

O_{ut of} T_{ime} O_{rdered}

Beni Yoshida (Perimeter)

[1] Chaos in quantum channel, Hosur, Qi, Roberts, BY 1511.04021

[2] Complexity by design, (in prep, with Daniel Roberts)

Scrambling

Out-of-time ordered correlation function

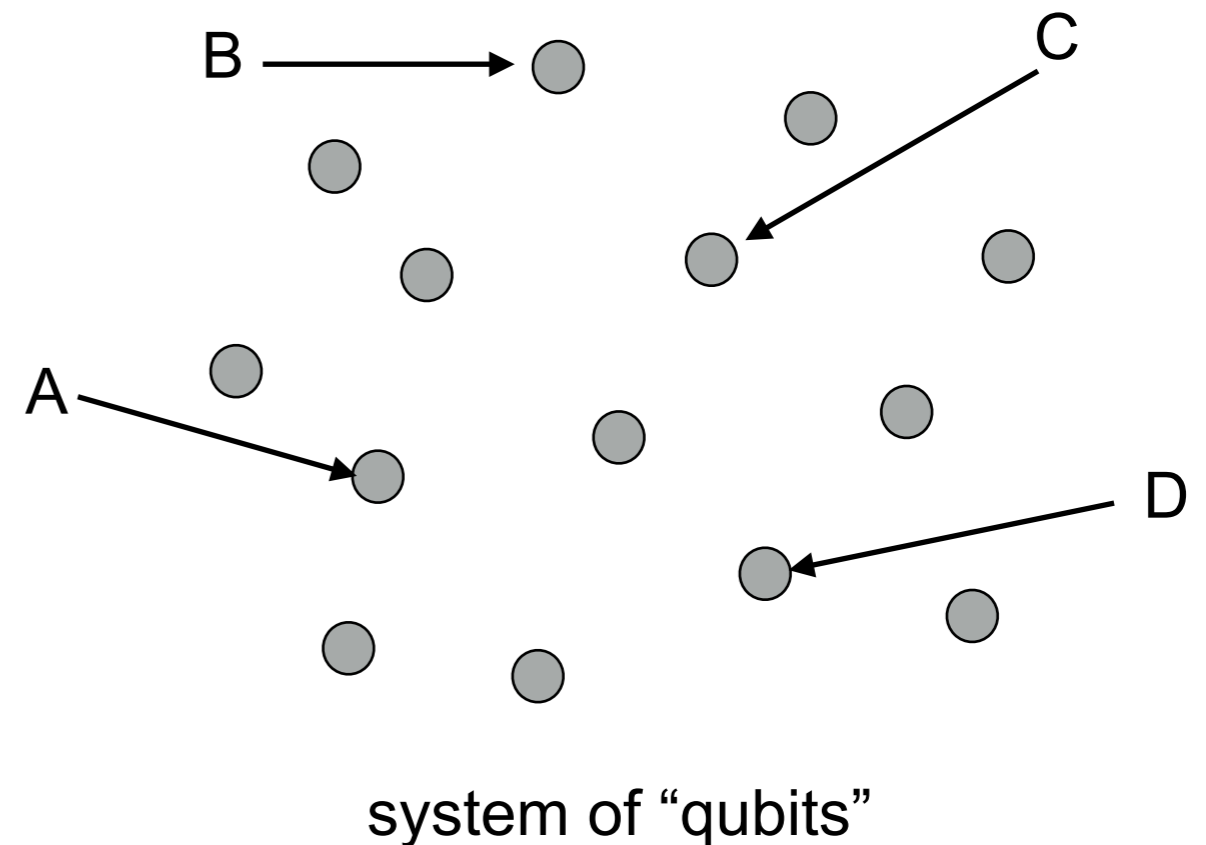
- Scrambling: **delocalization of quantum information**, hidden into non-local DOF.
- **Fast scrambling** is the defining feature of black holes.
- Out-of-time ordered (**OTO**) correlators can detect scrambling.
- Definition

$$\text{OTO} = \langle A(0)B(t)C(0)D(t) \rangle$$

local operators

$$B(t) = e^{-iHt} B(0) e^{iHt}$$

$$D(t) = e^{-iHt} D(0) e^{iHt}$$



Scrambling implies decay of OTO

- **Local perturbation** to an initial state cannot be detected by **any local measurement** on an output state.

- Consider $\text{OTO} = \langle A(0)B(t)A^\dagger(0)B^\dagger(t) \rangle$ (group commutator)

$$[A, B] = 0 \quad \text{then} \quad \text{OTO} = 1$$

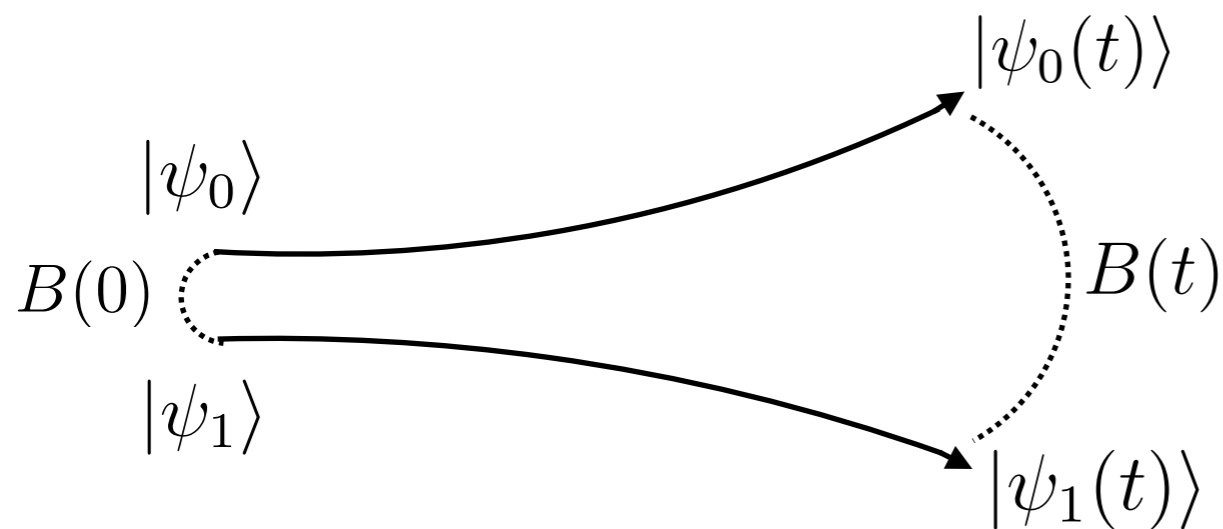
Non-commutativity between $A(0)$ and $B(t)$

$$\{A, B\} = 0 \quad \text{then} \quad \text{OTO} = -1$$

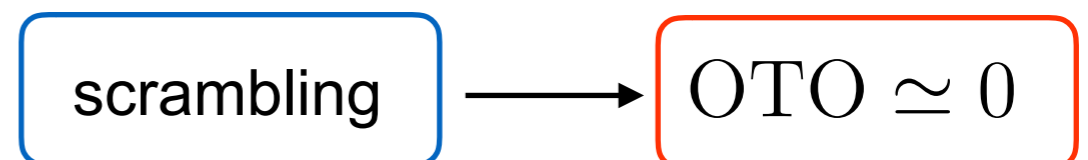
- Expand $B(t)$:

$$B(t) = e^{-iHt} B e^{iHt} = \sum_j \alpha_j P_j$$

high-weight Pauli operators



scrambling/chaos (butterfly effect)

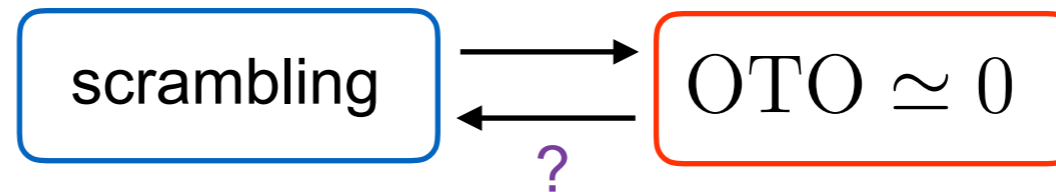


[Roberts-Susskind-Stanford]

Key Questions

- How do we define scrambling ?
- Quantum information theoretic meaning of OTO ?

- Is the converse true ?



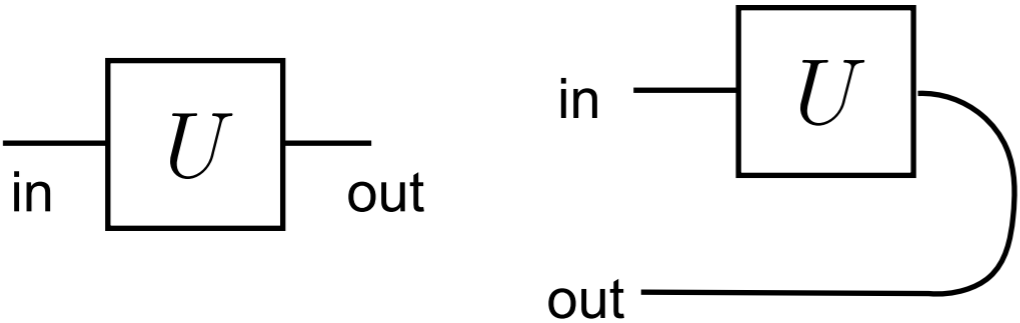
- Relation to entanglement entropy (and geometric quantities) ?

State-Channel duality

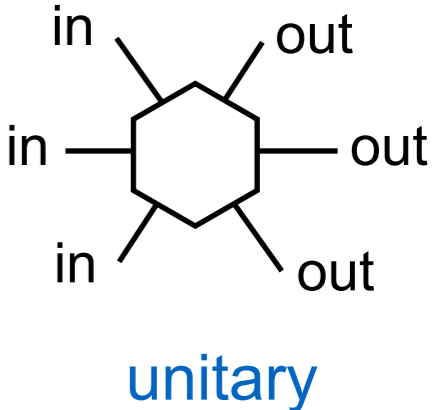
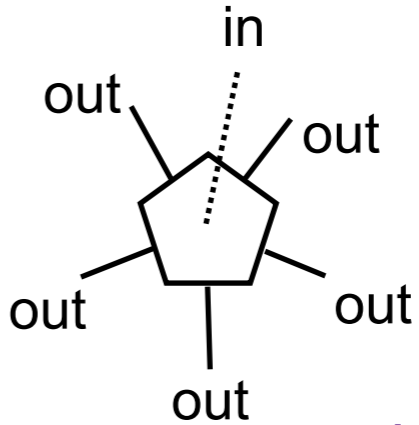
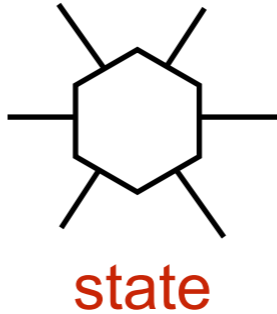
- Quantum channel on n qubits can be viewed as a state on 2n qubits.

unitary operator as a state (T=infty)

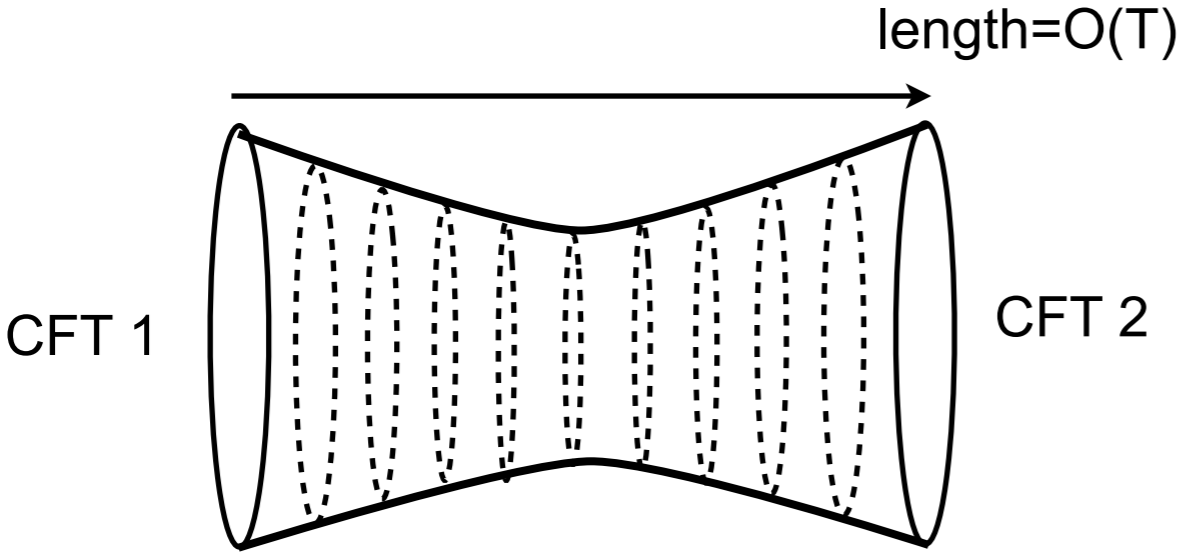
$$U = \sum_{i,j} U_{i,j} |i\rangle\langle j| \quad |U\rangle = \sum_{i,j} U_{i,j} |i\rangle \otimes |j\rangle$$



cf) perfect tensor



- Thermofield double state (finite T)



ER bridge = quantum channel

$$\rho_{in} = \frac{e^{-\beta H}}{\text{Tr } e^{-\beta H}} \quad \rho_{out} = \frac{e^{-\beta H}}{\text{Tr } e^{-\beta H}}$$

$$|\Psi\rangle = \sum_j e^{-\beta E_j/2} e^{-iE_j t} |\psi_j\rangle \otimes |\psi_j\rangle$$

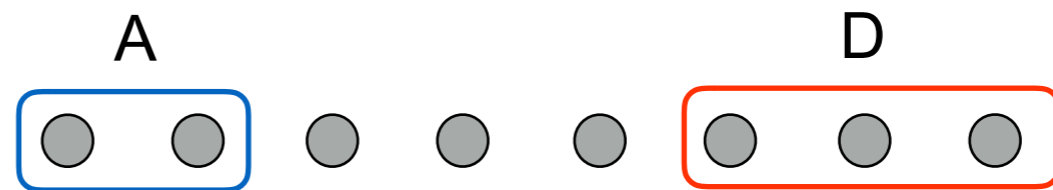
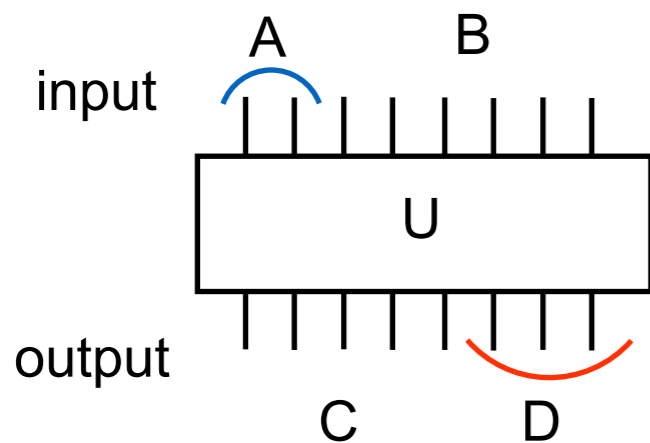
[Choi, Jamiolkowski, Hayden-Preskill]

Average value of OTO

- Average of OTO over local operators A and D at T=infty

$$\begin{aligned}
 \left| \langle A(0) D(t) A^\dagger(0) D^\dagger(t) \rangle \right| &= \frac{1}{4^{a+d}} \sum_{A,D} \langle A(0) D(t) A^\dagger(0) D^\dagger(t) \rangle \\
 &= 2^{n-a-d} S_{BD}^{(2)}
 \end{aligned}$$

average over A, D Pauli operators (unitary 1-design)
Renyi-2 entropy



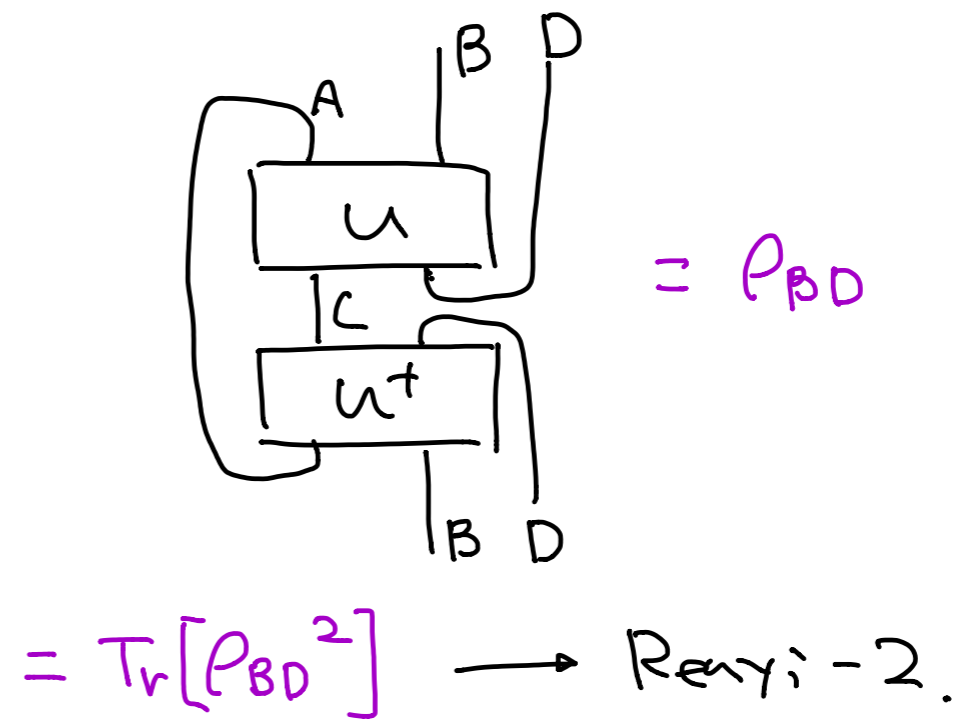
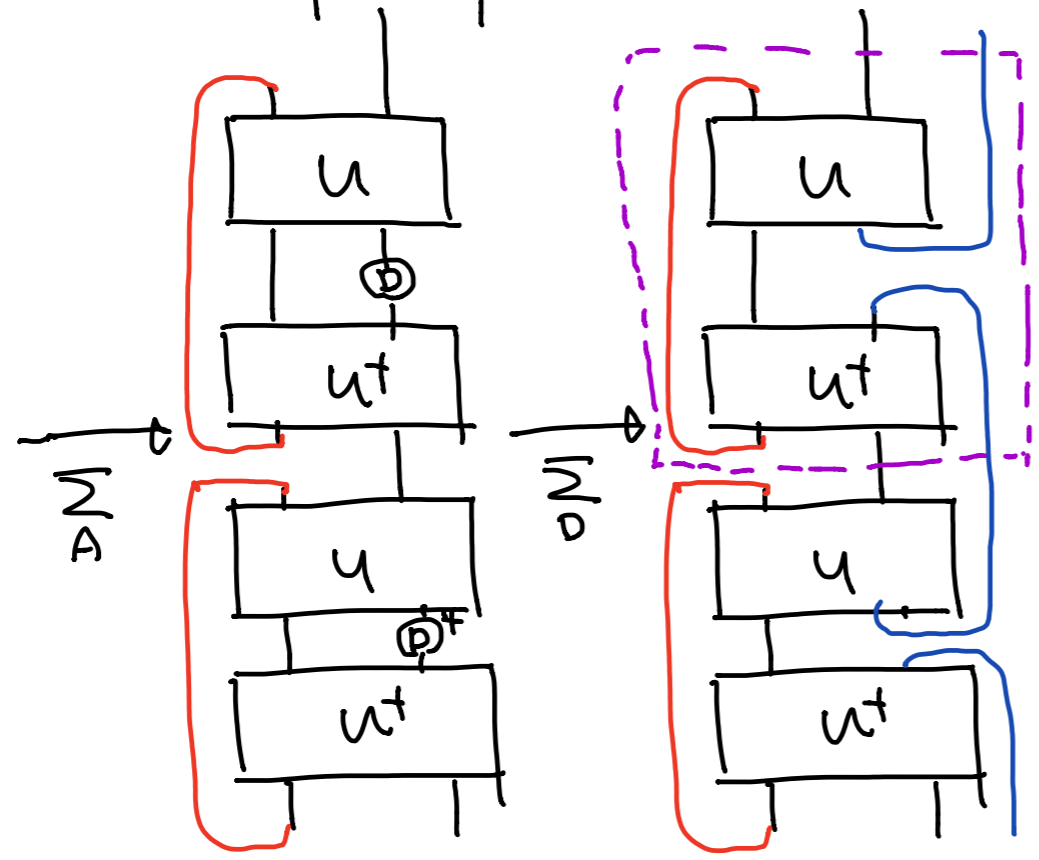
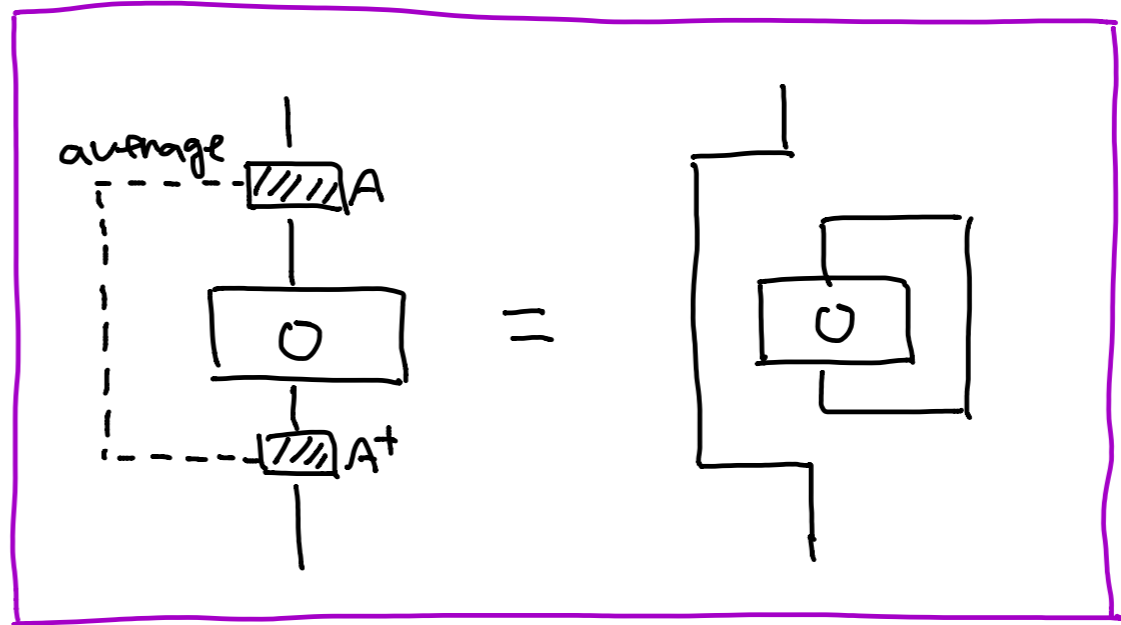
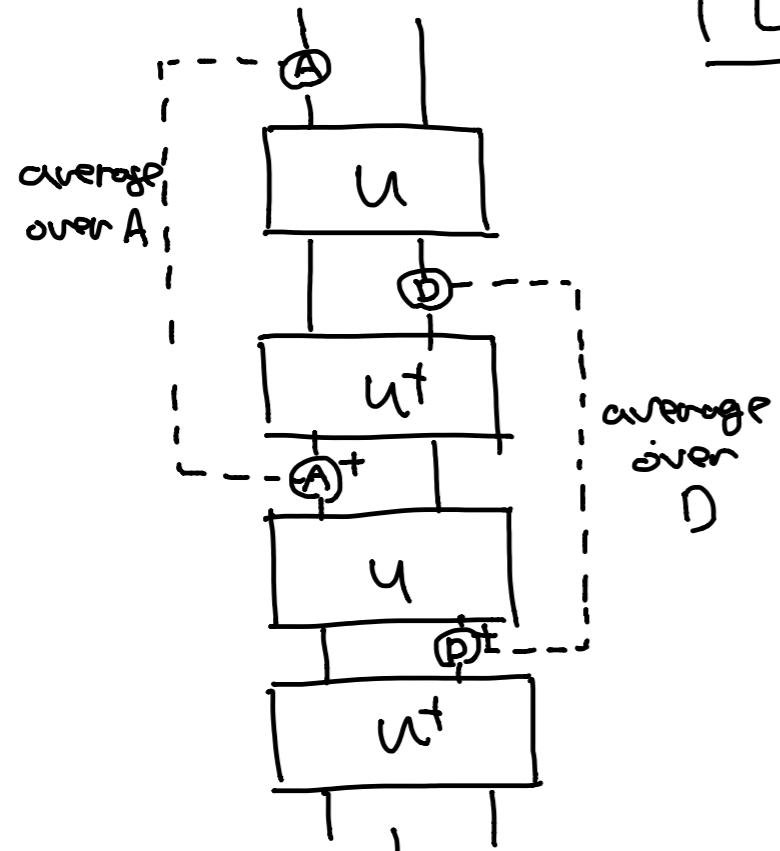
- If $OTO \simeq 0$ then, $S_{BD}^{(2)}$ is large

This implies the **mutual information** $I_{BD}^{(2)} = S_B^{(2)} + S_D^{(2)} - S_{BD}^{(2)}$ is small

B and D are not correlated, so the system is **scrambling**.

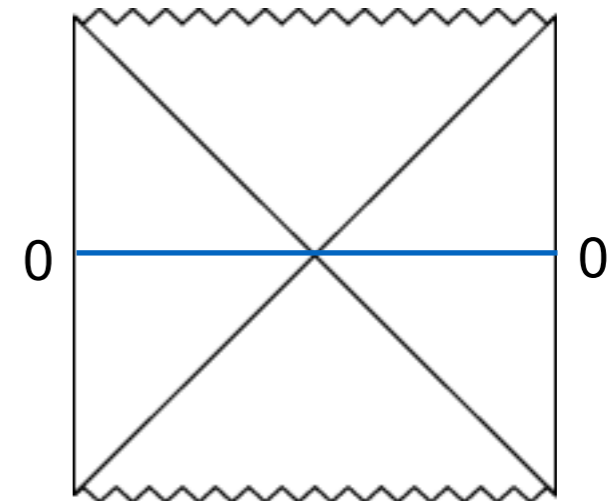
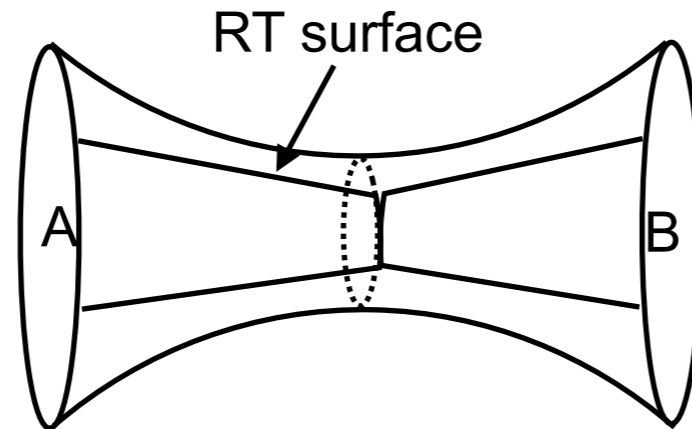
- For finite T, we consider **TFD state**.

Derivation



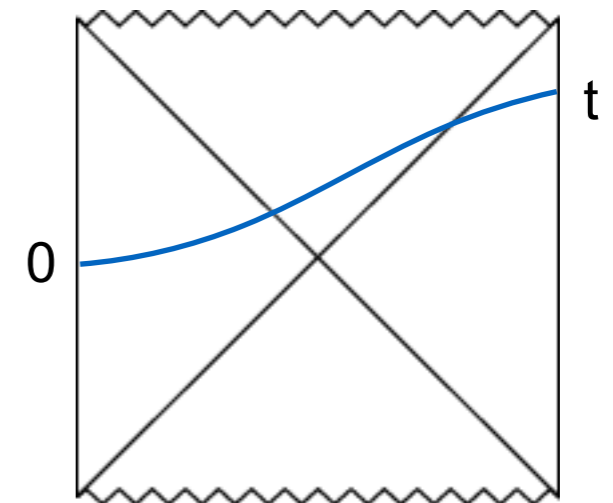
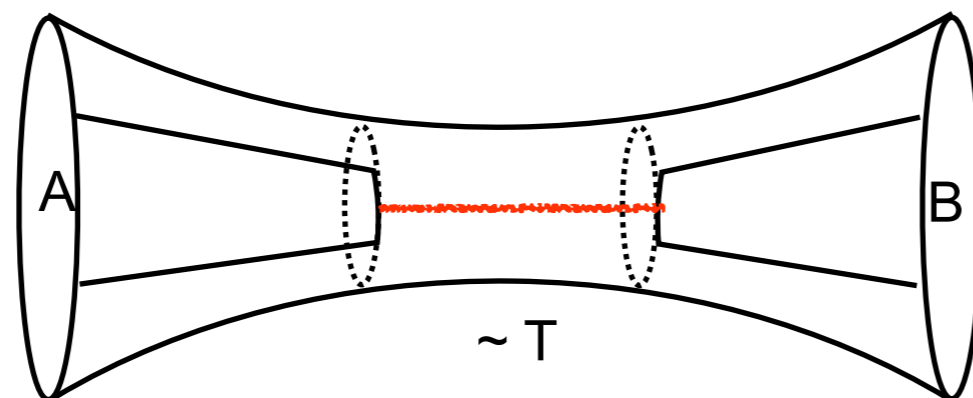
Scrambling in AdS black hole

TFD(0)



mutual information $I(A, B)$ is large

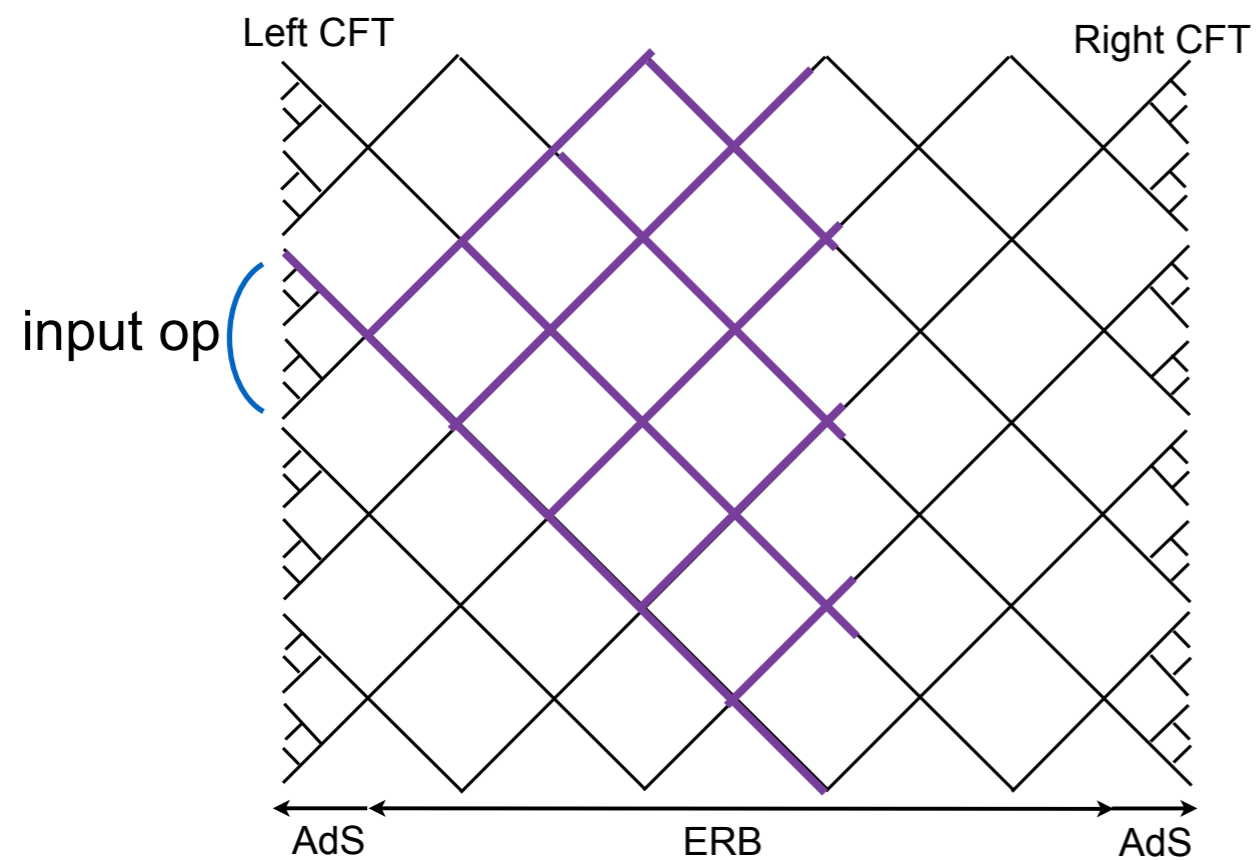
TFD(T)



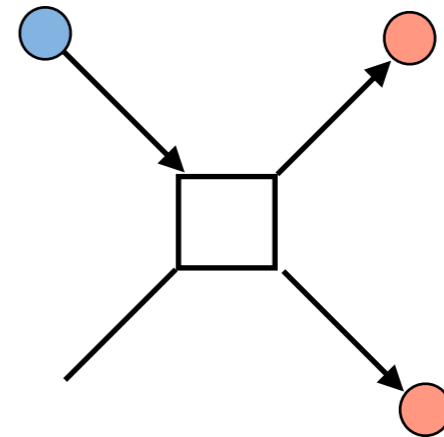
mutual information $I(A, B)$ is (almost) zero

Toy model of the ER bridge

- Operators grow ballistically, leading to decay of OTO correlators.



network of perfect
(or random) tensors

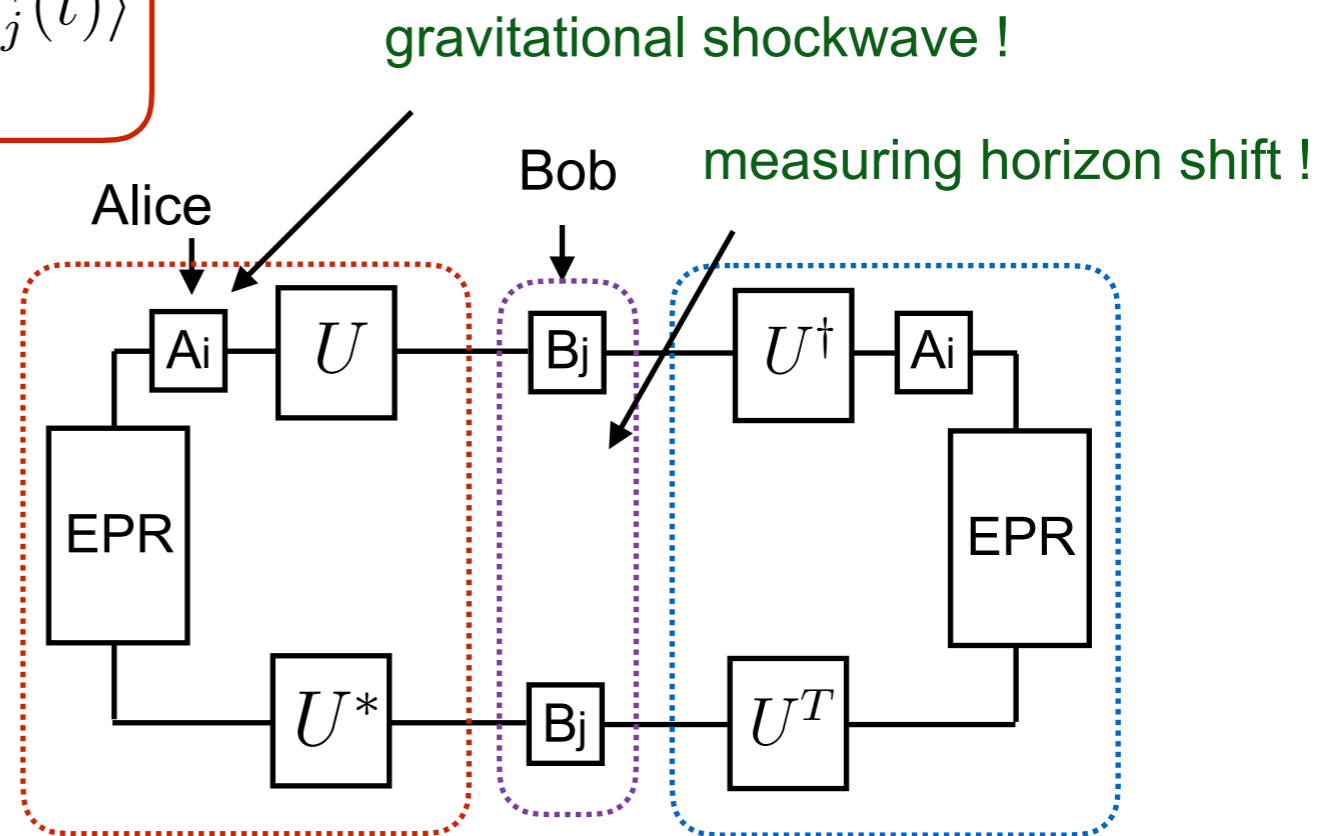
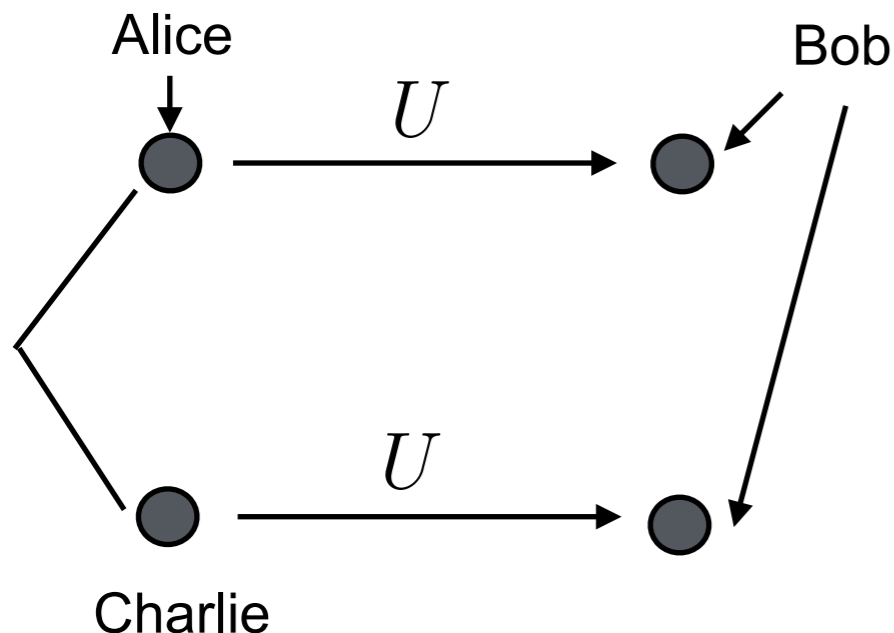


OTO average for Alice and Bob

- Alice and Bob are playing a catch ball. Alice may apply some **perturbation**
- Bob asks **Charlie** to throw the same ball, and compare it with Alice's ball.

- Alice applies **A_i with equal probability**. \longrightarrow superoperator $\sum_i A_i(\cdot)A_i^\dagger$
- Bob performs a **joint measurement on two systems**. \longrightarrow $\sum_j B_j \otimes B_j^\dagger$
SWAP operator

The outcome : $\sum_{i,j} \langle A_i(0)B_j(t)A_i^\dagger(0)B_j^\dagger(t) \rangle$



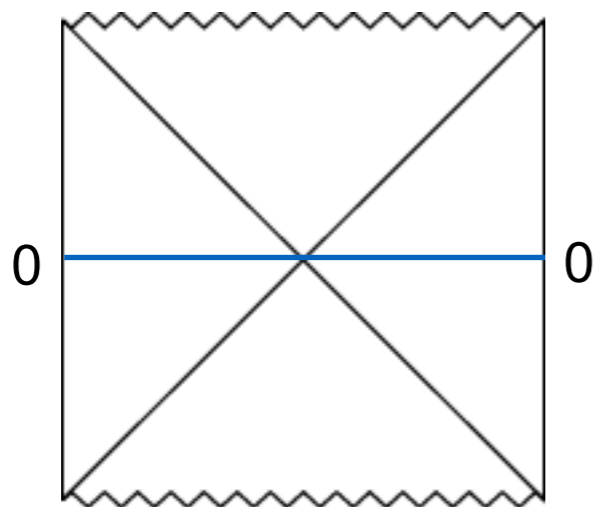
Complexity

Can we detect complexity growth ?

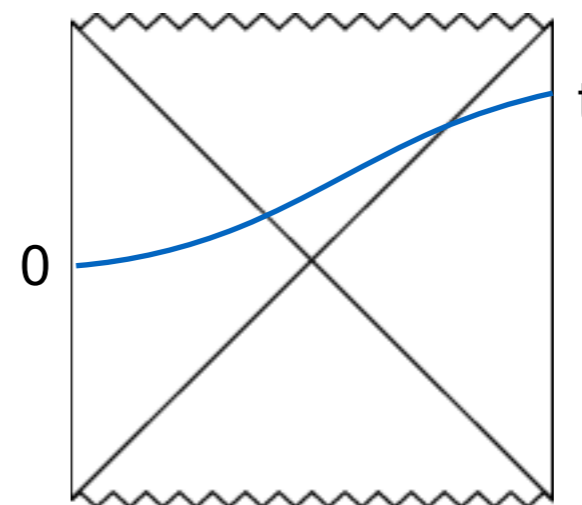
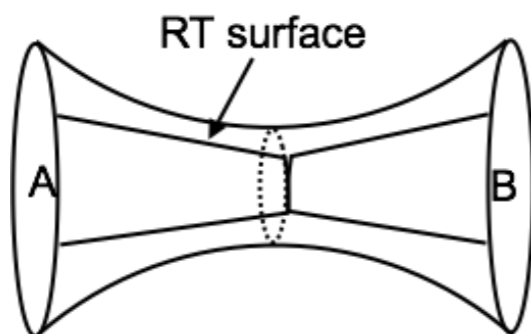
- The **complexity** of the TFD state still keeps growing ?

Reference state \longrightarrow Target state

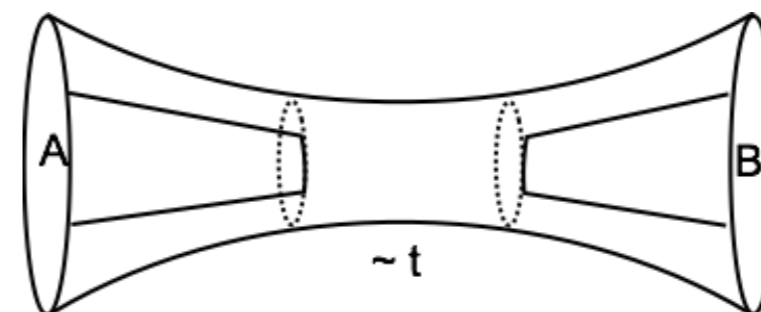
How many **quantum gates** do we need ?



$$|\Psi(0)\rangle = \sum_j e^{-\beta E_j/2} |\psi_j\rangle \otimes |\psi_j\rangle$$

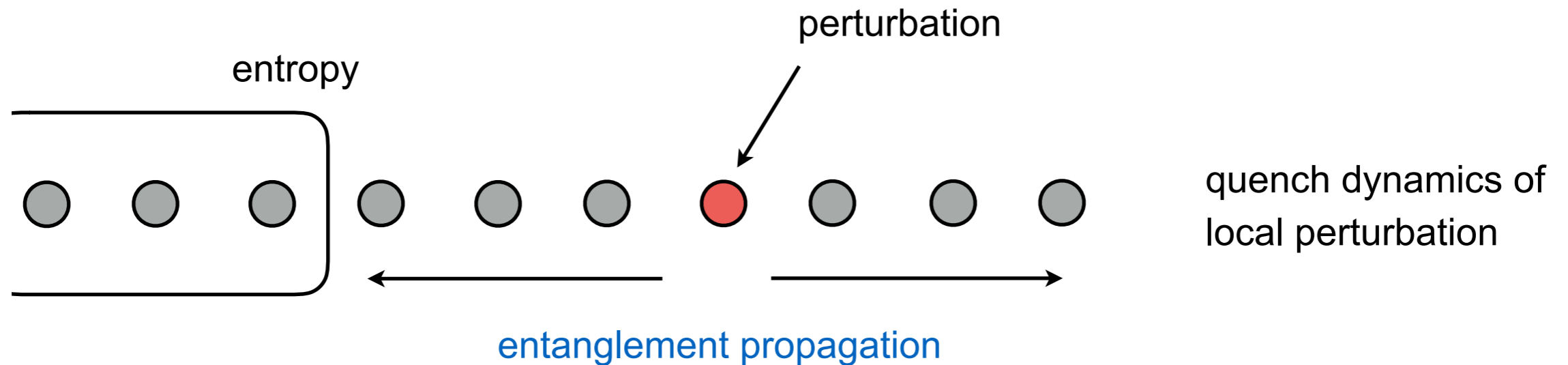


$$|\Psi(t)\rangle = \sum_j e^{-\beta E_j/2} e^{-iE_j t} |\psi_j\rangle \otimes |\psi_j\rangle$$

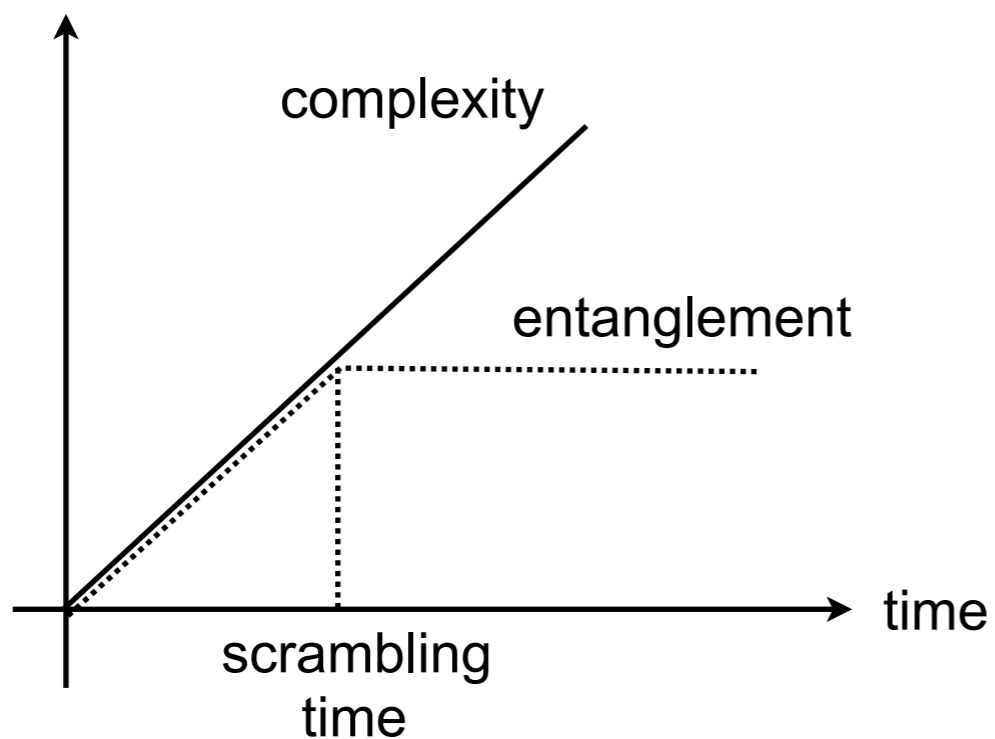


Entanglement can detect complexity ?

- Entropy grows as the complexity grows.



- After the scrambling time, **entanglement entropies get saturated**.



Entanglement is not enough !

Unitary k-design (complexity of randomness)

- Imagine an ensemble of **Haar** (uniformly distributed) random unitary (on n qubits)

A **typical operator** in the Haar ensemble has **$\exp(n)$ complexity**.

- Consider **k copies** of the system (kn qubits in total) and consider “k-fold twirl” :

$$\rho \rightarrow K(\rho) = \int dU \underbrace{(U \otimes \dots \otimes U)}_{k \text{ copies}} \rho (U^\dagger \otimes \dots \otimes U^\dagger)$$

kn qubits
n qubits

- We think of **approximating the Haar random ensemble** by some ensembles which are easier to generate.

$$\{p_j, U_j\} \quad \rho \rightarrow \Phi(\rho) = \sum_j p_j (U_j \otimes \dots \otimes U_j) \rho (U_j^\dagger \otimes \dots \otimes U_j^\dagger)$$

- If $\Phi = K$, then $\{p_j, U_j\}$ is said to form a **unitary k-design**,
i.e. it is as good as Haar up to k-th moment.

Examples of k-design

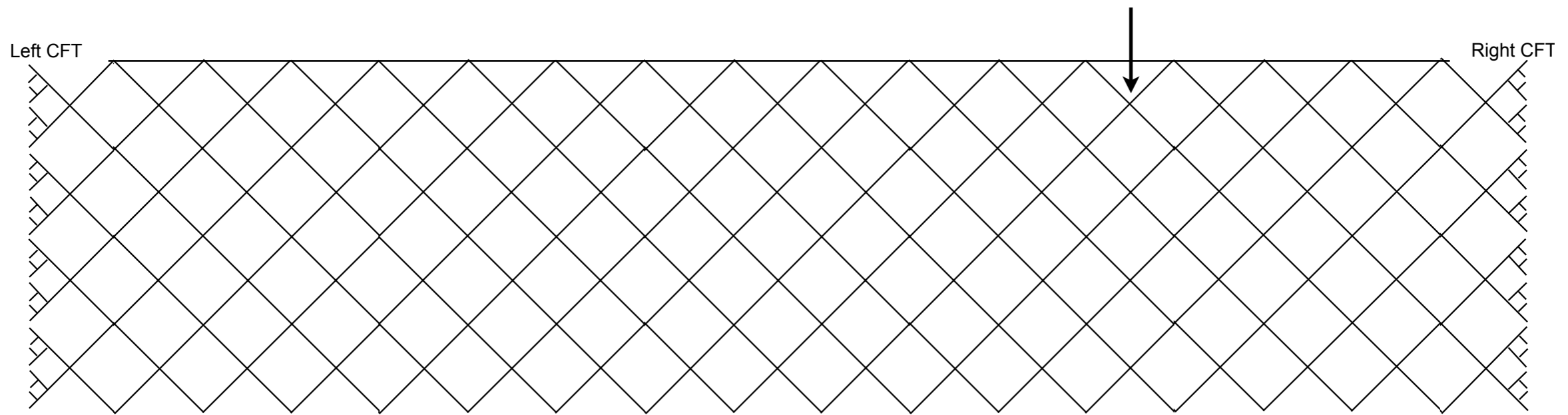
- A group of **Pauli operators** is 1-design.

I, X, Y, Z, ZZ, YYZY,

- A group of **Clifford operators** is 2-design.

Clifford operators can prepare arbitrary stabilizer states (eg **perfect tensors**).

- A toy model of the **wormhole** (Hosur, Qi, Roberts, BY) random unitary



this object forms an **approximate k-design**.

Lower bound on complexity

- Imagine a system of n qubits. ($d=2^n$ states in the Hilbert space).
- If an ensemble of unitary operators formed a k -design, the ensemble must contain at least

$$|\text{Supp}(\mathcal{E})| \geq \binom{d+k-1}{k}^2. \quad (\text{due to the Schur-Weyl duality})$$

- At each step, the number of implementable quantum gates is

$$\simeq gn^2 \quad (\text{g: number of different 2-qubit gates})$$

- In T step, the number of implementable quantum gates is $(gn^2)^T$

- A typical operator in k -design has a complexity of at least

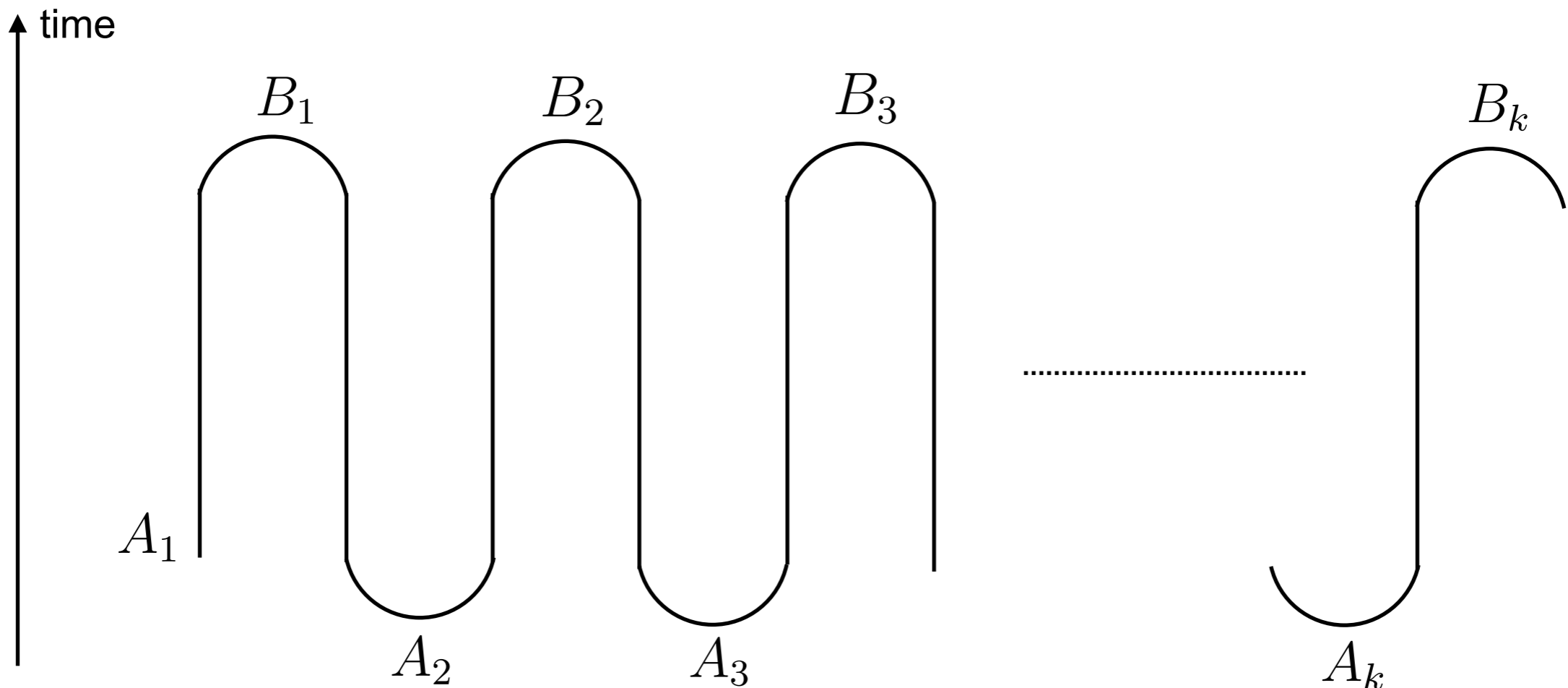
$$\frac{2kn \log(2)}{\log(gn^2)}$$

roughly linear in k and n !

How do we detect the design ?

- Answer : Out-of-time ordered correlation functions
- 2k-point OTO correlators can detect k-design.

$$\langle A_1(0)B_1(t)A_2(0)B_2(t) \cdots A_k(0)B_k(t) \rangle$$



OTO determines k-fold channel

Consider a **k-fold twirl** over an arbitrary ensemble $\mathcal{E} = \{p_j, U_j\}$

$$\Phi_{\mathcal{E}}(\rho) = \sum_j p_j (U_j \otimes \cdots \otimes U_j) \rho \underbrace{(U_j^\dagger \otimes \cdots \otimes U_j^\dagger)}_{k \text{ copies}}. \quad (\text{quantum channel})$$

The density matrix can be expanded by Pauli operators, so we are interested in

$$\Phi_{\mathcal{E}}(\underbrace{B_1 \otimes \cdots \otimes B_k}_{\text{Pauli op}}) = \sum_{C_1, \dots, C_k} \underbrace{\gamma_{C_1, \dots, C_k}}_{\text{knowing these numbers give complete characterization of the channel.}} (C_1 \otimes \cdots \otimes C_k)$$

Assume that we know averages of **2k-point** OTO correlators for Pauli operators

$$\underbrace{\alpha_{A_1, \dots, A_k}}_{\text{We know these numbers.}} = \left| \left\langle A_1(0) B_1(t) \cdots A_k(0) B_k(t) \right\rangle_{T=\infty} \right|_{\mathcal{E}} \quad B_i(t) = U B_j U^\dagger$$

Question

Can we determine γ_{C_1, \dots, C_n} from α_{A_1, \dots, A_k} ?

Yes

Theorem : OTO and k-fold twirl

Goal : $\alpha_{A_1, \dots, A_k} \rightarrow \gamma_{C_1, \dots, C_k}$

$$\alpha_{A_1, \dots, A_k} = \left| \left\langle A_1(0) B_1(t) \cdots A_k(0) B_k(t) \right\rangle_{T=\infty} \right|_{\mathcal{E}}$$

$$\Phi_{\mathcal{E}}(B_1 \otimes \cdots \otimes B_k) = \sum_{C_1, \dots, C_k} \gamma_{C_1, \dots, C_k} (C_1 \otimes \cdots \otimes C_k)$$

Define : $M_{A_1, \dots, A_k}^{C_1, \dots, C_k} = \text{Tr}[A_1 C_1 \cdots A_k C_k]$

$$M_{C_1, \dots, C_k}^{\dagger A_1, \dots, A_k} = \text{Tr}[C_k^{\dagger} A_k^{\dagger} \cdots C_1^{\dagger} A_1^{\dagger}],$$

$$\gamma_{C_1, \dots, C_k} \propto M_{C_1, \dots, C_k}^{\dagger A_1, \dots, A_k} \cdot \alpha_{A_1, \dots, A_k}$$

Effective “design” of an ensemble

- We need to know OTO values for Haar random (or k-design) **in advance**.

This is possible by using some heavy math machineries.

- 4m-point OTO correlation functions that are related to **shockwave geometries**.

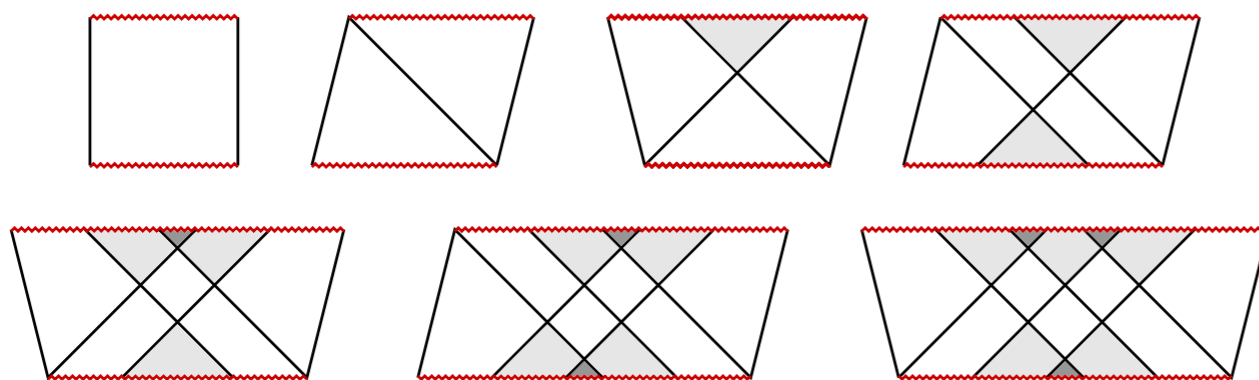
$$|4m\text{-point}|_{\text{Haar}} \sim \frac{1}{d^{2m}}$$

$$|4m\text{-point}|_{\text{Clifford}} \sim \frac{1}{d^2}$$

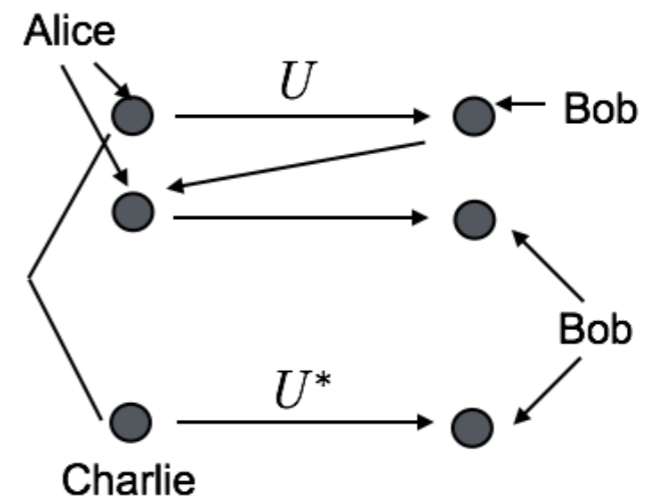
2-design

$$|4m\text{-point}| \sim \frac{1}{d^{\mathcal{D}}}$$

“effective design” of the ensemble



shockwave geometries



catch ball setup

Growth of design in an ER bridge ?

How do we define “design” in an ER bridge ?

- Unitary t-design considers an **ensemble of unitary operators**.
- Time-evolution of an ER bridge is given by **a single Hamiltonian H**.

Maybe, we can consider an **ensemble of Hamiltonians** ?

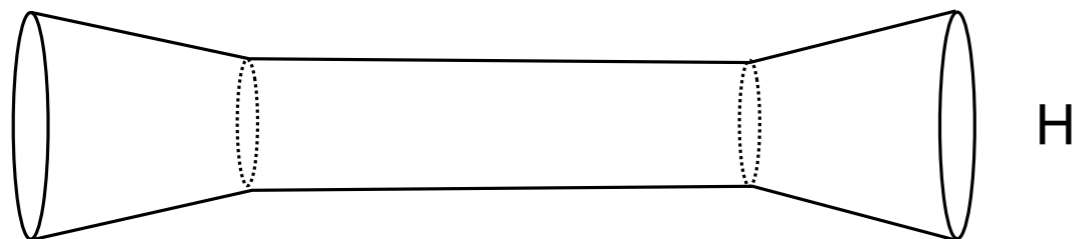
eg) **Sachdev-Ye-Kitaev** model

$$H = \frac{1}{4!} \sum_{i,j,k,l=1}^N J_{ijkl} \chi_i \chi_j \chi_k \chi_l$$

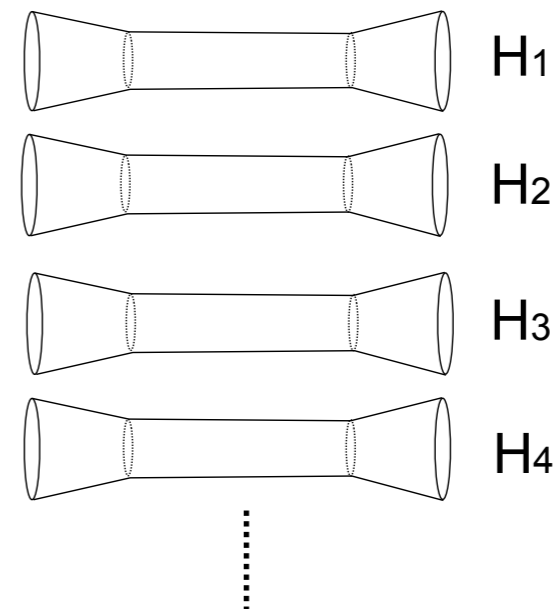
random variables

we can compute **disorder average** analytically

Or, we can imagine very high-energy DOF, which can be integrated out.



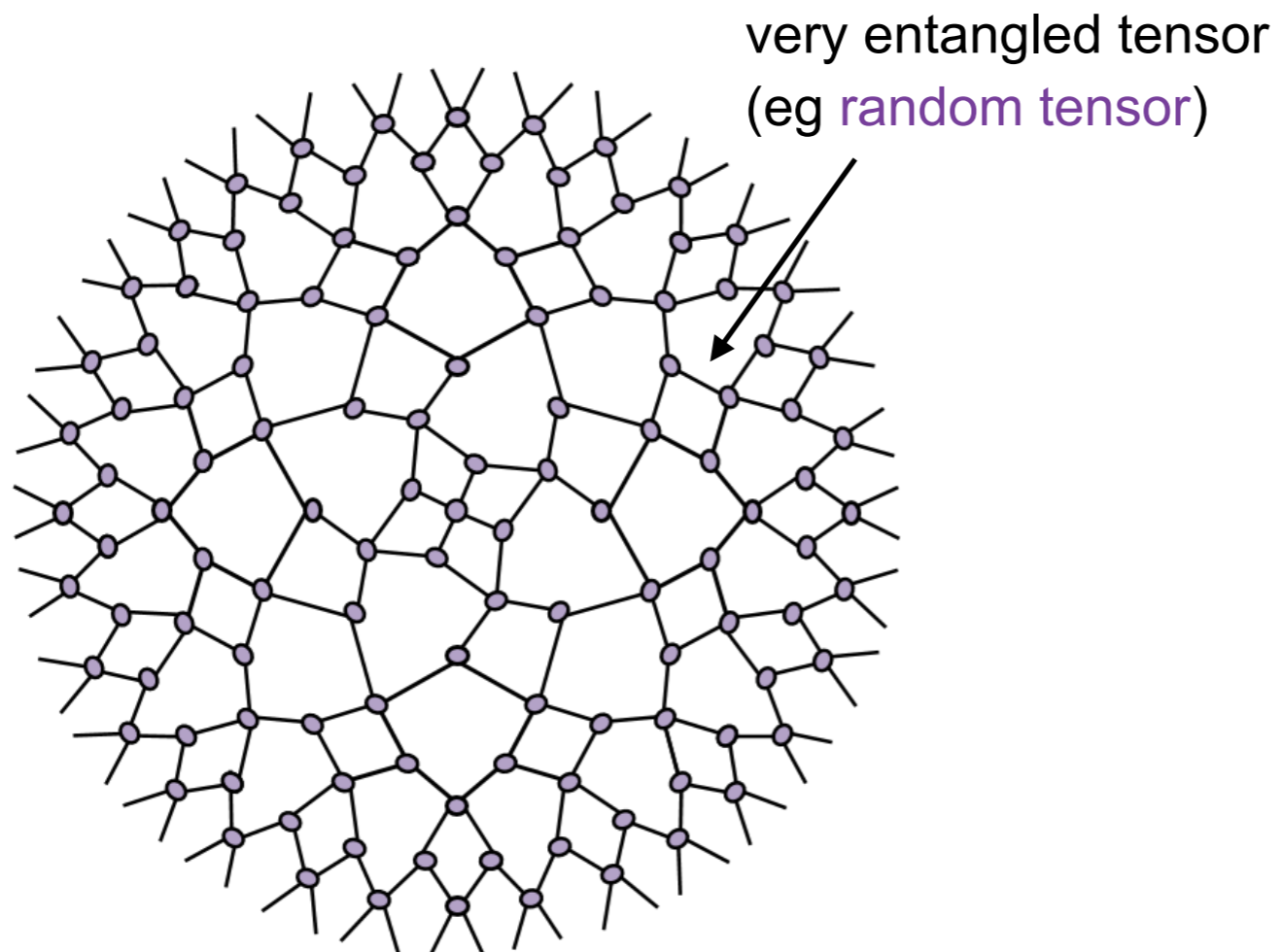
H



Conclusion / Speculation

toy model of AdS/CFT

out-of-time ordered correlator



probe of space-time

$$\text{OTO} = \langle A(0)B(t)C(0)D(t) \rangle$$

scrambling / chaos

complexity / design