Quantum Error Correction Threshold and von Neumann Algebra Types

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Threshold of Quantum Error Correction Codes

- Errors occur during the storage, computation, error correction and decoding processes.
- Only codes with error threshold could be used for large scale and long time quantum computing. [Aharonov, Ben-Or 99, Knill; Laflamme, Zurek, 98; Kitaev 97]

$$\lim_{n \to \infty} \Pr(\text{decoding success}) = \begin{cases} 1, & p < p_0 \\ 1/D, & p > p_0 \end{cases}$$



Thresholds are difficult to extract

- Concatenation of 3-repetition code [Aliferis, Gottesman, Preskill 05]
 Recursion relation between logical errors on different layers of concatenation.
- Toric code or planar code [Dennis, Kitaev, Landahl, Preskill 01; Chubb, Flammia 19]
 Map to a random bond Ising model and study its phase transition.
- Oscillator to oscillator codes do not have a threshold. [Hänggli, König 22]
- Can we try to have a better understanding of the underlying structure of the existence of the threshold of codes?



von Neumann algebras

We provide a plausible conjecture that the existence of threshold is related to von Neumann algebra change.

- von Neumann algebra \mathcal{A} is a set of *-algebra of bounded operators acting on a Hilbert space and closed under WOT. $\mathcal{A} = \mathcal{A}^{"}$.
- Relevant to our purpose: hyperfinite algebra

$$\mathcal{A}_{\Gamma} = \overline{\cup_{\Lambda_0 \in \Gamma} \mathcal{A}_{\Lambda}}, \quad \mathcal{A}_{\Lambda_0} = \otimes_{x \in \Lambda_0} \mathcal{B}(\mathcal{H}_X).$$

- Γ: the whole lattice; Λ₀: a finite region; C: an infinite cone-like subregion
- The type depends on the state used when taking the closure and encodes the entanglement structure of the states in the Hilbert space.
- The type shows the properties of the projectors in the place prage magnitude

More motivations

- von Neumann algebra is the natural mathematical language of systems of infinite degrees of freedom.
- There are examples in which algebra type can change depending on one parameter of the state. [Witten 22]

$$\Psi_{\mathsf{TFD}} = \frac{1}{\sqrt{Z(\beta)}} \bigotimes_{n=1}^{\infty} (|0_L 0_R\rangle + n^{-\beta} |1_L 1_R\rangle), \ Z(\beta) = \prod_{n=1}^{\infty} (1+n^{-\beta}).$$

When $\beta > 1$, $Z(\beta) < \infty$, A_L is type I_{∞} ; when $\beta < 1$, $Z(\beta) > \infty$, A_L is type III₁.

 On ferromagnetic and antiferromagnetic phases of Ising model on a Bethe tree, the algebra types are different. [Mukhamedov 00]



More motivations from gravity

- There exists a Hawking-Page transition from thermal AdS space-time to AdS black hole.
- The types of algebras in thermal AdS and AdS black hole are conjectured to be Type I and Type III respectively. [Leutheusser, Liu 22]
- This transition is dual to the confinement/deconfinement phase transition in the dual large N gauge theory. [Witten 98]
- The holographic error correction code is microscopic model for AdS/CFT duality of quantum gravity. [Pastawski, Yoshida, Harlow, Preskill 15]
- In the holographic error correction code, this transition is conjectured to dual to the threshold transition . [Bao, Cao, Zhu 22]
- These facts suggests that a good candidate of holographic code should have threshold and the type change should in the gravity.

States in the logical subspace?

- Stabilizer formalism gives logical subspace as ground states of a gapped Hamiltonian and they are pure states.
- If consider the algebra of operators on the full lattice constructed on this pure state, it is always a type I factor.
- How about some infinite subregion on the lattice?
- Trivial example: for repetition code, it is again type I;
- For any codes that can be realized by a 1D gapped Hamiltonian ground state, the algebra of its infinite half line subregion is type I. [Matsui 13]
- Non-trivial example: the algebra of a cone-shaped region of planar code ground state is type II_∞.[Ogata 22]



Which state? Which algebra?

States in the logical subspace dressed by error channel?

- Let the error channel on one lattice site x be $\varepsilon_x(\cdot) = \sum_i K_i(\cdot)K_i^{\dagger}$ and the total error channel be $\varepsilon_{\Gamma}(\cdot) = \bigotimes_{x \in \Gamma} \varepsilon_x(\cdot)$
- Consider the state ω ∘ ε_Γ(·), where ω is a pure state in the logical subspace.
- For example, when the error model is the bit flip error $\varepsilon_x(\cdot) = (1-p)I(\cdot) + pX(\cdot)X$, $0 , the algebra <math>\mathcal{A}_{\Gamma}$ is type III_{λ} with $\lambda = \frac{p}{1-p}$. If p = 1/2, \mathcal{A}_{Γ} is type II₁. [Araki, Woods 68]
- How about considering \mathcal{A}_C with state $\omega_C \circ \varepsilon_C(\cdot)$?
- For 1D, it is essentially the same case as ω ∘ ε_Γ(·). For 2D planar code, we are not sure but suspect it is type III for generic p.

The thermal state in the different phases of the dual statistical model? [Dennis, Kitaev, Landahl, Preskill 01; Chubb, Flammia 19]

- For each error configuration, there is an Hamiltonian H_E such that its partition function corresponds to the probability of the error configuration class conjugated by stabilizers. $Z_E = \Pr(\bar{E}).$
- For toric code, this is the random bond Ising model.
- The threshold corresponds to a ordered/disordered phase transition in this model.
- An order parameter is the quenched free energy difference between error *E* and *EL_m*, *L_m* is a non-trivial logical operator.
- The free energy difference diverges below the threshold. When above threshold, the configurations where it diverges become negligible.

We may take a thermal state of a given H_E , then averaging over different error configurations

- $\sum_{E} \Pr(E) \operatorname{tr}(e^{-\beta H_{E}}(\cdot)) = \langle \langle (\cdot) \rangle_{\beta} \rangle_{p}$.
- Hint that this might be the correct choice:
- In the disordered phase,

$$\lim_{r\to\infty}\langle\langle\sigma_o\sigma_r\rangle_\beta\rangle_p=0.$$

- This is a necessary condition of the split property of the state in on two different regions. If the global state were pure, this ensures that the algebra on the cone-like subregion is type I.
- Still far from a complete proof!



Summary and outlook

- We provide a conjecture that the existence of the threshold of code implies the type of algebra changes along a one parameter family of states related to the code.
- We give some plausible arguments supporting this conjecture.
- We provide some trial states constructions.
- We hope this new probe of von Neumann algebra type could help to characterize different types of threshold transition.
- If the conjecture is (partially) true, we hope this may guide the search of realistic holographic codes.



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Thank you!



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