A gauge theory of measurement-based quantum computation

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Outline

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Part I:

MBQC



The 2D cluster state is a computationally universal "material"

Quantum states can have computational power

Measurement-based quantum computation

Unitary transformation



deterministic, reversible

Projective measurement



probabilistic, irreversible

Measurement-based quantum computation



measurement of Z (\odot), X (\uparrow), $\cos \alpha X + \sin \alpha Y$ (\nearrow)

- Information written onto the resource state, processed and read out by one-qubit measurements only.
- Universal computational resources exist: cluster state, AKLT state.
- R. Raussendorf, H.-J. Briegel, Physical Review Letters 86, 5188 (2001).

Consider a 2D cluster state:



At the end of any MBQC, all that is left is a tensor product state and the measurement record \mathbf{s} .

- The computational output **o** is extracted from the measurement record $\mathbf{s} = (s_1, s_2, .., s_N)$.
- Individually, each bit s_i of measurement record is completely random; *information is only in the correlations*.

Consider a 2D cluster state:



Some cluster qubits simulate the preparation of a quantum register, some its measurement, and some implement quantum gates.

Do these different types of qubits contribute differently to the MBQC output?

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Do these different types of qubits contribute differently to the MBQC output?

No

The classical side-processing relations

Every MBQC has a classical side-processing relation that convert measurement record ${\mbox{s}}$ into

- computational output **o**, and
- the choice **q** of local measurement bases,

$$\mathbf{o} = Z\mathbf{s} \mod 2, \ \mathbf{q} = T\mathbf{s} \mod 2.$$
 (1)

Every bit of computational output is a parity of individual local measurement outcomes.

What determines the support of these parities?

Which measurement outcomes s_i contribute to what output bit o_i ?

The classical side-processing relations

Every MBQC has a classical side-processing relations

 $\mathbf{o} = Z\mathbf{s} \mod 2, \ \mathbf{q} = T\mathbf{s} \mod 2.$

Every bit of computational output is a parity of individual local measurement outcomes.

What determines the support of these parities?

The measurement record \mathbf{s} is a gauge field, and each bit of the computational output \mathbf{o} is a holonomy of the gauge field.

How the side-processing comes about

... it describes how to counteract the randomness of measurement outcomes.



• Propagate the random heralded *byproduct operators* forward in time past the end of the computation.

Part II:

A short history of

"computational phases of quantum matter"



... all across support quantum computation.

Motivation: MBQC and symmetry



Can MBQC schemes be classified by symmetry, in a similar way as, say, elementary particles can?

If so, does this have a bearing on quantum algorithms?

I. Symmetry protects computation





we observe low-maintenance features of the ground-code MQC in that this computation is doable without an exact (classical) description of the resource ground state as well as without an initialization to a pure state. It

It

turns out these features are deeply intertwined with the physics of the 1D Haldane phase (cf. Fig. 1), that is well characterized as the symmetry-protected topological order in a modern perspective [6, 7]. We believe our approach must bring the study of MQC, conventionally based on the analysis of the model entangled states (e.g., [1, 8, 9]), much closer to the condensed matter physics, which is aimed to describe characteristic physics based on the Hamiltonian.

A. Miyake, Phys. Rev. Lett. 105, 040501 (2010).

II. Symmetry-protected wire in MBQC



Else Schwartz Doherty Bartlett



- Computational wire persists throughout symmetry-protected phases in 1D.
- Imports group cohomology from the classification of SPT phases.

D.V. Else, I. Schwartz, S.D. Bartlett and A.C. Doherty, PRL 108 (2012).

F. Pollmann *et al.*, PRB B 81, 064439 (2010); N. Schuch, D. Perez-Garcia, and I. Cirac, PRB 84, 165139 (2011); X. Chen, Z.-C. Gu, and X.-G. Wen, PRB 83, 035107 (2011).



Theorem. Consider a symmetry-protected phase characterized by a finite abelian symmetry group and a maximally noncommutative cohomology class $[\omega]$. In the symmetry-respecting basis $\mathcal{B} = \{|i\rangle\}$, the MPS tensor of a given state in the SPT phase described by $[\omega]$ has the form

$$A[i] = B_i \otimes A_{\mathsf{junk}}[i].$$

The byproduct operators B_i are all unitary and constant throughout the phase.

D.V. Else, I. Schwartz, S.D. Bartlett and A.C. Doherty, PRL 108 (2012).





III. First quantum computational phase



- 1-qubit universal MBQC on a chain of spin-1 particles protected by an S_4 symmetry.
- J. Miller and A. Miyake, Phys. Rev. Lett. 114, 120506 (2015).

IV. The SPT \Rightarrow MBQC meat grinder



A classification of MBQC schemes by symmetry in 1D.

A. Prakash and T.-C. Wei, Phys. Rev. A (2016).

RR, A.Prakash, D.-S. Wang, T.-C. Wei, D.T. Stephen, Phys. Rev. A (2017).

cohomology class $[\omega]$

\downarrow	The given SPT phase determines the byproduct operators [Else et al., RRL 2012]
B_i	
\downarrow	The byproduct operators determine MBQC computational power [RR et al., RRA 2017]

set of gates

Byproduct operators: What entered as a nuisance that could fortunately be dealt with, becomes the central object governing MBQC computational power.

V: Computationally universal SPT phase in 2D



• The 2D cluster state is inside the phase

Result. For a spin-1/2 lattice on a torus all ground states in the 2D cluster phase, except a possible set of measure 0, are universal resources for measurement-based quantum computation.

R. Raussendorf, C. Okay, D.S. Wang, D.T. Stephen, H.P. Nautrup, *A computationally universal quantum phase of matter*, Phys. Rev. Lett. 122, 090501 (2019).

Part III:

A gauge theory of MBQC

The discussion of computational phases of quantum matter is ongoing. It has so far taught us

- The computational capability of MBQC schemes is determined and classified by symmetry.
- To harness this capability, however, a breaking of symmetry is required, namely through the local measurements.

We want a complete picture of the role of symmetry in MBQC. Gauge symmetry seems part of that. Recall the MBQC elevator thought experiment:



All cluster qubits contribute in the same fashion to the computational output, namely through parities.

What determines the support of these parities?

Answer: The measurement record **s** can be regarded as a gauge field, and the computational output **o** corresponds bit-wise to holonomies of that gauge field.

Slightly modified setting—dodging boundaries

cluster state on a chain: concepondes W2 1+> Ιų μĮ



What is implemented on the cluster ring?



What do the two output bits represent?



outcome	occurs with probability
00	$ \alpha ^2$
01	$ \beta ^2$
10	$ \gamma ^2$
11	$ \delta ^2$

How are the two output bits obtained from s?



$$o_1 = \sum_{i \text{ even}} s_i \mod 2, \quad o_2 = \sum_{i \text{ odd}} s_i \mod 2.$$
 (2)

Recap: cluster state stabilizer



Pauli/ stabilizer action on s and q



Gauge transformations

Recall:

$$\begin{array}{rcl} X^{\dagger}O_{\alpha}(q)X &= O_{\alpha}(q \oplus 1), \ \textit{flips q} \\ Z^{\dagger}O_{\alpha}(q)Z &= -O_{\alpha}(q). \ \textit{flips s} \\ \textit{Kj} &= \textit{Zj-1} \ \textit{Xj} \ \textit{Zjr1} \\ \langle \mathcal{C}|O_{\alpha_{i}}^{(i)}(q_{i})|\mathcal{C} \rangle = \langle \mathcal{C}|K_{j}^{\dagger}O_{\alpha_{i}}^{(i)}(q_{i})K_{j}|\mathcal{C} \rangle, \ \forall i,j \ \textit{Insertion f} \ \textit{Kj} \\ \textit{Loss noeffed } \end{cases}$$

Thus we have the following equivalence/ gauge transformations

- Those equivalence transformations are local.
- For j = 2, .., 2N 1 those equivalence transformations preserve the classical processing relations.

Cohomological interpretation



Gauge transformations in cohomological form





standard form

cohomological form

$$K_{j}: \begin{array}{ccc} s_{j-1} & \mapsto & s_{j-1} \oplus 1, \\ k_{j}: & s_{j+1} & \mapsto & s_{j+1} \oplus 1, \\ & q_{j} & \mapsto & q_{j} \oplus 1. \end{array}$$

$$s \mapsto s \oplus d\Lambda, \\ q \mapsto q \oplus \Lambda, \\ \overline{s} \mapsto \overline{s} \oplus d\overline{\Lambda}, \\ \overline{q} \mapsto \overline{q} \oplus \overline{\Lambda}.$$

Gauge-invariant quantities

Find a 1-chain e such that, for all λ

$$s(e) \mapsto [s \oplus d\Lambda](e) = s(e).$$

Requires

$$d\Lambda(e)=0.$$

Hence,

$$0 = d\Lambda(e) = \Lambda(\partial e), \ \forall \Lambda,$$

and thus

 $\partial e = 0.$

The 1-chain e is, in fact, a 1-cycle.

- s(e) is the corresponding holonomy.
- s(e) is gauge invariant.
- Have the same invariant $\overline{s}(\overline{e})$ on the dual complex \overline{C} .





Gauge-invariant quantities

- s(e) is gauge invariant.
- Have the same invariant $\overline{s}(\overline{e})$ on the dual complex \overline{C} .

These holonomies have previously been identified as the computational output. Recall:



The MBQC "elevator thought experiment"

Consider a 2D cluster state:



Outlook

🕖 Measurmmt -3 Gauge heary of based guantum computation MBQC (nBQC) 2 Computation OnKook phases of quantum maker '(SPTO)

* Extend gange-Merchical description from bixed-point state to surrounding SPT phase - Problem: Can nolonger use the stabilizers - Gord sign: Classical proceeding relations extend mo the SPT phase unchanged

- In 1D so far, MBQC can be understood in terms of gauge theory, with the measurement record as the gauge field.
- In that picture, the MBQC computational output arises as holonomies of the gauge field.
- Extension to computational phases of quantum matter is suggestive.

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