

21 Sep. 2023

Clifford Group and Unitary Designs under Symmetry

arXiv:2306.17559

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Motivation for our work

What emerges from the interplay of randomness and symmetry?



Randomness

- ✓ Investigation in fundamental physics
e.g. Thermalization
- ✓ Useful in quantum algorithms
e.g. Randomized benchmarking

Symmetry

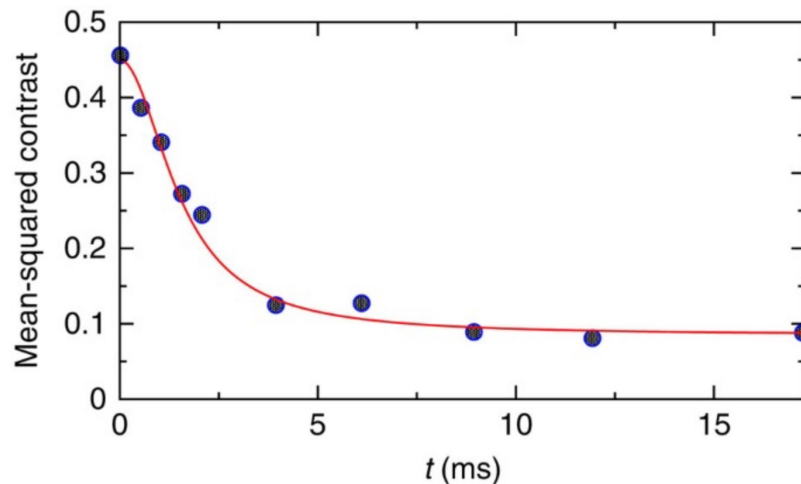
- ✓ Efficient analysis of systems
e.g. Integrability
- ✓ Opens new physics
e.g. SPT phase

Uniform random unitary

Generalized notion of uniform distribution on the unitary group

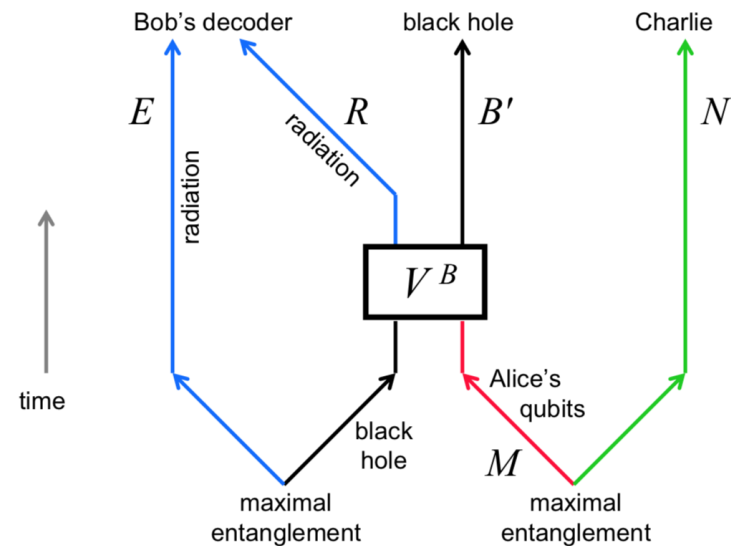
✓ Appears in many fields of fundamental physics.

Thermalization ← Typicality



P. Reimann, Nat. Commun. 7, 10821 (2016).

Black hole dynamics



P. Hayden and J. Preskill, J. High Energy Phys. 09 (2007) 120.

Unitary t -design

Simulating Haar random unitary is a hard task.

➔ We need a set of unitary operators approximating uniform random unitary.

A compact subgroup $\mathcal{X} \subset \mathcal{U}_N$ is a unitary t -design if

$$\Phi_{t,\mathcal{X}} = \Phi_{t,\mathcal{U}_N},$$

← Unitary group on N qubits

where $\Phi_{t,\mathcal{X}}$ is the t -fold twirling channel over \mathcal{X} :

$$\Phi_{t,\mathcal{X}}(L) := \int_{U \in \mathcal{U}_N} U^{\otimes t} L U^{\dagger \otimes t} d\mu_{\mathcal{X}}(U) \quad \forall L \in \mathcal{L}_{tN}.$$

↑
Haar measure on \mathcal{X}

↑
all linear operators on tN qubits

Unitary t -designs for large t are desirable.

Applications

e.g. Noise twirling

J. Emerson *et al.*, J. Opt. B: Quantum
Semiclass. Opt. **7**, S347-S352 (2005).

Randomized benchmarking

C. Dankert *et al.*,
Phys. Rev. A **80**, 012304 (2009).

} Unitary 2-designs are needed.

Partial quantum tomography

H.-Y. Huang, R. Kueng, and J. Preskill,
Nat. Phys. **16**, 1050-1057 (2020).

← Unitary 3-designs are needed.

Unitary designs play essential roles
in various quantum information processing tasks.

Clifford group

Pauli group $\mathcal{P}_N := \{\pm 1, \pm i\} \cdot \{I, X, Y, Z\}^{\otimes N}$

$$I = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \quad X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \quad Y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix} \quad Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

Clifford group $\mathcal{C}_N := \{U \in \mathcal{U}_N \mid U\mathcal{P}_N U^\dagger = \mathcal{P}_N\}$

Hadamard, S, CNOT $\in \mathcal{C}_N$

✓ Quantum circuits with Clifford gates can be simulated classically.

D. Gottesman, arXiv:quant-ph/9807006.

✓ The Clifford group is a unitary 3-design.

Z. Webb, Quantum Inf. Comput. **16**, 1379 (2016).

H. Zhu, Phys. Rev. A **96**, 062336 (2017).

Symmetry in physics

A physical system has symmetry $\mathcal{G} \iff \forall G \in \mathcal{G} [H, G] = 0$.

H : Hamiltonian

Transverse-field Ising model

$$H = g \sum X_j X_k + h \sum Z_j$$

$$\mathcal{G} = \{I^{\otimes N}, Z^{\otimes N}\} \cong \mathbb{Z}_2$$

Heisenberg XXZ model

$$H = \sum [g(X_j X_k + Y_j Y_k) + g' Z_j Z_k]$$

$$\mathcal{G} = \{(e^{i\theta Z})^{\otimes N} \mid \theta \in \mathbb{R}\} \cong U(1)$$

Heisenberg XXX model

$$H = g \sum (X_j X_k + Y_j Y_k + Z_j Z_k)$$

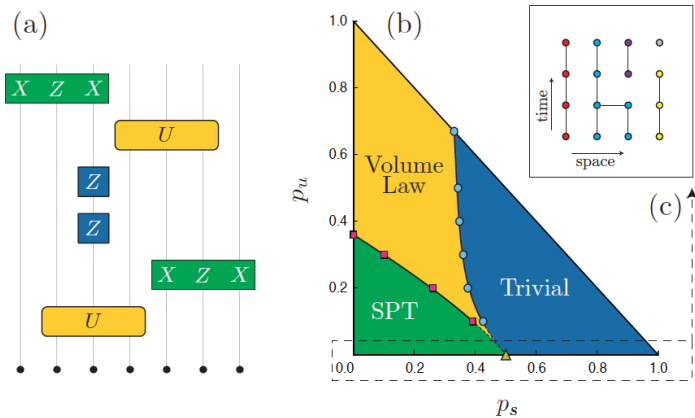
$$\mathcal{G} = \{(e^{i(\theta_X X + \theta_Y Y + \theta_Z Z)})^{\otimes N} \mid \theta_\alpha \in \mathbb{R}\} \cong SU(2)$$

✓ Useful for efficient analysis of the systems

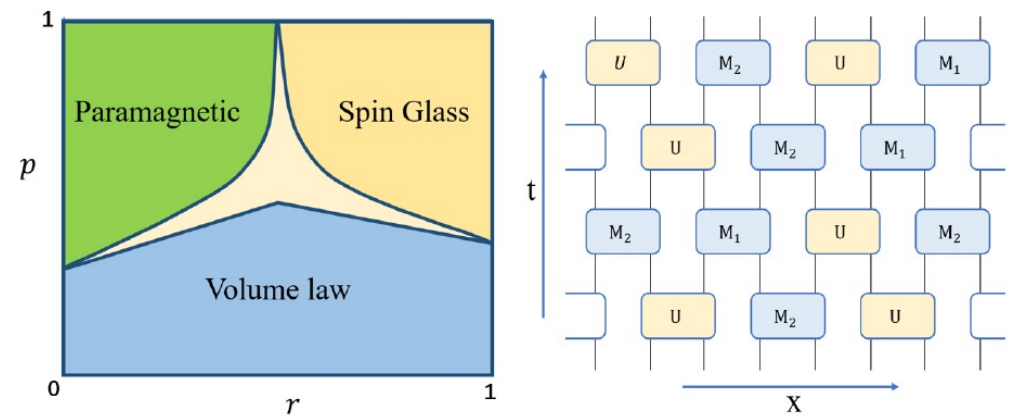
Symmetry in physics

Symmetry also enriches physics.

e.g. In measurement-induced phase transitions, interesting features emerge by introducing symmetry constraints.



A. Lavasani, Y. Alavirad, and M. Barkeshli,
Nat. Phys. **17**, 342-347 (2021).



S. Sang and T. H. Hsieh,
Phys. Rev. Research **3**, 023200 (2021).

Core questions in our work

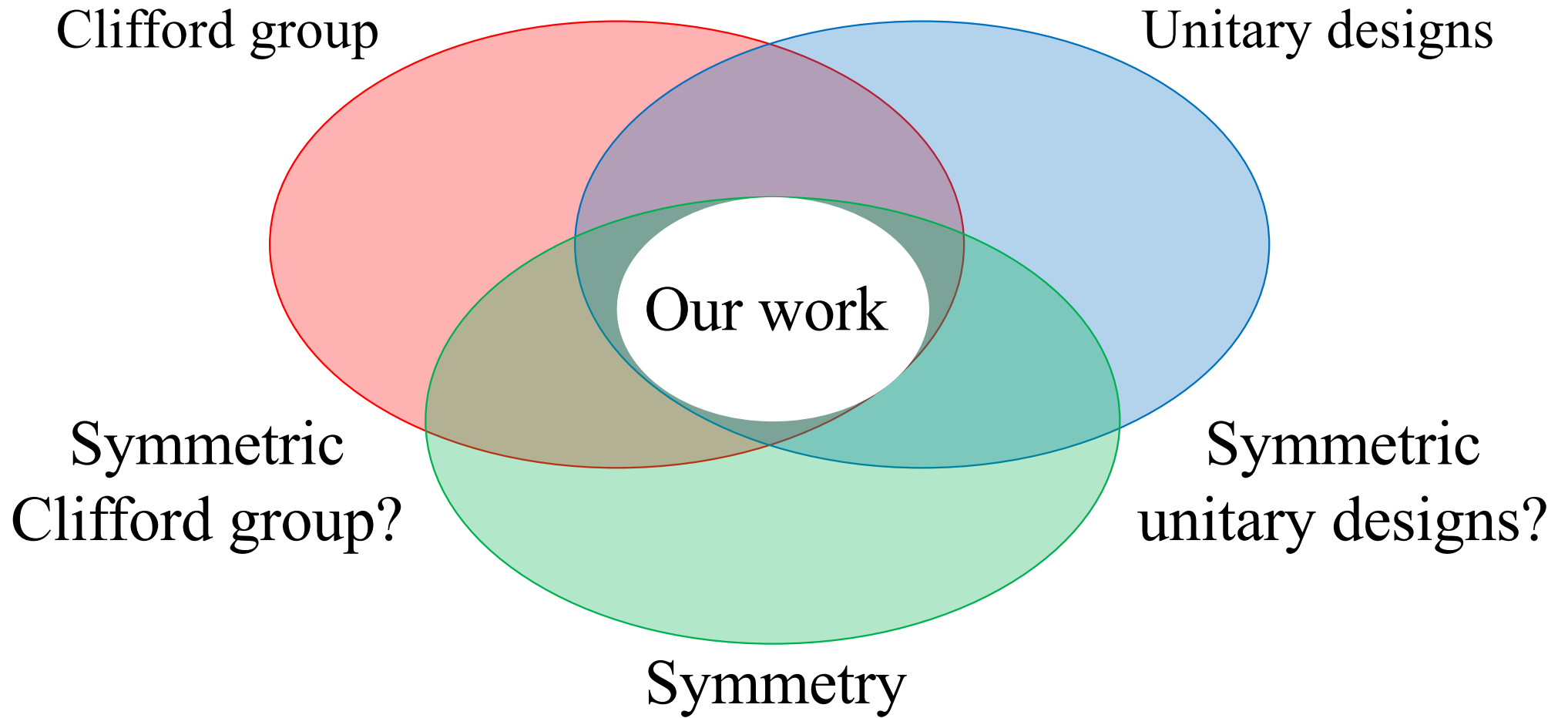
1. Can we characterize the pseudo-randomness of the Clifford group under symmetry?

 First result

2. What is the explicit construction?

 Second result

Setup

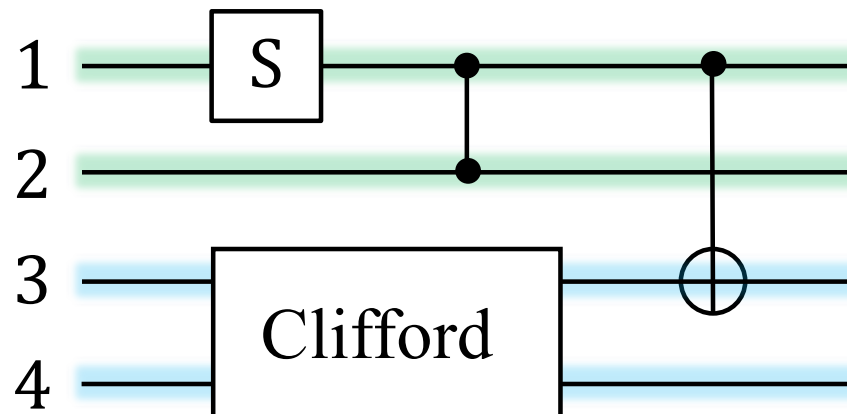


Symmetric Clifford group

\mathcal{G} -symmetric unitary group $\mathcal{U}_{N,\mathcal{G}} := \{U \in \mathcal{U}_N \mid \forall G \in \mathcal{G} [U, G] = 0\}$

\mathcal{G} -symmetric Clifford group $\mathcal{C}_{N,\mathcal{G}} := \mathcal{C}_N \cap \mathcal{U}_{N,\mathcal{G}}$

e.g. \mathcal{G} -symmetric Clifford gates for $\mathcal{G} = \{I, Z_1, Z_2, Z_1Z_2\}$



Symmetric unitary t -design

Unitary designs approximate uniform distribution on \mathcal{U}_N .



\mathcal{G} -symmetric unitary designs approximate uniform distribution on $\mathcal{U}_{N,\mathcal{G}}$.

$$\mathcal{U}_{N,\mathcal{G}} := \{U \in \mathcal{U}_N \mid \forall G \in \mathcal{G} [U, G] = 0\}$$

A compact subgroup $\mathcal{X} \subset \mathcal{U}_N$ is a \mathcal{G} -symmetric unitary t -design if

$$\Phi_{t,\mathcal{X}} = \Phi_{t,\mathcal{U}_{N,\mathcal{G}}}. \quad \Phi_{t,\mathcal{X}}(L) := \int_{U \in \mathcal{U}_N} U^{\otimes t} L U^{\dagger \otimes t} d\mu_{\mathcal{X}}(U)$$

The standard unitary designs are included as the case of $\mathcal{G} = \{I\}$.

Problem

The Clifford group is a unitary 3-design.

Is the \mathcal{G} -symmetric Clifford group
a \mathcal{G} -symmetric unitary 3-design?

Does it always hold?

If not, for what symmetry \mathcal{G} does it hold?

1st result: randomness of symmetric Clifford group

\mathcal{G} : subgroup of \mathcal{U}_N

$\mathcal{C}_{N,\mathcal{G}}$ is a \mathcal{G} -symmetric unitary 3-design

if and only if \exists subgroup $\mathcal{Q} \subset \mathcal{P}_N$ s.t. $\mathcal{U}_{N,\mathcal{G}} = \mathcal{U}_{N,\mathcal{Q}}$

Symmetry constraint is described by some Pauli subgroup.

$$\mathcal{C}_{N,\mathcal{G}} := \mathcal{C}_N \cap \mathcal{U}_{N,\mathcal{G}}$$

\mathcal{C}_N : the Clifford group

$$\mathcal{U}_{N,\mathcal{G}} := \{U \in \mathcal{U}_N \mid \forall G \in \mathcal{G} [U, G] = 0\}$$

✓ Characterized Pauli-symmetric Clifford group in terms of unitary designs

Examples

Pauli subgroup symmetry

$$\mathcal{G} = \{I^{\otimes N}, Z^{\otimes N}\} \cong \mathbb{Z}_2$$

→ $\mathcal{C}_{N,\mathcal{G}}$ is a \mathcal{G} -symmetric unitary 3-design.

(but not a 4-design)

Non-Pauli subgroup symmetry

$$\begin{cases} \mathcal{G} = \{(e^{i\theta Z})^{\otimes N} \mid \theta \in \mathbb{R}\} & \cong U(1) \\ \mathcal{G} = \{(e^{i(\theta_X X + \theta_Y Y + \theta_Z Z)})^{\otimes N} \mid \theta_X, \theta_Y, \theta_Z \in \mathbb{R}\} & \cong SU(2) \end{cases} \quad \text{for } N \geq 2$$

→ $\mathcal{C}_{N,\mathcal{G}}$ is not a \mathcal{G} -symmetric unitary 3-design.

(a 1-design but not a 2-design)

Proof sketch

Theorem S1. $\mathcal{C}_{N,\mathcal{G}}$ is a \mathcal{G} -symmetric unitary 3-design iff $\exists Q \subset \mathcal{P}_N$ s.t. $\mathcal{U}_{N,\mathcal{G}} = \mathcal{U}_{N,Q}$.

“if” part

Proposition S1. $\mathcal{C}_{N,Q}$ is a Q -symmetric unitary 3-design for $Q \subset \mathcal{P}_N$.

Lemma S1. $\forall \mathcal{D} \in \mathfrak{C}_{t,Q}, \Phi_{3,\mathcal{X}} \circ \mathcal{D} = \Phi_{3,\mathcal{X}}$ for $Q \subset \mathcal{P}_N, \mathcal{X} = \mathcal{C}_{N,Q}$ or $\mathcal{U}_{N,Q}$.

Lemma S2. $\forall L \in \mathcal{M}_Q, \Phi_{3,\mathcal{X}}(L) = L$ for $Q \subset \mathcal{P}_N$.

Lemma S3. $\exists \mathcal{D} \in \mathfrak{C}_{t,Q}$ s.t. $\forall L \in \mathcal{L}(\mathcal{H}^{\otimes 3}), \mathcal{D}(L) \in \mathcal{M}_Q$ for $Q \subset \mathcal{P}_N$.

← Lemmas S4, S16

← Lemma S14 ← Lemmas S12, S13 ↘ Lemma S15

Lemma S5.

S gates

Lemma S6.

CZ gates

Lemma S7.

Clifford gates

Lemma S8.

Controlled-Pauli gates

“only if” part

Proposition S2. If $\mathcal{C}_{N,\mathcal{G}}$ is a \mathcal{G} -symmetric unitary 3-design, then $\mathcal{U}_{N,\mathcal{G}} = \mathcal{U}_{N,Q}$, where $Q = \{\{Q \in \mathcal{P}_N^+ \mid \exists G \in \mathcal{G} \text{ s.t. } \text{tr}(GQ) \neq 0\}\}$.

Lemma S9. If $L = \sum_{P \in \mathcal{P}_N^+} \gamma_P(G) P^{\otimes 3}$, then $\forall U \in \mathcal{C}_{N,\mathcal{G}}, U^{\otimes 3} L U^{\dagger \otimes 3} = L$.

← Lemma S17

Lemma S10. If $\mathcal{C}_{N,\mathcal{G}}$ is a \mathcal{G} -symmetric unitary 3-design and $\forall U \in \mathcal{C}_{N,\mathcal{G}}, U^{\otimes 3} L U^{\dagger \otimes 3} = L$, then $\forall U \in \mathcal{U}_{N,\mathcal{G}}, U^{\otimes 3} L U^{\dagger \otimes 3} = L$.

Lemma S11. If $L = \sum_{P \in \mathcal{P}_N^+} \gamma'_P P^{\otimes 3}$ and $\forall U \in \mathcal{U}_{N,\mathcal{G}}, U^{\otimes 3} L U^{\dagger \otimes 3} = L$, then $\forall U \in \mathcal{U}_{N,\mathcal{G}}$ and $P \in \mathcal{P}_N^+$ satisfying $\gamma'_P \neq 0, U P U^\dagger = P$.

← Lemma S18

“if” part

Suppose $\mathcal{U}_{N,\mathcal{G}} = \mathcal{U}_{N,Q}$.

We explicitly construct a channel \mathcal{D} s.t.

$$\mathcal{D} = \Phi_{3,\mathcal{X}} \text{ and } \mathcal{D} = \Phi_{3,\mathcal{U}_{N,\mathcal{G}}}$$

$$\rightarrow \Phi_{3,\mathcal{X}} = \Phi_{3,\mathcal{U}_{N,\mathcal{G}}}$$

“only if” part

Suppose $\mathcal{C}_{N,\mathcal{G}}$ is a \mathcal{G} -symmetric unitary 3-design.

$$\text{e.g. } \mathcal{G} = \left\{ I^{\otimes 3}, \frac{X^{\otimes 3} + Z^{\otimes 3}}{\sqrt{2}} \right\}.$$

$$\rightarrow \mathcal{U}_{N,\mathcal{G}} = \mathcal{U}_{N,Q}.$$

$$\text{with } Q = \{\pm 1, \pm i\} \cdot \{I^{\otimes 3}, X^{\otimes 3}, Y^{\otimes 3}, Z^{\otimes 3}\}$$

2nd result: construction of symmetric Clifford gates

Q : Pauli subgroup

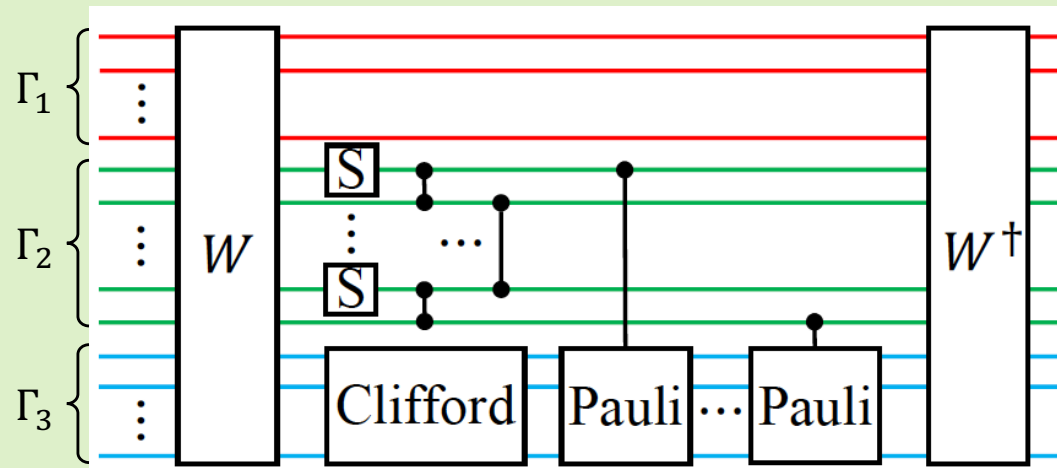
$$\exists W \in \mathcal{C}_N \text{ s.t. } W Q W^\dagger = \{I, X, Y, Z\}^{\otimes N_1} \otimes \{I, Z\}^{\otimes N_2} \otimes \{I\}^{\otimes N_3} \quad (\text{up to phase})$$

Γ_1

Γ_2

Γ_3 indices for qubits

Every Q -symmetric Clifford gate can be uniquely expressed as:



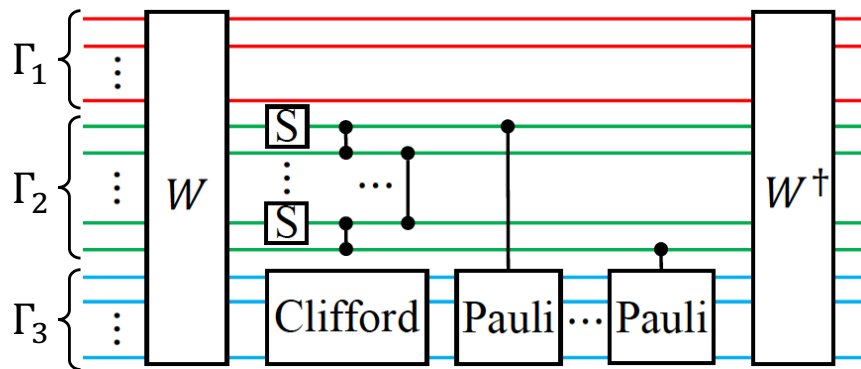
$$W^\dagger \left(\prod_{j \in \Gamma_2} \prod_{k \in \Gamma_3} C^{(P_{j,k})_{j,k}} \right) V$$

$$\times \left(\prod_{j,k \in \Gamma_2, j < k} CZ_{j,k}^{v_{j,k}} \right) \left(\prod_{j \in \Gamma_2} S_j^{\mu_j} \right) W$$

with $\mu_j \in \{0, 1, 2, 3\}$, $v_{j,k} \in \{0, 1\}$,
 $V \in \mathcal{C}_{N_3}$, $P_{j,k} \in \{I, X, Y, Z\}$.

✓ Efficient method for generating all the symmetric Clifford gates

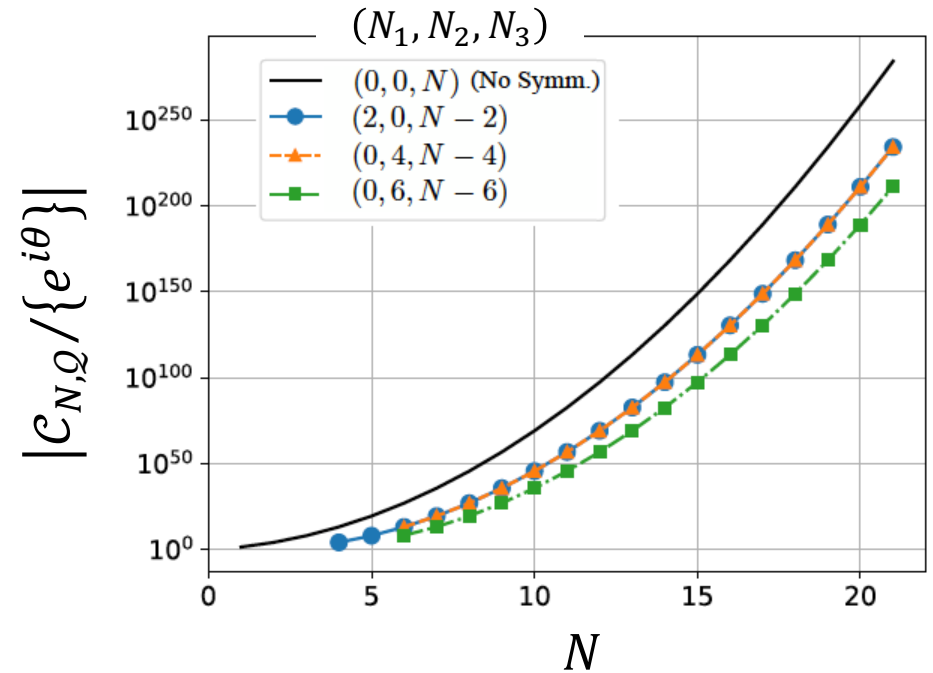
Size of symmetric Clifford gates



$$|C_{N,Q}/\{e^{i\theta}\}| \sim 2^{2(N_3+N_2/2)^2+3(N_3+N_2/2)}$$

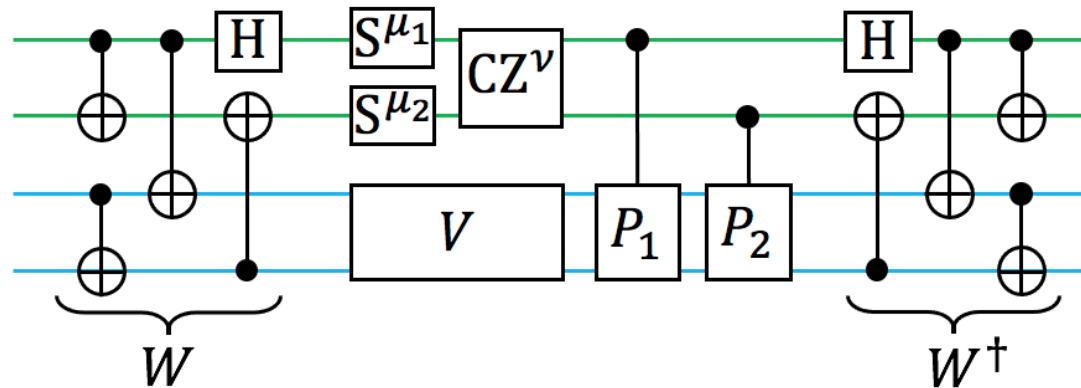
$$\sim |C_{N_3+N_2/2}/\{e^{i\theta}\}|$$

$|C_{N,Q}/\{e^{i\theta}\}|$ is super-exponential,
but it is exponentially smaller than $|C_N/\{e^{i\theta}\}|$.



Example

For $Q = \{I^{\otimes 4}, X^{\otimes 4}, Y^{\otimes 4}, Z^{\otimes 4}\}$, every $U \in \mathcal{C}_{4,Q}$ is uniquely expressed as:



with $\mu_j \in \{0, 1, 2, 3\}$, $v_{j,k} \in \{0, 1\}$, $V \in \mathcal{C}_2$, $P_j \in \{I, X, Y, Z\}^{\otimes 2}$.

$$|\mathcal{C}_{4,Q}/\{e^{i\theta}\}| \sim 10^8, |\mathcal{C}_4/\{e^{i\theta}\}| \sim 10^{13}$$

➔ Generating $\mathcal{C}_{4,Q}$ by this method is 10^5 times more efficient than generating \mathcal{C}_4 and choosing symmetric gates.

Proof

Sufficient to consider $Q = \mathcal{P}_0\{I, X, Y, Z\}^{\otimes N_1} \otimes \{I, Z\}^{\otimes N_2} \otimes \{I\}^{\otimes N_3}$.

Completeness Take the Heisenberg picture.

U is Q -symmetric

$$U \cdot U^\dagger$$

$$j \in \Gamma_1$$

$$Z_j \mapsto Z_j$$

$$X_j \mapsto X_j$$

$$j \in \Gamma_2$$

$$Z_j \mapsto Z_j$$

$$X_j \mapsto ?$$

inductively construct U'
with the 4 types of gates s.t.

$$U'U \cdot U^\dagger U'^\dagger$$

$$Z_j \mapsto Z_j$$

$$X_j \mapsto X_j$$

$$Z_j \mapsto Z_j$$

$$X_j \mapsto X_j$$

➔ $U'U$ is a Clifford gate on the third subsystem.

➔ U can be constructed by the 4 types of gates.

Uniqueness

Apply U to states $|\Psi\rangle \otimes \left(\otimes_j |x_j\rangle\right) \otimes |\mathbb{E}\rangle$
s.t. the number of nonzero x_j is at most 2.

$|x_j\rangle$: Z-basis

Summary and future perspectives

Summary

arXiv:2306.17559

Proved the \mathcal{G} -symmetric Clifford group is a \mathcal{G} -symmetric unitary 3-design if and only if the symmetry constraint is described by some Pauli subgroup.

Found a complete and unique construction of the symmetric Clifford groups.

Perspectives

New direction in quantum information processing tasks under symmetry?

Foundation for understanding to many-body systems?