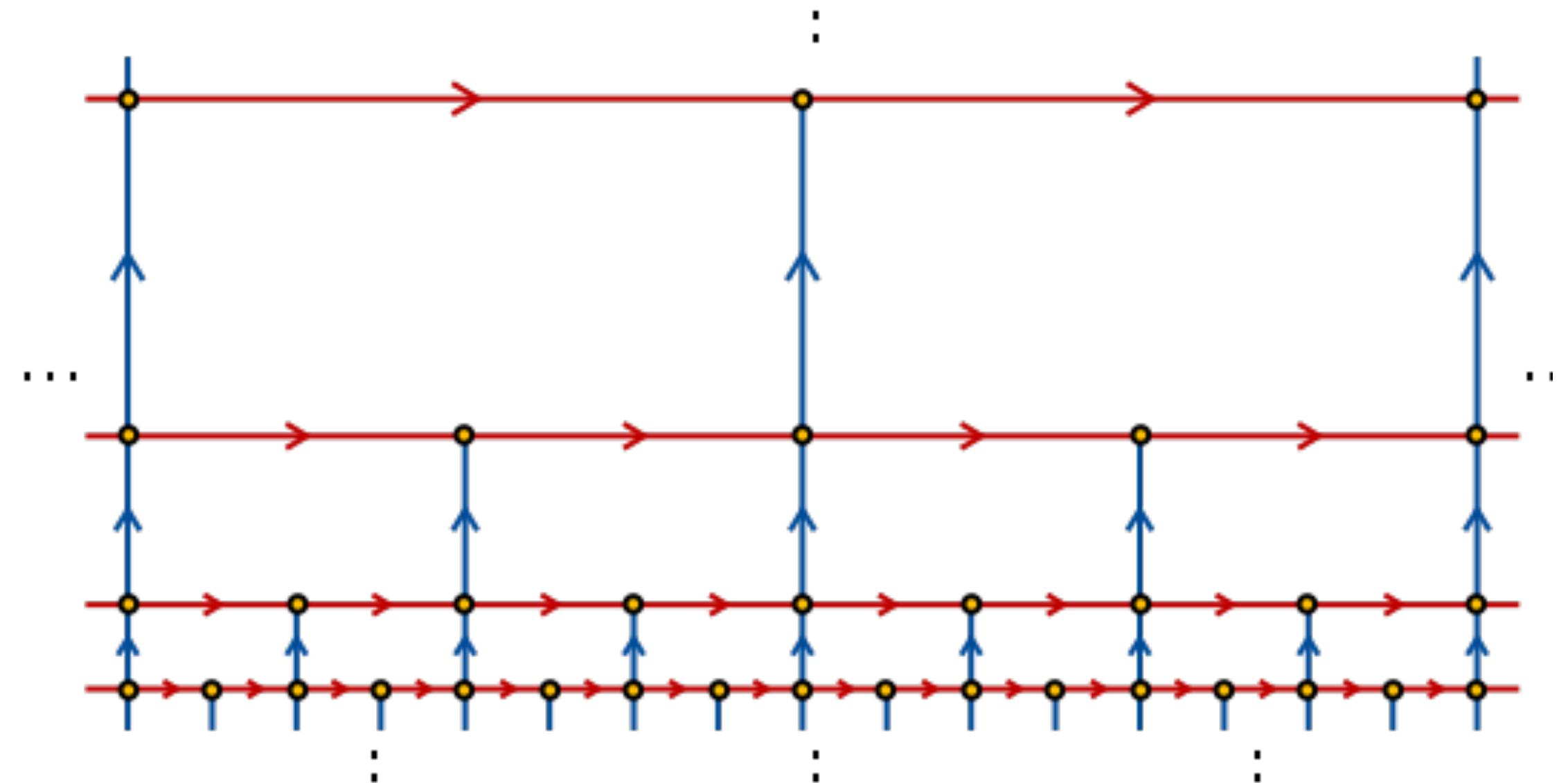
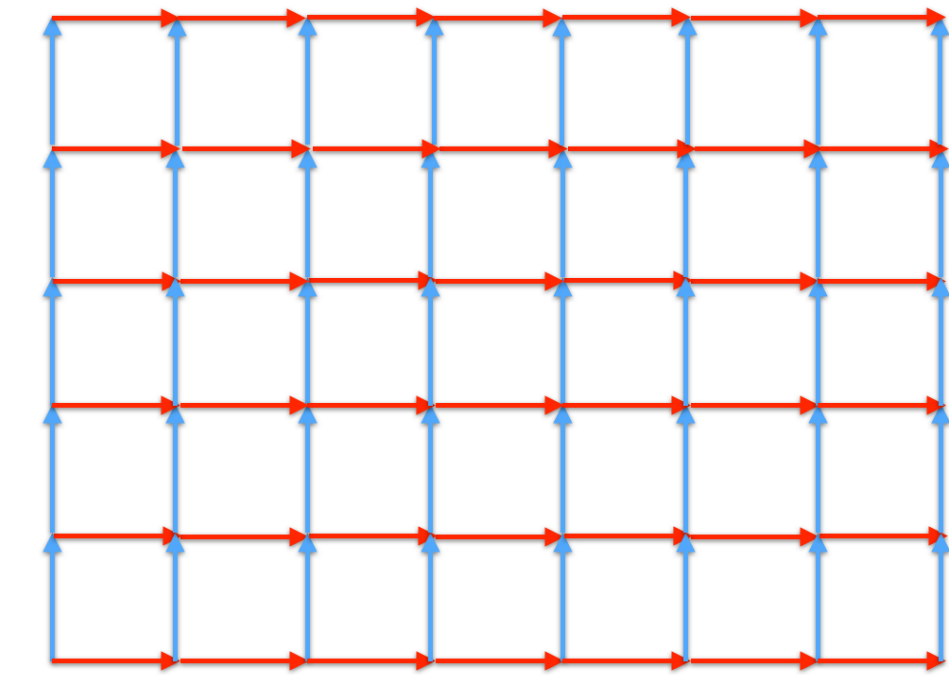
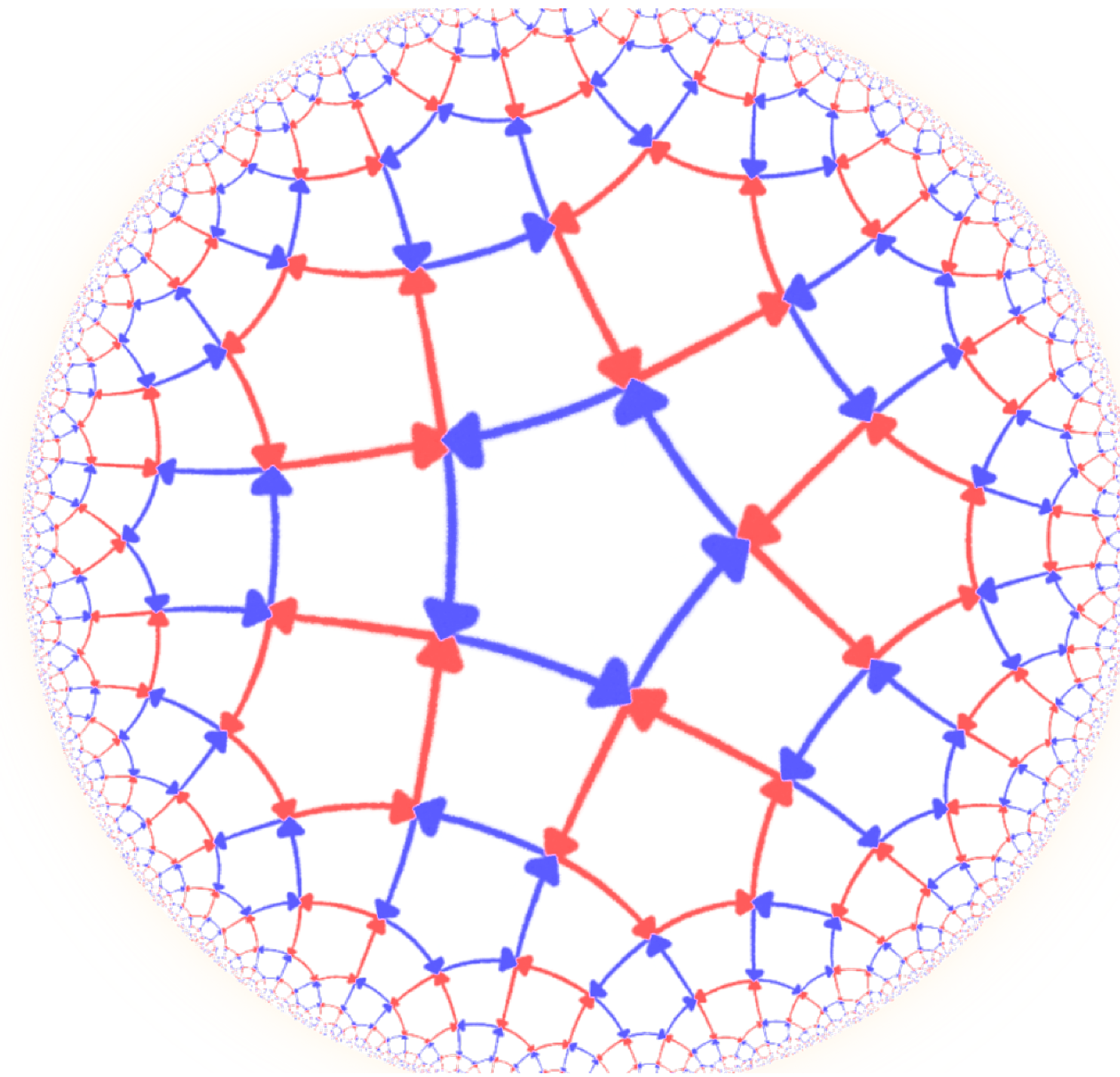
- We collect all the GURs \mathcal{T} such that $R_i \bowtie_{\mathcal{T}} R_j$, for every $R_i \ni g_i$, $R_j \ni g_j$, and $g_i, g_j \in G$



Translation group

From homogeneity

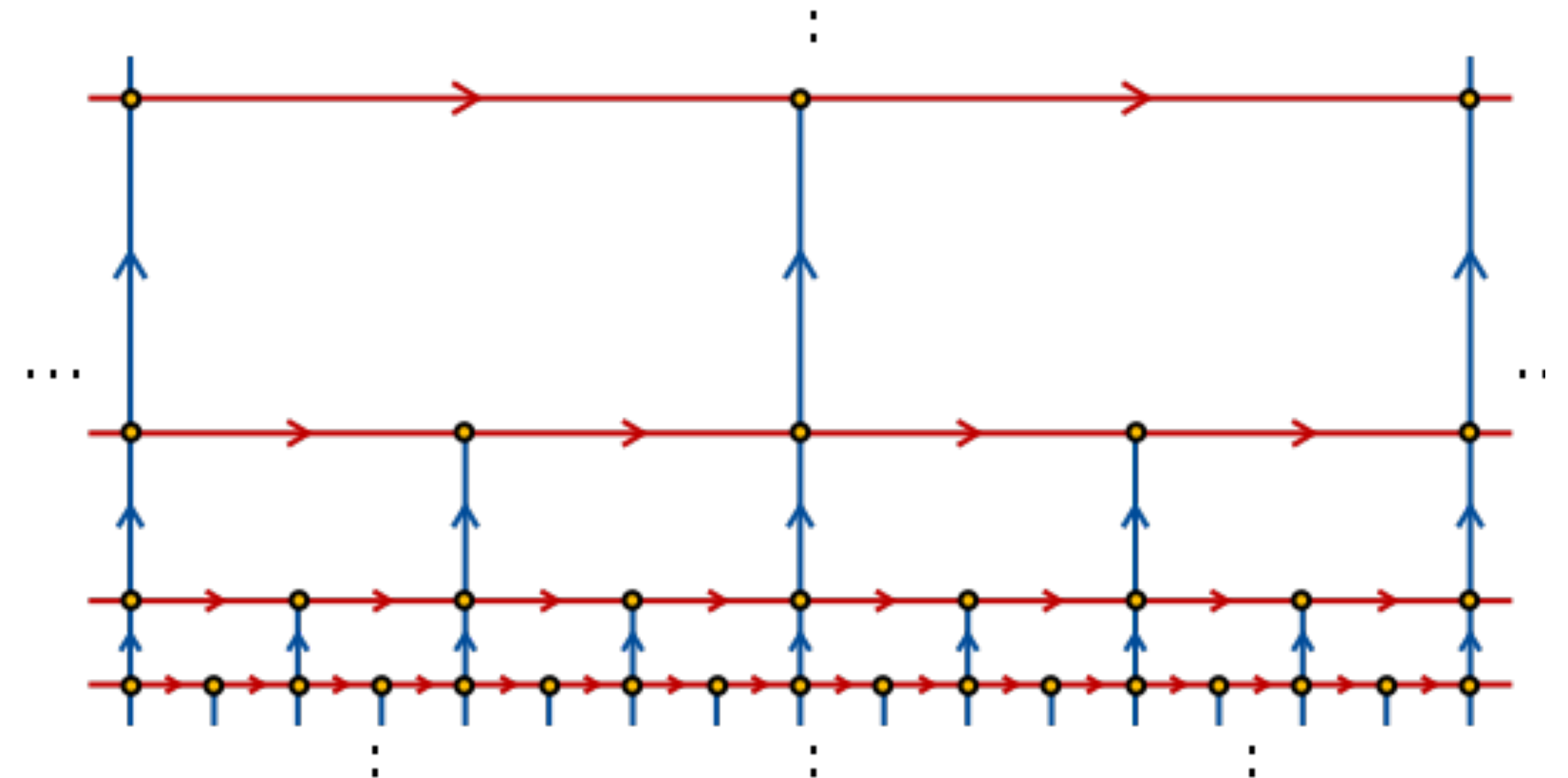
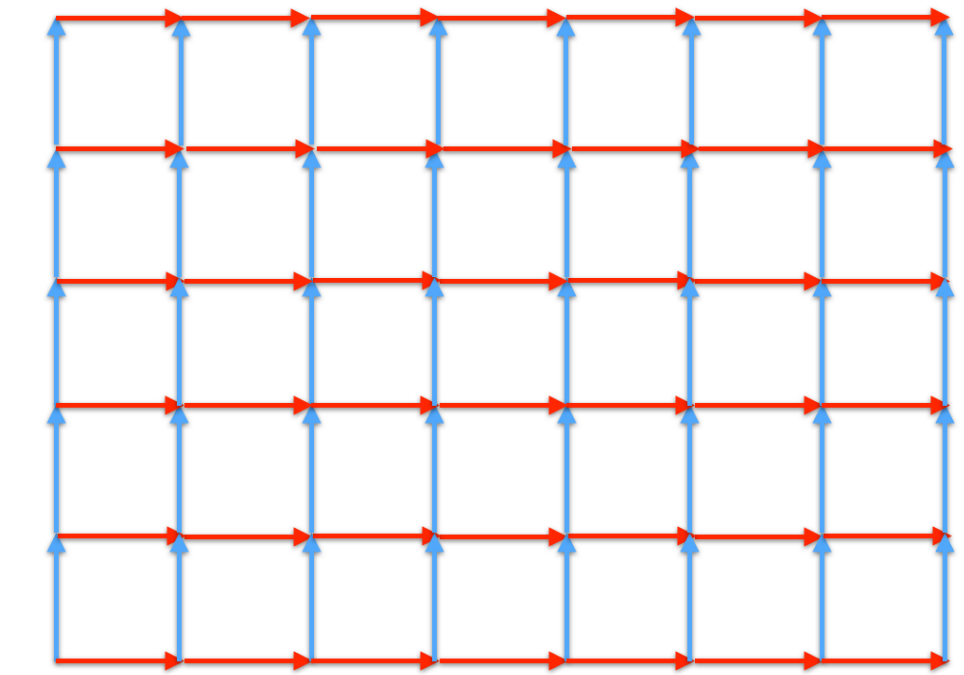
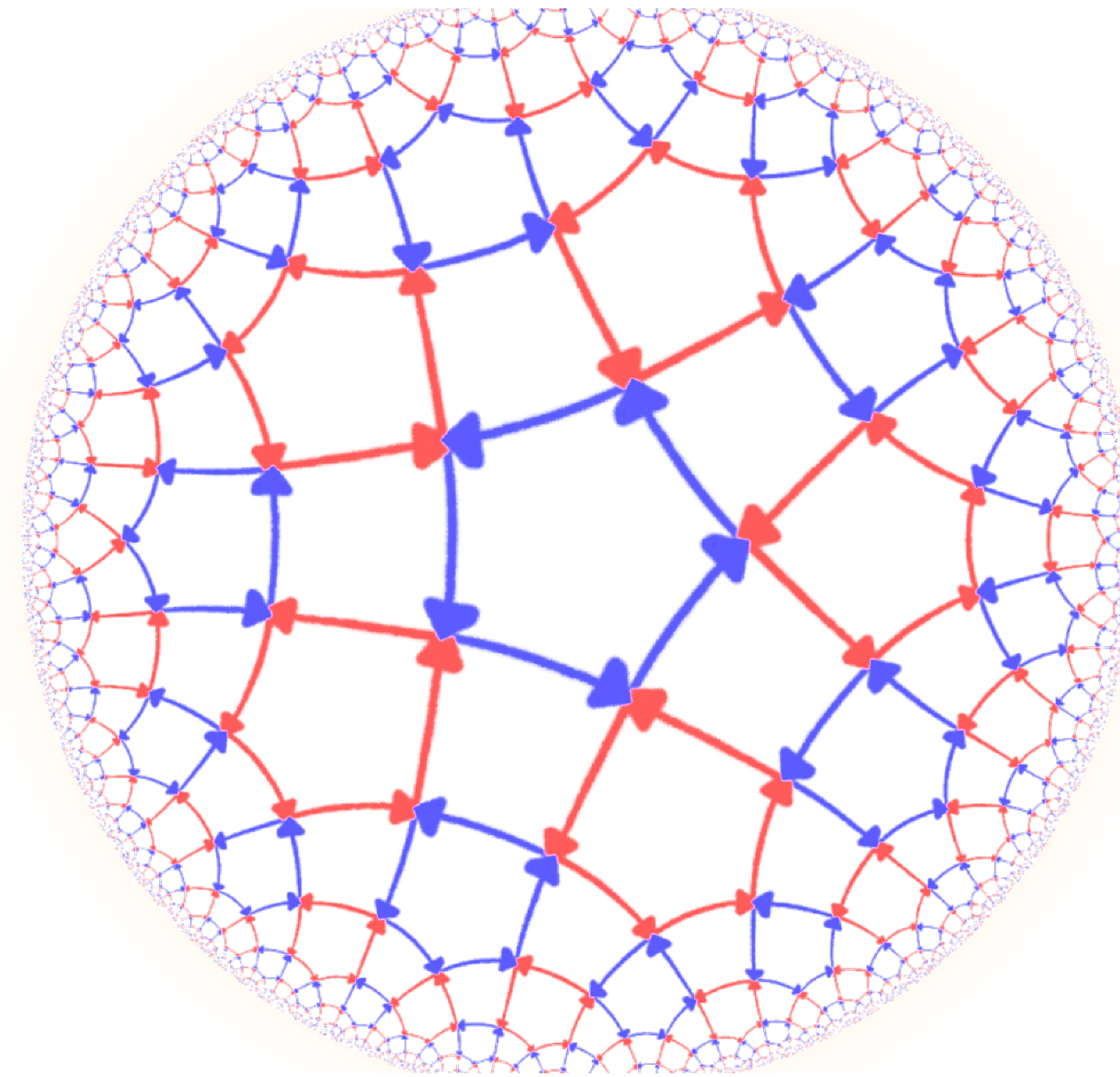
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- They form a group of permutations of G



Translation group

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- Since cells can be discriminated relatively to some third cell, there can be no fixed point unless $\mathcal{T} = \mathcal{I}$



Translation group

From homogeneity

- We collect all the GURs \mathcal{T} such that $R_i \bowtie_{\mathcal{T}} R_j$, for every $R_i \ni g_i$, $R_j \ni g_j$, and $g_i, g_j \in G$
- They form a group of permutations of G
- Since cells can be discriminated relatively to some third cell, there can be no fixed point unless $\mathcal{T} = \mathcal{I}$
- The group acts on itself transitively and freely: it is a group of translations

